

# CS 113 – Computer Science I

# Lecture 23 – Sorting

Adam Poliak 04/18/2023

# Announcements

#### • HW08:

- Due 04/21 will be released by Friday
- Inheritance and interfaces fully autograded
- HW09:
  - Due 04/28
  - Building a fancyArray class fully autograded

#### Midterm

- Overall class did well:
  - Median 79% (Midterm 1: 77%)
  - Mean 71% (Midterm 1: 68%)

- When grading we'll lower the maximum
  - Instead of being out of 77, the denominator will be lower
- You can make a regrade request on Gradescope

#### Course evaluations

What do you see as the major strengths of Adam Poliak in this course? What areas do you see for improvement in instruction and/or in content?

How prepared were you to take this course? What courses, if any, would you have found useful to take before this course? Is this course listed at the appropriate level?

How did Adam Poliak effectively create an accessible and inclusive course experience? What areas do you see for commendation and/or improvement in the instructor's attention to accessibility and inclusivity?

Would you recommend this course, as taught by Adam Poliak, to other students? Why or why not?

Title	Unique ID	Instructor	Enrollments	Responded	Response Rate
Computer Science I - bmc.CMSC.B113.001.SP23	bmc.CMSC.B11 3.001.SP23	Adam Poliak	32	2	6.25%
Computer Science I - bmc.CMSC.B113.00A.SP23	bmc.CMSC.B11 3.00A.SP23	Adam Poliak	32	4	12.5%

# Outline

- Announcements
- Sorting two algorithms
  - BubbleSort
  - SelectionSort

# Who do we care about sorting?

Makes searching faster!

Efficient sorting is important for optimizing the efficiency of other algorithms (such as search and merge algorithms) that require input data to be in sorted lists. Sorting is also often useful for canonicalizing data and for producing humanreadable output.



# ChatGPT's response:

Covered in data structures, next semester!



Sorting is a fundamental operation in computer science because it is a common and essential task in many applications. Sorting refers to arranging a collection of elements in a particular order, such as numerical, alphabetical, chronological, or any other custom order. Here are a few reasons why we care about sorting in computer science:

- Searching: Sorted data is easier to search. Searching an unsorted list requires examining every element in the list, which can be slow and inefficient for large data sets. In contrast, searching a sorted list can be done using efficient algorithms like binary search that require examining only a logarithmic number of elements.
- 2. Data analysis: Sorting can help in data analysis. Sorting data can help identify patterns and relationships that may not be immediately apparent in an unsorted data set. For example, sorting a list of sales data by date can reveal trends and seasonality.
- 3. Optimization: Sorting is often a pre-requisite for other algorithms, such as finding the median or calculating the mode of a data set. Sorting can also be used as a building block for other algorithms, such as merge sort or quick sort, which are used in many

computational applications.

4. Database management: Sorting is important in managing databases. Databases often contain large amounts of data that need to be sorted in order to perform operations such as searching, filtering, and data analysis.

In summary, sorting is an essential operation in computer science that helps in searching, data analysis, optimization, and database management. It is a basic building block for many algorithms and is crucial for efficient and effective computation.

# Sorting

#### How might we sort the list of numbers below. Can we come up with an algorithm?

0	1	2	3	4	5
10	4	3	0	11	8

Compare two adjacent items, and swap if needed

Repeat until largest item is at the back

Repeat process until done

0	1	2	3	4	5
10	4	3	0	11	8

#### What do we do first?

0	1	2	3	4	5
10	4	3	0	11	8



j - 1



j - 1

 $\bigcap$ 

len = 6



Compare j-1 and j; Swap if L[j-1] > L[j]

len = 6

0	1	2	3	4	5
4	10	3	0	11	8



j - 1

1

j

2

len = 6



Compare j-1 and j; Swap if L[j-1] > L[j]



_	0	1	2	3	4	5
	4	3	10	0	11	8
			Î	Î	<u>-</u>	
			j - 1	j		
			2	3		

Compare j-1 and j; Swap if L[j-1] > L[j]

len = 6

0	1	2	3	4	5
4	3	0	10	11	8
		Î Î	Î		
		j - 1 2	j 3		
			1 、 」 [:]		

Compare j-1 and j; Swap if L[j-1] > L[j]

len = 6



#### Compare j-1 and j; Swap if L[j-1] > L[j]

len = 6

0	1	2	3	4	5
4	3	0	10	11	8

#### Compare j-1 and j; Swap if L[j-1] > L[j]

What next?

j-1 j 4 5

len = 6

0	1	2	3	4	5
4	3	0	10	8	11

#### Compare j-1 and j; Swap if L[j-1] > L[j]

What next?

j 5

j - 1

4

Bubble Sort	
-------------	--

j - 1

 $\mathbf{O}$ 



0	1	2	3	4	5
4	3	0	10	8	11
Î			L	ast elemer	nt has

Reset and compare pairs with shorter list!

j

1

What next?

largest element!





What next?

len = 5





What next?

len = 5

# Bubble Sort len = 5 0 1 2 3 4 5 3 0 4 10 8 11

j - 1

3

#### Compare j-1 and j; Swap if L[j-1] > L[j]



0	1	2	3	4	5
3	0	4	8	10	11
			f j - 1 3	1 j 4	

#### Compare j-1 and j; Swap if L[j-1] > L[j]



j - 1

 $\mathbf{O}$ 



0	1	2	3	4	5
3	0	4	8	10	11

Reset and check pairs with shorter list

j

1

j - 1

 $\bigcap$ 





# Compare j-1 and j; Swap if L[j-1] > L[j]



0	1	2	3	4	5
0	3	4	8	10	11

#### Compare j-1 and j; Swap if L[j-1] > L[j]

j - 1

1

j

2



0	1	2	3	4	5
0	3	4	8	10	11
					,
		j - 1 2	j 3		

Compare j-1 and j; Swap if L[j-1] > L[j]





0	1	2	3	4	5
0	3	4	8	10	11









1

j - 1



j - 1

 $\mathbf{O}$ 



0	1	2	3	4	5
0	3	4	8	10	11

Reset; Compare j-1 and j; Swap if L[j-1] > L[j]

j

1

Idea: bubble highest values to the end of the list; Check a shorter sublist each time

bubbleSort(L):

```
for len in range(len(L), 1, -1):
    for j in range(1, len): # bubble
    if L[j-1] > L[j]:
        swap(j-1, j, L)
```





Repeatedly find the smallest item and put it at front of list

selectionSort(L):

```
for startIdx in range(len(L)):
```

```
minIdx = findMinimum(startIdx, L)
```

```
swap(startIdx, minIdx, L)
```

0	1	2	3	4	5
10	4	3	0	11	8

#### What do we do first?



Find minimum element idx between start to end



Swap the elements at start and minIdx



#### start 1

Decrease the interval.



start minIdx 1 2

Find minimum element between start to end



Swap the elements at start and minIdx



start 2

#### Decrease the interval.



Find minimum element idx between start to end



Swap the elements at start and minIdx



3

#### Decrease the interval.



Find minimum element idx between start to end



Swap the elements at start and minIdx

0	1	2	3	4	5
0	3	4	8	11	10

Decrease the interval.

What next?

start

4



#### Find minimum element idx between start to end



#### Swap the elements at start and minIdx

0	1	2	3	4	5
0	3	4	8	10	11

start 5

Decrease the interval.

We're done!

```
findMinimum(startIdx, L):
```

```
minIdx = startIdx
```

```
for i in range(startIdx, len(L)):
    if L[i] < L[minIdx]:
        minIdx = i</pre>
```

return minIdx

# Swap

swap(i, j, L):
 temp = L[i] # step 1
 L[i] = L[j] # step 2
 L[j] = temp # step 3



## Selection sort and Bubble sort are O(N<sup>2</sup>)

