CS 113 - Computer
Science I

Lecture 23 - Sorting

Adam Poliak
04/18/2023

## Announcements

- HW08:
- Due 04/21 - will be released by Friday
- Inheritance and interfaces - fully autograded
- HW09:
- Due 04/28
- Building a fancyArray class - fully autograded


## Midterm

- Overall class did well:
- Median 79\% (Midterm 1: 77\%)
- Mean 71\% (Midterm 1: 68\%)
- When grading we'll lower the maximum
- Instead of being out of 77, the denominator will be lower
- You can make a regrade request on Gradescope


## Course evaluations

What do you see as the major strengths of Adam Poliak in this course? What areas do you see for improvement in instruction and/or in content?

How prepared were you to take this course? What courses, if any, would you have found useful to take before this course? Is this course listed at the appropriate level?

How did Adam Poliak effectively create an accessible and inclusive course experience? What areas do you see for commendation and/or improvement in the instructor's attention to accessibility and inclusivity?

Would you recommend this course, as taught by Adam Poliak, to other students? Why or why not?

## Outline

- Announcements
- Sorting - two algorithms
- BubbleSort
- SelectionSort


## Who do we care about sorting?

## Makes searching faster!

Efficient sorting is important for optimizing the efficiency of other algorithms (such as search and merge algorithms) that require input data to be in sorted lists. Sorting is also often useful for canonicalizing data and for producing humanreadable output.


## ChatGPT's response:

Sorting is a fundamental operation in computer science because it is a common and essential task in many applications. Sorting refers to arranging a collection of elements in a particular order, such as numerical, alphabetical, chronological, or any other custom order. Here are a few reasons why we care about sorting in computer science:

1. Searching: Sorted data is easier to search. Searching an unsorted list requires examining every element in the list, which can be slow and inefficient for large data sets. In contrast, searching a sorted list can be done using efficient algorithms like binary search that require examining only a logarithmic number of elements.
2. Data analysis: Sorting can help in data analysis. Sorting data can help identify patterns and relationships that may not be immediately apparent in an unsorted data set. For example, sorting a list of sales data by date can reveal trends and seasonality.
3. Optimization: Sorting is often a pre-requisite for other algorithms, such as finding the median or calculating the mode of a data set. Sorting can also be used as a building block for other algoritnms, sucn as merge sort or quick sort, which are used in many computational applications.
4. Database management: Sorting is important in managing databases. Databases often contain large amounts of data that need to be sorted in order to perform operations such
as searching, filtering, and data analysis.

In summary, sorting is an essential operation in computer science that helps in searching, data analysis, optimization, and database management. It is a basic building block for many algorithms and is crucial for efficient and effective computation.

## Sorting

How might we sort the list of numbers below.
Can we come up with an algorithm?

| 0 | 1 | 2 | 3 | 4 | 5 |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 10 | 4 | 3 | 0 | 11 | 8 |

## Bubble Sort

Compare two adjacent items, and swap if needed

Repeat until largest item is at the back

Repeat process until done

## Bubble Sort

| 0 | 1 | 2 | 3 | 4 | 5 |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 10 | 4 | 3 | 0 | 11 | 8 |

What do we do first?

## Bubble Sort

$$
\text { len }=6
$$



Compare $\mathrm{j}-1$ and j ; Swap if $\mathrm{L}[\mathrm{j}-1]>\mathrm{L}[\mathrm{j}]$

## Bubble Sort

$$
\text { len }=6
$$



Compare j-1 and j; Swap if $\mathrm{L}[\mathrm{j}-1]>\mathrm{L}[\mathrm{j}]$

## Bubble Sort

$$
\text { len }=6
$$



Compare $\mathrm{j}-1$ and j ; Swap if $\mathrm{L}[\mathrm{j}-1]>\mathrm{L}[\mathrm{j}]$

## Bubble Sort

$$
\text { len }=6
$$



Compare $\mathrm{j}-1$ and j ; Swap if $\mathrm{L}[\mathrm{j}-1]>\mathrm{L}[\mathrm{j}]$

## Bubble Sort

$$
\text { len }=6
$$



Compare j-1 and j; Swap if $\mathrm{L}[\mathrm{j}-1]>\mathrm{L}[\mathrm{j}]$

## Bubble Sort

$$
\text { len }=6
$$



Compare j-1 and j; Swap if L[j-1] > L[j]

## Bubble Sort

$$
\text { len }=6
$$



Compare j-1 and j; Swap if L[j-1] > L[j]

## Bubble Sort

$$
\text { len }=6
$$



Compare j-1 and j; Swap if $\mathrm{L}[\mathrm{j}-1]>\mathrm{L}[\mathrm{j}]$

## Bubble Sort

$$
\text { len }=6
$$



Compare j-1 and j; Swap if $\mathrm{L}[\mathrm{j}-1]>\mathrm{L}[\mathrm{j}]$

## Bubble Sort

$$
\text { len = } 5
$$



Reset and compare pairs with shorter list!

## Bubble Sort

len $=5$


Compare $j-1$ and $j$; Swap if $L[j-1]>L[j]$

## Bubble Sort

len $=5$


Compare $\mathrm{j}-1$ and j ; Swap if $\mathrm{L}[\mathrm{j}-1]>\mathrm{L}[\mathrm{j}]$

## Bubble Sort

len $=5$


Compare j-1 and j; Swap if $\mathrm{L}[\mathrm{j}-1]>\mathrm{L}[\mathrm{j}]$

## Bubble Sort

len $=5$


Compare j-1 and j; Swap if L[j-1] > L[j]

## Bubble Sort

len $=5$


Compare $\mathrm{j}-1$ and j ; Swap if $\mathrm{L}[\mathrm{j}-1]>\mathrm{L}[\mathrm{j}]$

## Bubble Sort

len $=5$


Compare j-1 and j; Swap if L[j-1] > L[j]

## Bubble Sort

len $=4$


Reset and check pairs with shorter list

## Bubble Sort

$$
\text { len }=4
$$



Compare $j-1$ and $j$; Swap if $L[j-1]>L[j]$

## Bubble Sort

$$
\text { len }=4
$$



Compare $j-1$ and $j$; Swap if $L[j-1]>L[j]$

## Bubble Sort

len $=4$


Compare $j-1$ and $j$; Swap if $L[j-1]>L[j]$

## Bubble Sort

len $=3$


Reset; Compare j-1 and j; Swap if $L[j-1]>L[j]$
What next?

## Bubble Sort

len $=3$


Reset; Compare j-1 and j; Swap if $L[j-1]>L[j]$
What next?

## Bubble Sort

len $=2$


Reset; Compare j-1 and j; Swap if $L[j-1]>L[j]$
What next?

## Bubble Sort

Idea: bubble highest values to the end of the list; Check a shorter sublist each time
bubbleSort(L):
for len in range(len(L), 1, -1):
for j in range(1, len): \# bubble if $L[j-1]>L[j]:$
$\operatorname{swap}(j-1, j, L)$

## Bubble sort

 $\operatorname{swap}(i, j, L):$temp $=\mathrm{L}[\mathrm{i}] \#$ step 1
$\mathrm{L}[\mathrm{i}]=\mathrm{L}[\mathrm{j}] \quad$ \# step 2
$\mathrm{L}[\mathrm{j}]=$ temp \# step 3


## Selection Sort

## Selection sort

Repeatedly find the smallest item and put it at front of list
selectionSort(L):
for startldx in range(len(L)):
$\min I d x=$ findMinimum(startldx, L)
swap(startldx, minldx, L)

## Selection Sort

| 0 | 1 | 2 | 3 | 4 | 5 |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 10 | 4 | 3 | 0 | 11 | 8 |

What do we do first?

## Selection Sort



Find minimum element idx between start to end

## Selection Sort



Swap the elements at start and minldx

## Selection Sort



Decrease the interval.

## Selection Sort



Find minimum element between start to end

## Selection Sort



Swap the elements at start and minldx

## Selection Sort



Decrease the interval.

## Selection Sort



Find minimum element idx between start to end

## Selection Sort



Swap the elements at start and minldx
What next?

## Selection Sort



Decrease the interval.

## Selection Sort

| 0 | 1 | 2 | 3 | 4 | 5 |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 3 | 4 | 10 | 11 | 8 |

Find minimum element idx between start to end

## Selection Sort

| 0 | 1 | 2 | 3 | 4 | 5 |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 3 | 4 | 8 | 11 | 10 |

Swap the elements at start and minldx
What next?

## Selection Sort

| 0 | 1 | 2 | 3 | 4 | 5 |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 3 | 4 | 8 | 11 | 10 |

Decrease the interval.
What next?

## Selection Sort



Find minimum element idx between start to end

## Selection Sort



Swap the elements at start and minldx
What next?

## Selection Sort

| 0 | 1 | 2 | 3 | 4 | 5 |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 3 | 4 | 8 | 10 | 11 |

Decrease the interval.
We're done!

## Selection sort

findMinimum(startldx, L):
$\min I d x=s t a r t l d x$
for i in range(startldx, len(L)):
if $\mathrm{L}[\mathrm{i}]<\mathrm{L}[$ minidx $]$ :
$\operatorname{minld}=\mathrm{i}$
return minldx

## Swap <br> $\operatorname{swap}(i, j, L):$ <br> temp $=\mathrm{L}[\mathrm{i}]$ \# step 1 <br> $\mathrm{L}[\mathrm{i}]=\mathrm{L}[\mathrm{j}] \quad$ \# step 2 <br> $\mathrm{L}[\mathrm{j}]=$ temp \# step 3



## Selection sort and Bubble sort are $\mathrm{O}\left(\mathrm{N}^{2}\right)$



