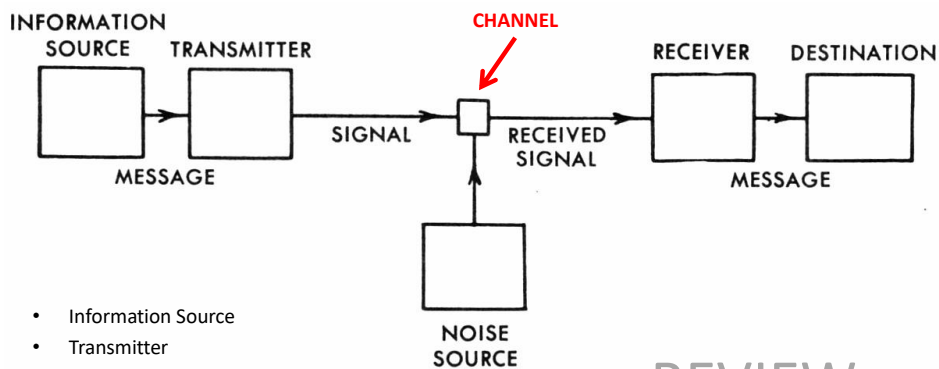


Introduction to Information Theory

Part 3

A General Communication System



- Information Source
- Transmitter
- Channel
- Receiver
- Destination

REVIEW...

Shannon's First Theorem

- By coding sequences of independent symbols (in S^n), it is possible to construct codes such that

$$\lim_{n \rightarrow \infty} \frac{L_n}{n} = H$$

The price paid for such improvement is increased coding complexity (due to increased n) and increased delay in coding.

REVIEW...

Entropy & Coding: Central Ideas

- Use short codes for highly likely events. This shortens the average length of coded messages.
- Code several events at a time. Provides greater flexibility in code design.
- Shannon's Source Coding Theorem
- Algorithms
- Applications

REVIEW...

Source Coding

- **Efficient** codes
- **Singular** codes
- **Nonsingular** codes
- **Instantaneous** codes
- **Universal Codes**

Codes that do not require knowledge of probability distribution but act in the limit as if they did have the knowledge.

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REVIEW...

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Huffman Codes

- Nonsingular
- Instantaneous
- Efficient
- Non-unique
- Powers of a source lead closer to H
- Requires knowledge of symbol probabilities

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REVIEW...

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Entropy & Coding: Central Ideas

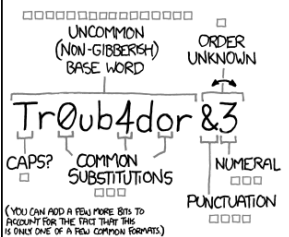

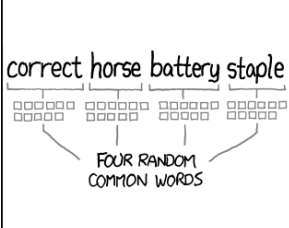
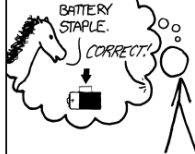
- Use short codes for highly likely events. This shortens the average length of coded messages.
- Code several events at a time. Provides greater flexibility in code design.
- Shannon's Source Coding Theorem
- Algorithms: **Huffman Encoding**, ...
- Applications: **Compression...**

REVIEW...

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xkcd(#936): Password Strength

 <p>UNCOMMON (NON-GIBBERISH) BASE WORD</p> <p>ORDER UNKNOWN</p> <p>Tr0ub4dor & 3</p> <p>CAPS? COMMON SUBSTITUTIONS NUMERAL PUNCTUATION</p> <p>(YOU CAN ADD A FEW MORE BITS TO ACCOUNT FOR THE FACT THAT THIS IS ONLY ONE OF A FEW COMMON REPAIRS)</p>	<p>~28 BITS OF ENTROPY</p> <p>$2^{28} = 3 \text{ DAYS AT } 1000 \text{ GUESSES/SEC}$</p> <p>(PLAUSIBLE ATTACK ON A WEAK REMOTE WEB SERVICE: YES, CRACKING A STOLEN PASSWORD IS FASTER, BUT IT'S NOT WHAT THE PASSWORD WAS SPECIALLY ABOUT.)</p> <p>DIFFICULTY TO GUESS: EASY</p>	<p>WAS IT TROMBONE? NO, TROUBADOR. AND ONE OF THE 0s WAS A ZERO?</p> <p>AND THERE WAS SOME SYMBOL...</p>  <p>DIFFICULTY TO REMEMBER: HARD</p>
 <p>correct horse battery staple</p> <p>FOUR RANDOM COMMON WORDS</p>	<p>~44 BITS OF ENTROPY</p> <p>$2^{44} = 580 \text{ YEARS AT } 1000 \text{ GUESSES/SEC}$</p> <p>DIFFICULTY TO GUESS: HARD</p>	<p>THAT'S A BATTERY STAPLE.</p> <p>CORRECT!</p>  <p>DIFFICULTY TO REMEMBER: YOU'VE ALREADY MEMORIZED IT</p>

THROUGH 20 YEARS OF EFFORT, WE'VE SUCCESSFULLY TRAINED EVERYONE TO USE PASSWORDS THAT ARE HARD FOR HUMANS TO REMEMBER, BUT EASY FOR COMPUTERS TO GUESS.

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Some Observations

- Successive symbols from a source are not always independent.
E.g. in english,

h is more likely to occur following a *t*.
- This **intersymbol dependency** must be accounted for in an accurate measure of entropy.

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Lossless Compression: English Text

- How much can we compress a given text?
- What is the accurate measure of entropy of English texts?
- Zeroth-Order entropy: S

$$H_0 \leq \log\left(\frac{1}{27}\right) \\ \leq 4.755 \text{ bits/letter}$$

- First-Order Entropy: S^2
- Second-Order Entropy: S^3

$$H_1 = 3.3$$

$$H_2 = 3.1$$

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Zipf's Law, 1949

$$P_n \sim 1/n^a$$

P_n is the frequency of occurrence of the n^{th} ranked item and a is close to 1.

- The most frequent word will occur approximately twice as often as the second most frequent word, three times as often as the third most frequent word, etc.
- For example, in a corpus of over 1 million words:

the	69,971	7%
of	36,411	3.5%
and	28,852	2.9%

- For English text:

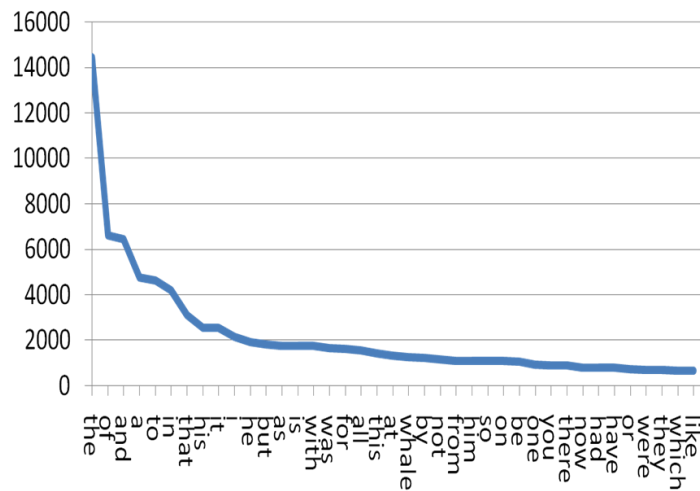
$$p_m = \frac{A}{m}$$

where m is the word's rank, p_m is the probability of the m^{th} rank word, A is a constant that depends on the number of active words in the language.

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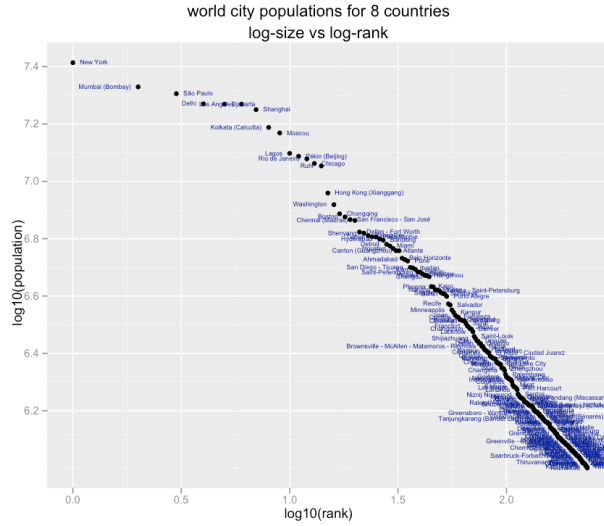
The "Long Tail" of Moby Dick



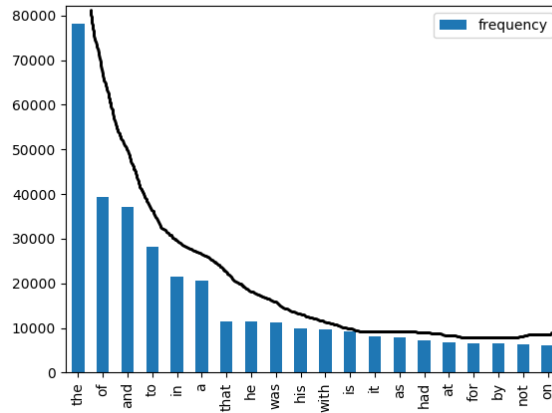
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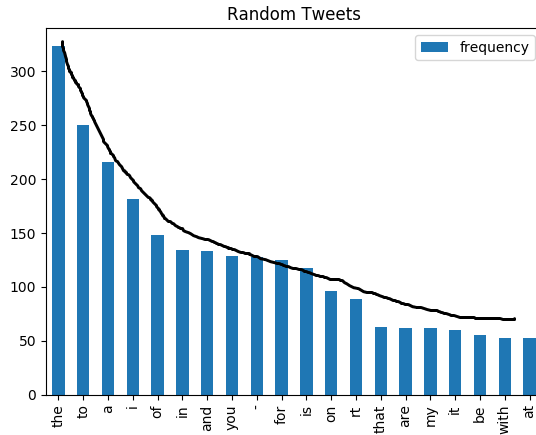
Zipf's Law and Populations



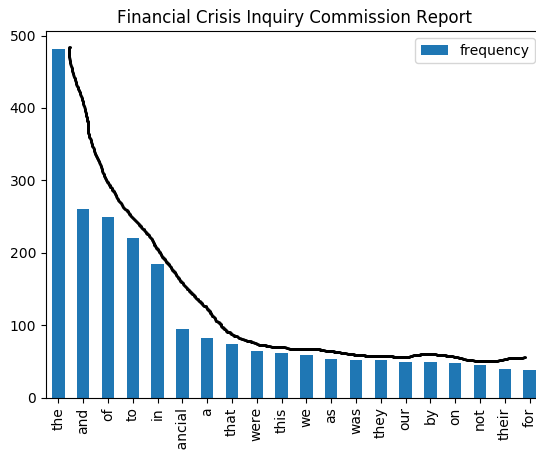
The Adventures of Sherlock Holmes – A. C. Doyle



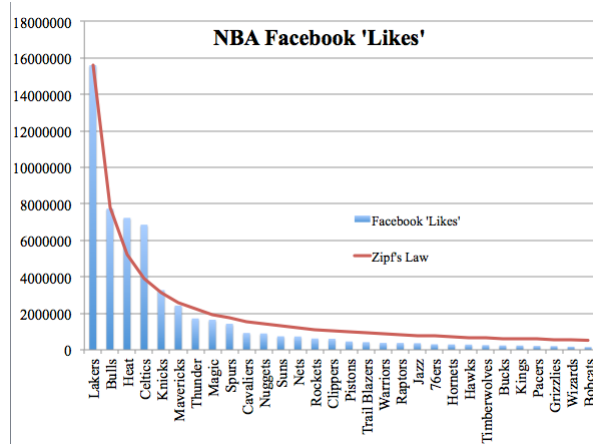
Random Tweets



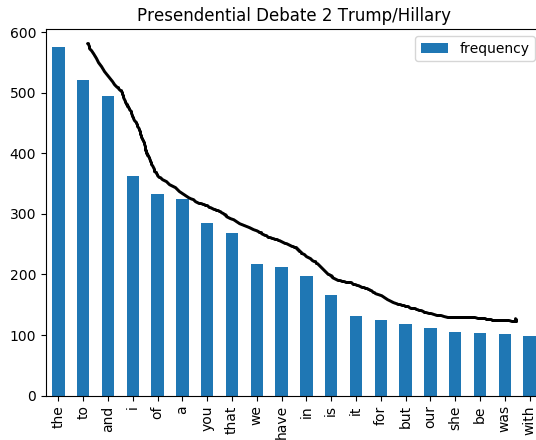
Financial Crisis Enquiry Commission Report



NBA Facebook Likes



2nd Presidential Debate between Clinton & Trump



Estimating Entropy of English Text

- For English text, W with M words:

$$p_m = \frac{A}{m}$$

where m is the word's rank, p_m is the probability of the m^{th} rank word, A is a constant that depends on the number of active words in the language.

- If $A = 0.1$, so that

$$\sum_{m=1}^M p_m = 1$$

We need, $M = 12,368$.

$$H_W = \sum_{m=1}^{m=12,368} \frac{1}{m} \log\left(\frac{m}{.1}\right) = 9.716391 \text{ bits/word}$$

If $\bar{w}=4.7$ letters/word

$$H = \frac{9.716391}{4.7} = 2.067 \text{ bits}$$

Shannon Redundancy

$$R = 1 - \frac{H}{\log M}$$

Where H is the per letter entropy, M is the size of the source alphabet. Thus redundancy of English is

$$1 - \frac{2.067}{\log 27} = 56.5\%$$

With an entropy of 1.5 we get 67% redundancy.

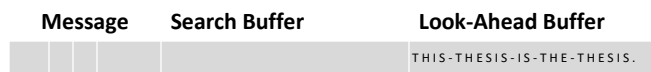
I.e. Huffman coding (even with an entropy of 3.3 or 3.1) will not get close to the theoretical limit.

Can we achieve compression rates close to 33%???

Lempel-Ziv Coding

- Sequences of text repeat patterns (words, phrases, etc)
- Construct a dictionary of common patterns
- Send references to patterns as triples (x, y, z)

Lempel-Ziv Coding (LZ77)



Lempel-Ziv Coding (LZ77)

Message	Search Buffer	Look-Ahead Buffer
		THIS-THE-SIS-IS-THE-THE-SIS.
0 0 T T		THIS-THE-SIS-IS-THE-THE-SIS.

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Lempel-Ziv Coding (LZ77)

Message	Search Buffer	Look-Ahead Buffer
		THIS-THE-SIS-IS-THE-THE-SIS.
0 0 T T		THIS-THE-SIS-IS-THE-THE-SIS.
0 0 H H		THIS-THE-SIS-IS-THE-THE-SIS.

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Lempel-Ziv Coding (LZ77)

Message				Search Buffer	Look-Ahead Buffer
					THIS-THE-SIS-IS-THE-THE-SIS.
0	0	T	T		THIS-THE-SIS-IS-THE-THE-SIS.
0	0	H	H		THIS-THE-SIS-IS-THE-THE-SIS.
0	0	I	I		THIS-THE-SIS-IS-THE-THE-SIS.

Lempel-Ziv Coding (LZ77)

Message				Search Buffer	Look-Ahead Buffer
					THIS-THE-SIS-IS-THE-THE-SIS.
0	0	T	T		THIS-THE-SIS-IS-THE-THE-SIS.
0	0	H	H		THIS-THE-SIS-IS-THE-THE-SIS.
0	0	I	I		THIS-THE-SIS-IS-THE-THE-SIS.
0	0	S	S		THIS-THE-SIS-IS-THE-THE-SIS.

Lempel-Ziv Coding (LZ77)

Message				Search Buffer	Look-Ahead Buffer
					THIS-THE-SIS-IS-THE-THE-SIS.
0	0	T	T		THIS-THE-SIS-IS-THE-THE-SIS.
0	0	H	H		THIS-THE-SIS-IS-THE-THE-SIS.
0	0	I	I		THIS-THE-SIS-IS-THE-THE-SIS.
0	0	S	S		THIS-THE-SIS-IS-THE-THE-SIS.
0	0	-	-		THIS-THE-SIS-IS-THE-THE-SIS.

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Lempel-Ziv Coding (LZ77)

Message				Search Buffer	Look-Ahead Buffer
					THIS-THE-SIS-IS-THE-THE-SIS.
0	0	T	T		THIS-THE-SIS-IS-THE-THE-SIS.
0	0	H	H		THIS-THE-SIS-IS-THE-THE-SIS.
0	0	I	I		THIS-THE-SIS-IS-THE-THE-SIS.
0	0	S	S		THIS-THE-SIS-IS-THE-THE-SIS.
0	0	-	-		THIS-THE-SIS-IS-THE-THE-SIS.
5	2	E	E		THIS-THE-SIS-IS-THE-THE-SIS.

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Lempel-Ziv Coding (LZ77)

Message				Search Buffer	Look-Ahead Buffer
					THIS-THE-SIS-IS-THE-THE-SIS.
0	0	T	T		THIS-THE-SIS-IS-THE-THE-SIS.
0	0	H	H		THIS-THE-SIS-IS-THE-THE-SIS.
0	0	I	I		THIS-THE-SIS-IS-THE-THE-SIS.
0	0	S	S		THIS-THE-SIS-IS-THE-THE-SIS.
0	0	-	-		THIS-THE-SIS-IS-THE-THE-SIS.
5	2	E	THE		THIS-THE-SIS-IS-THE-THE-SIS.
5	1	I	SI		THIS-THE-SIS-IS-THE-THE-SIS.

Lempel-Ziv Coding (LZ77)

Message				Search Buffer	Look-Ahead Buffer
					THIS-THE-SIS-IS-THE-THE-SIS.
0	0	T	T		THIS-THE-SIS-IS-THE-THE-SIS.
0	0	H	H		THIS-THE-SIS-IS-THE-THE-SIS.
0	0	I	I		THIS-THE-SIS-IS-THE-THE-SIS.
0	0	S	S		THIS-THE-SIS-IS-THE-THE-SIS.
0	0	-	-		THIS-THE-SIS-IS-THE-THE-SIS.
5	2	E	THE		THIS-THE-SIS-IS-THE-THE-SIS.
5	1	I	SI		THIS-THE-SIS-IS-THE-THE-SIS.
7	2	I	SI		THIS-THE-SIS-IS-THE-THE-SIS.

Lempel-Ziv Coding (LZ77)

	Message	Search Buffer	Look-Ahead Buffer
			THIS-THE-SIS-IS-THE-THE-SIS.
0	0 T T		THIS-THE-SIS-IS-THE-THE-SIS.
0	0 H H		THIS-THE-SIS-IS-THE-THE-SIS.
0	0 I I		THIS-THE-SIS-IS-THE-THE-SIS.
0	0 S S		THIS-THE-SIS-IS-THE-THE-SIS.
0	0 - -		THIS-THE-SIS-IS-THE-THE-SIS.
5	2 E THE		THIS-THE-SIS-IS-THE-THE-SIS.
5	1 I SI		THIS-THE-SIS-IS-THE-THE-SIS.
7	2 I S-I		THIS-THE-SIS-IS-THE-THE-SIS.
10	5 - S-THE-		THIS-THE-SIS-IS-THE-THE-SIS.

Lempel-Ziv Coding (LZ77)

	Message	Search Buffer	Look-Ahead Buffer
			THIS-THE-SIS-IS-THE-THE-SIS.
0	0 T T		THIS-THE-SIS-IS-THE-THE-SIS.
0	0 H H		THIS-THE-SIS-IS-THE-THE-SIS.
0	0 I I		THIS-THE-SIS-IS-THE-THE-SIS.
0	0 S S		THIS-THE-SIS-IS-THE-THE-SIS.
0	0 - -		THIS-THE-SIS-IS-THE-THE-SIS.
5	2 E THE		THIS-THE-SIS-IS-THE-THE-SIS.
5	1 I SI		THIS-THE-SIS-IS-THE-THE-SIS.
7	2 I S-I		THIS-THE-SIS-IS-THE-THE-SIS.
10	5 - S-THE-		THIS-THE-SIS-IS-THE-THE-SIS.
14	6 . THE-SIS.	THIS-THE-SIS-IS-THE-THE-SIS.	

Lempel-Ziv Coding

- Sequences of text repeat patterns (words, phrases, etc)
- Construct a dictionary of common patterns
- Send references to patterns as triples (x, y, z)
e.g. $(5, 3, F)$
go back 5 received chars
take the next 3 from there
add F to the end
- Size of Search Buffer and Look-Ahead Buffer is finite.
- Used by ZIP, PKZip, Lharc, PNG, gzip, ARJ
- Extended to LZ78 (uses dictionary), LZW (+Terry Welch)
- Achieves optimal rate of transmission in the long run w/o using probability dist.

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Decode

Message

0	0	I	
0	0	-	
0	0	M	
3	1	S	
1	1	-	
5	5	L	
5	3	Y	

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Decode

Message

0	0	I	
0	0	-	
0	0	M	
3	1	S	
1	1	-	
5	5	L	
5	3	Y	

Decode

Message

0	0	I	
0	0	-	-
0	0	M	
3	1	S	
1	1	-	
5	5	L	
5	3	Y	

Decode

Message

0	0	I	I
0	0	-	I -
0	0	M	I - M
3	1	S	
1	1	-	
5	5	L	
5	3	Y	

Decode

Message

0	0	I	I
0	0	-	I -
0	0	M	I - M
3	1	S	I - M I S
1	1	-	
5	5	L	
5	3	Y	

Decode

Message

0	0	I	I
0	0	-	I-
0	0	M	I-M
3	1	S	I-MIS
1	1	-	I-MISS-
5	5	L	
5	3	Y	

Decode

Message

0	0	I	I
0	0	-	I-
0	0	M	I-M
3	1	S	I-MIS
1	1	-	I-MISS-
5	5	L	I-MISS-MISS-L
5	3	Y	

Decode

Message

0	0	I	I
0	0	-	I -
0	0	M	I - M
3	1	S	I - M I S
1	1	-	I - M I S S -
5	5	L	I - M I S S - M I S S - L
5	3	Y	I - M I S S - M I S S - L I S S Y

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