

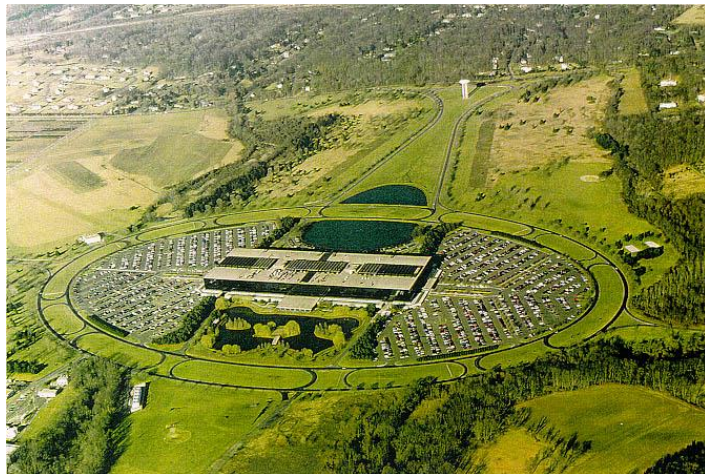
# Introduction to Information Theory Part 1

Deepak Kumar  
Bryn Mawr College

1

## 1948, Bell Labs

- Two very significant events
- Two neologisms



2

## 1948, Bell Labs

- Two very significant events
  1. Invention of an amazingly simple device that could do anything a vacuum tube could do and more efficiently.

*Semiconductor triode? lotatron? ...*

- Two neologisms
  1. *Transistor*

Spawned the electronic revolution!



3

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# Shockley Semiconductor Lab



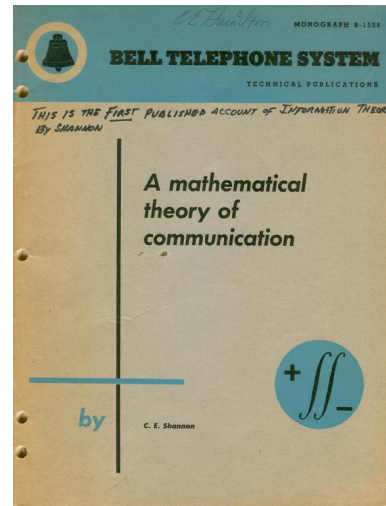
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# Shockley Semiconductor Lab

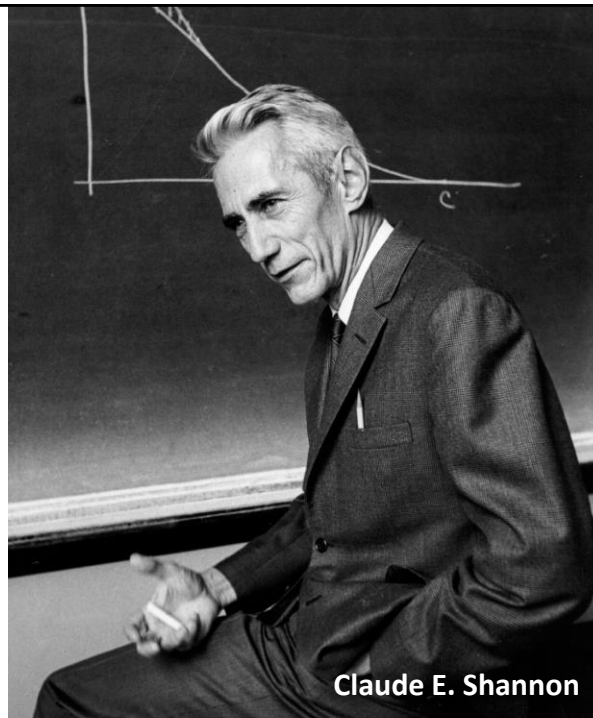


## 1948, Bell Labs

- Two very significant events
  1. Invention of transistor
  2. Publication of a monograph
  
- Two neologisms
  1. *Transistor*
  2. *bit* – unit of measuring information  
Not the same as binary digit.

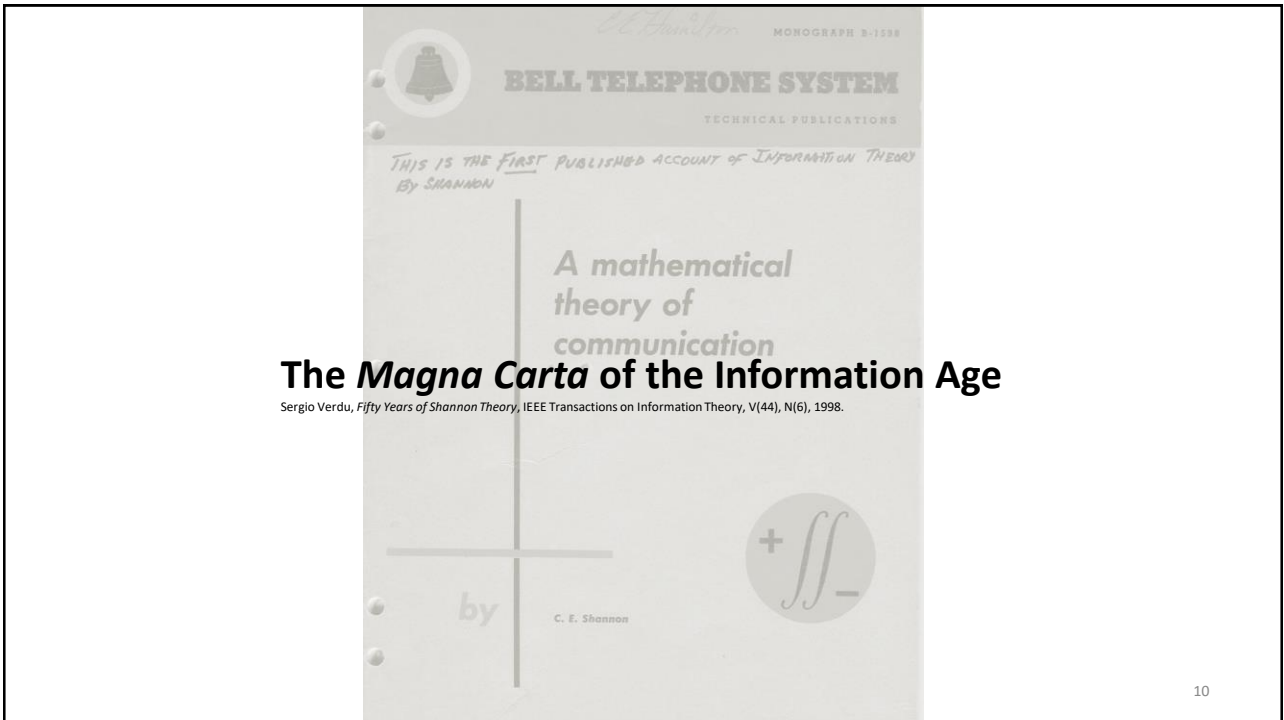
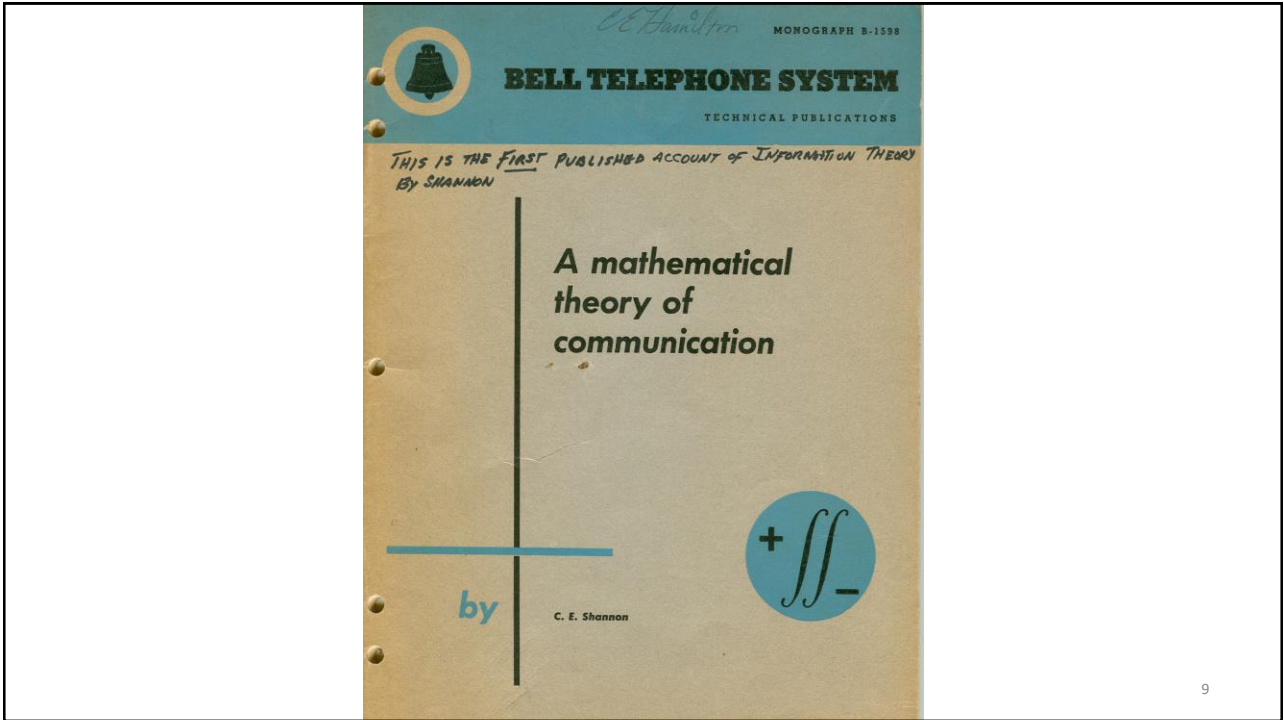


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Claude E. Shannon

8



## The Magna Carta of the Information Age

Sergio Verdu, *Fifty Years of Shannon Theory*, IEEE Transactions on Information Theory, V(44), N(6), 1998.

## A General Communication System

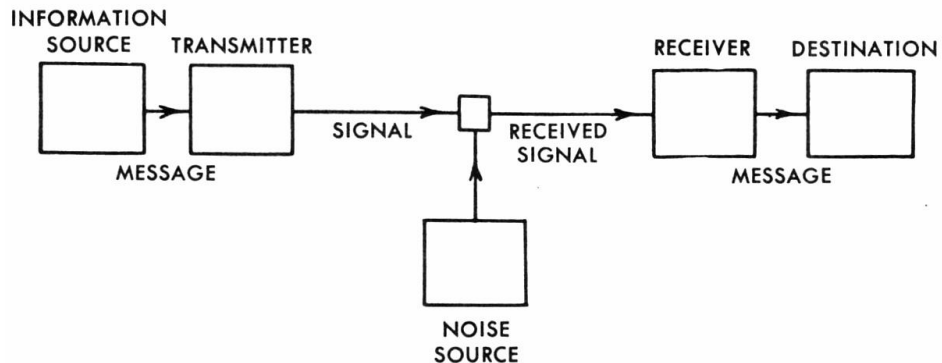
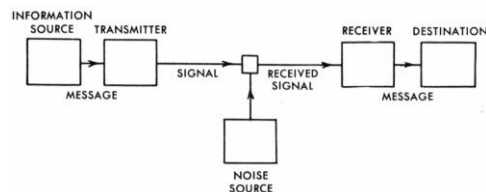


Figure 1.

*A Mathematical Theory of Communication, by Claude E. Shannon, The Bell System Technical Journal, Vol. 27, pp. 379–423, 623–656, July, October, 1948.*

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## Shannon's Information Theory



- Conceptualization of information & modeling of information sources
- Sending of information across the channel:
  - What are the limits on the amount of information that can be sent?
  - What is the effect of noise on this communication.

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# A General Communication System

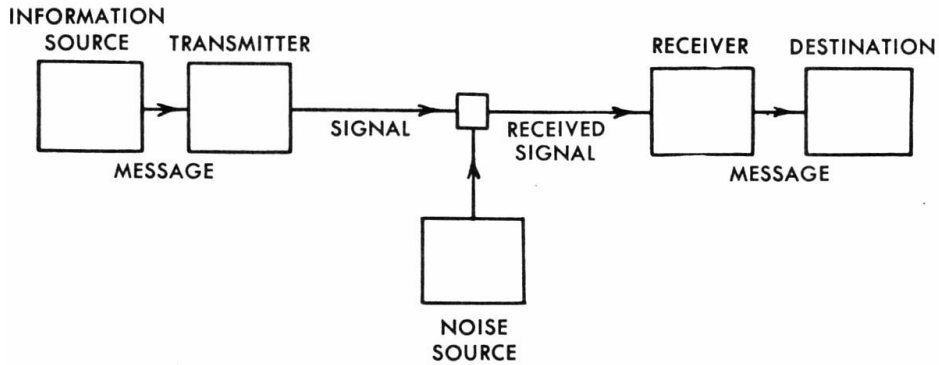
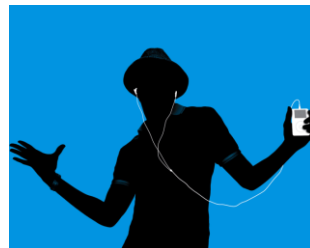
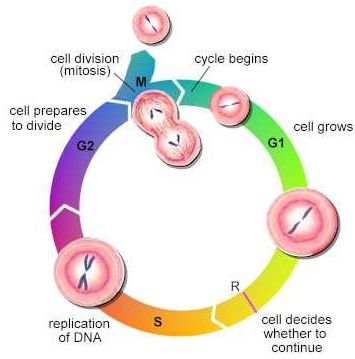
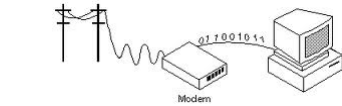


Figure 1.

*A Mathematical Theory of Communication*, by Claude E. Shannon, *The Bell System Technical Journal*, Vol. 27, pp. 379-423, 623-656, July, October, 1948.

# Examples of Communication Systems



# How much information?

ᚫᚱᚴᚢᚦᚢᚱ ᚩᚱ ᚲᚢᚱᚢᚱᚢᚱᚢᚱᚢᚱ

ᚠ ᚡ ᚢ ᚣ ᚤ ᚥ ᚦ ᚧ ᚨ ᚩ ᚪ ᚫ ᚬ  
a b c d e f g h i j k l m

ᚸ ᚹ ᚺ ᚻ ᚼ ᚽ ᚾ ᚿ ᛀ ᛁ ᛂ ᛃ  
n o p q r s t u v w x y z

9/6/2012

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# Some more Trivia: Origins of Bluetooth

ᚠ ᚡ ᚢ ᚣ ᚤ ᚥ ᚦ ᚧ ᚨ ᚩ ᚪ ᚫ ᚬ  
a b c d e f g h i j k l m

ᚸ ᚹ ᚺ ᚻ ᚼ ᚽ ᚾ ᚿ ᛀ ᛁ ᛂ ᛃ  
n o p q r s t u v w x y z

$$H (\text{ᚨ}) + B (\text{ᚢ}) = \text{Bluetooth logo}$$

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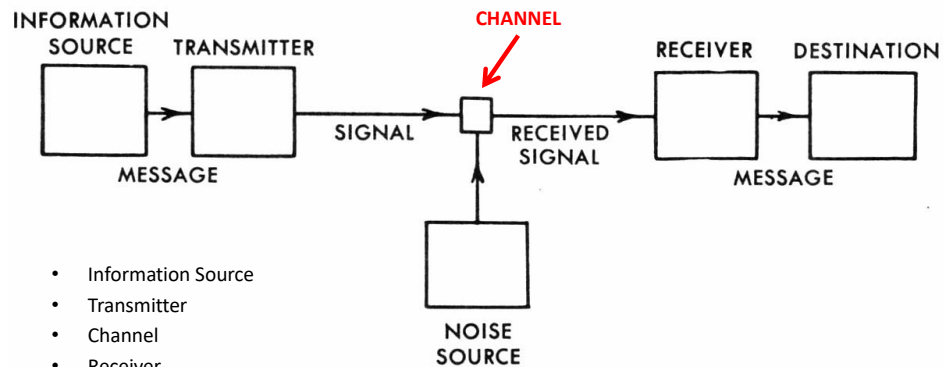


“The fundamental problem of communication is that of reproducing at one point either exactly or approximately a message selected at another point.”

*A Mathematical Theory of Communication*, by Claude E. Shannon, *The Bell System Technical Journal*, Vol. 27, pp. 379–423, 623–656, July, October, 1948.

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## A General Communication System



- Information Source
- Transmitter
- Channel
- Receiver
- Destination

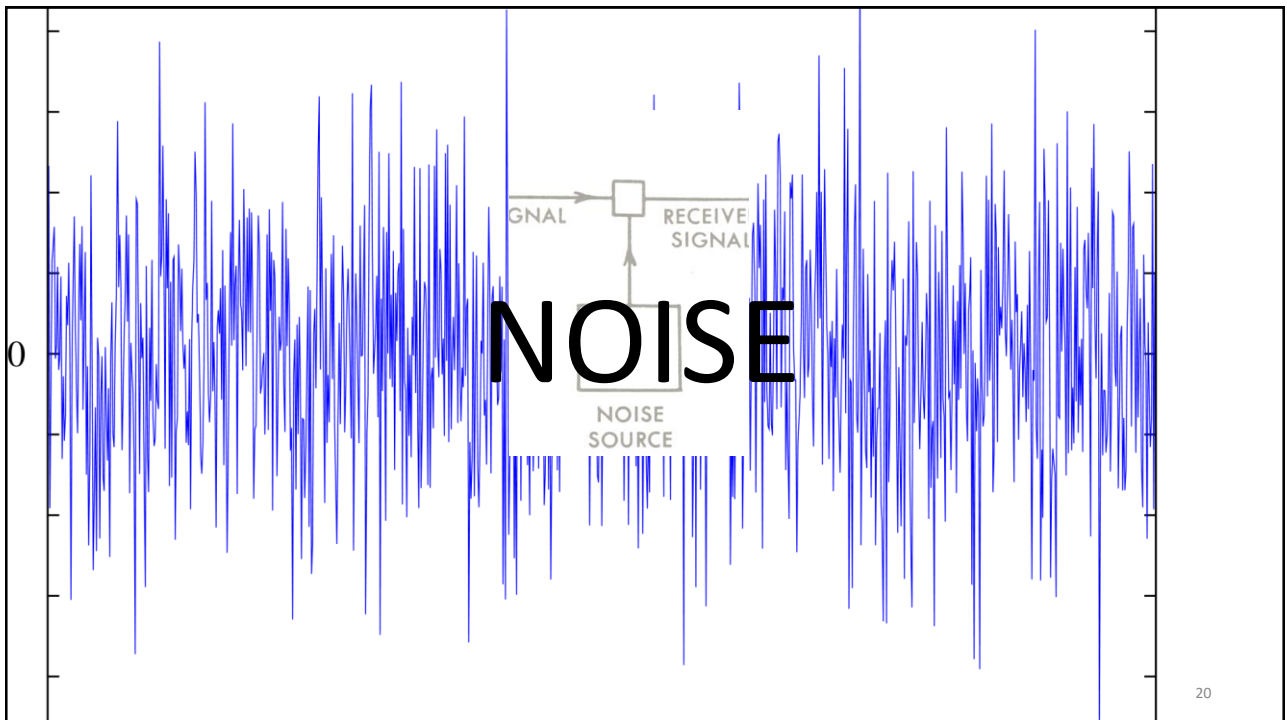
*A Mathematical Theory of Communication*, by Claude E. Shannon, *The Bell System Technical Journal*, Vol. 27, pp. 379–423, 623–656, July, October, 1948.

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## Perfect Communication (*Noiseless Channel*)

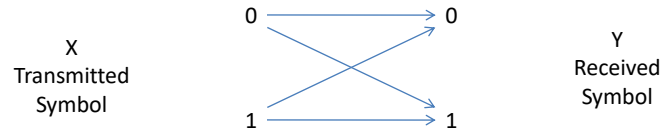


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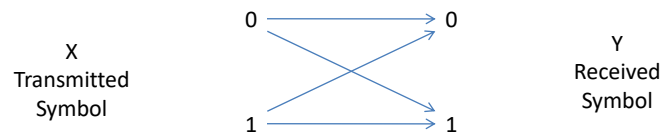
20

## Motivating Noise...



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## Motivating Noise...

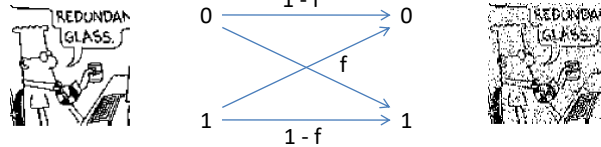


$$\begin{array}{ll}
 P(Y=1|X=0) = f & P(Y=0|X=0) = 1 - f \\
 P(Y=0|X=1) = f & P(Y=1|X=1) = 1 - f
 \end{array}$$

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## Motivating Noise...

$$f = 0.1, n = \sim 10,000$$



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## Question

How can we achieve perfect communication over an imperfect, noisy communication channel?

- Use more reliable components;
- Stabilize the environment;
- Use larger areas;
- Use power/cooling to reduce thermal noise.

These are all costly solutions.

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## Alternately...

How can we achieve perfect communication over an imperfect, noisy communication channel?

- Accept that there will be noise
- Add error detection and correction
- Introduce the concepts of ENCODER/DECODER

### **Information Theory**

Theoretical limitations of such systems

### **Coding Theory**

Creation of practical encoding/decoding systems

25

## Alternately...

How can we achieve perfect communication over an imperfect, noisy communication channel?

- Accept that there will be noise
- Add error detection and correction
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**REDUNDANCY IS KEY!**

### **Information Theory**

Theoretical limitations of such systems

### **Coding Theory**

Creation of practical encoding/decoding systems

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## Shannon's Insight

High Reliability → Low Transmission Rate

I.e. Perfect reliability → Zero Transmission Rate

For a given level of noise there is an associated rate of transmission that can be achieved with arbitrarily good reliability.

e.g. Sending a lone T versus THIS...

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## Meaning? What meaning?

“Frequently the messages have *meaning*; that is they refer to or are correlated according to some system with certain physical or conceptual entities.

These semantic aspects of communication are irrelevant to the engineering problem.”

A Mathematical Theory of Communication, by  
Claude E. Shannon, *The Bell System Technical Journal*,  
Vol. 27, pp. 379–423, 623–656, July, October, 1948.

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## What is information?

- Just the physical aspects...Shannon, 1948
- The General Definition of Information...Floridi, 2010.

GDI)  $\sigma$  is an instance of information, understood as semantic content, if and only if:

GDI.1)  $\sigma$  consists of *n data*, for  $n \geq 1$ ;

GDI.2) the data are *well formed*;

GDI.3) the well-formed data are *meaningful*.

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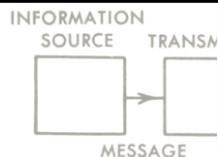
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## Shannon Information is...

- **Uncertainty**  
Can be measured by counting the number of possible messages.
- **Surprise**  
Some messages are more likely than others.
- **Difficult**  
What is significant is the difficulty of transmitting the message from one point to the other.
- **Entropy**  
A fundamental measure of information.

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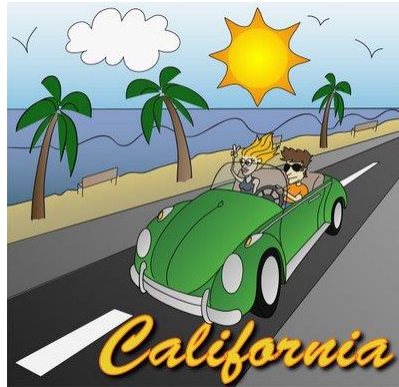


## Information Source

- An information source generates a finite number of messages (or symbols).
- Information is quantified using probabilities.
- Given a finite set of possible messages, associate a probability with each message.
- A message with low probability represents more information than one with high probability.

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## Understanding Information



*It is sunny in California today!*

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## Information: Definition

- Information is quantified using probabilities.
- Given a finite set of possible messages, associate a probability with each message.
- A message with low probability represents more information than one with high probability.

**Definition of Information:**

$$I(p) = \log\left(\frac{1}{p}\right) = -\log(p)$$

Where  $p$  is the probability of the message

Base 2 is used for the logarithm so  $I$  is measured in **bits**

**Trits** for base 3, **nats** for base  $e$ , **Hartleys** for base 10...

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## Information in a coin flip

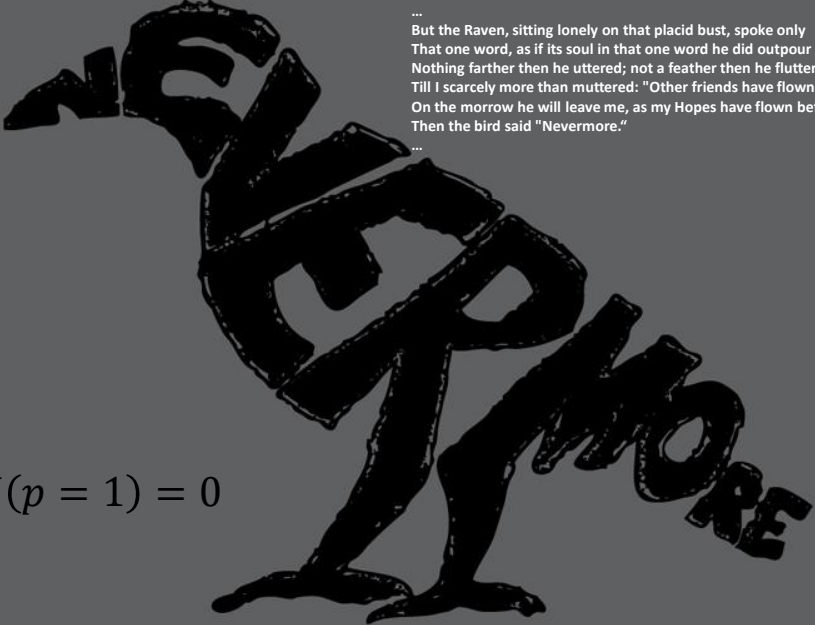
$$P(\text{HEADS}) = \frac{1}{2}$$

$$I = -\log\left(\frac{1}{2}\right) = 1 \text{ bit of information}$$

Given a sequence of 14 coin flips: hthhththhtht

We will need 14 bits: 10110010110010

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The Raven  
Edgar Allan Poe

...

But the Raven, sitting lonely on that placid bust, spoke only  
That one word, as if its soul in that one word he did outpour  
Nothing farther then he uttered; not a feather then he fluttered--  
Till I scarcely more than muttered: "Other friends have flown before--  
On the morrow he will leave me, as my Hopes have flown before."  
Then the bird said "Nevermore."  
...

$I(p = 1) = 0$

Picture From: <http://adrianamelnic.ro/2012/06/theraven/>

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$$I(p) = \log(1/p) = -\log(p)$$

Some properties of  $I$

1.  $I(p) \geq 0$

Information is non-negative.

2.  $I(p_1 * p_2) = I(p_1) + I(p_2)$

Information we get from observing two independent events occurring is the sum of two information(s).

3.  $I(p)$  is monotonic and continuous in  $p$

Slight changes in probability incur slight changes in information.

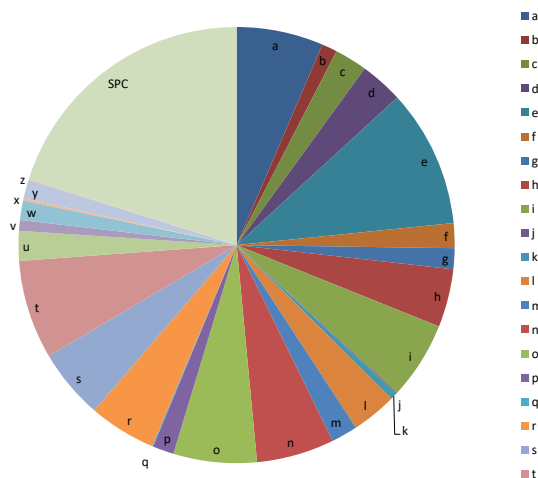
4.  $I(p = 1) = 0$

We get zero information from an event whose probability is 1.

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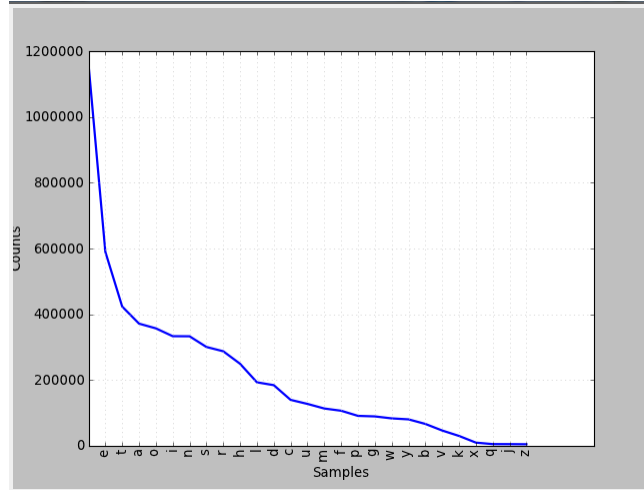
## Example: Text Analysis

a	0.06428
b	0.01147
c	0.02413
d	0.03188
e	0.10210
f	0.01842
g	0.01543
h	0.04313
i	0.05767
j	0.00082
k	0.00514
l	0.03338
m	0.01959
n	0.05761
o	0.06179
p	0.01571
q	0.00084
r	0.04973
s	0.05199
t	0.07327
u	0.02201
v	0.00800
w	0.01439
x	0.00162
y	0.01387
z	0.00077
SPC	0.20098

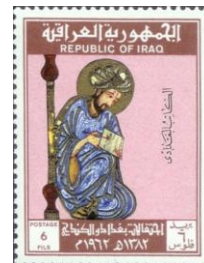


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# Example: Text Analysis



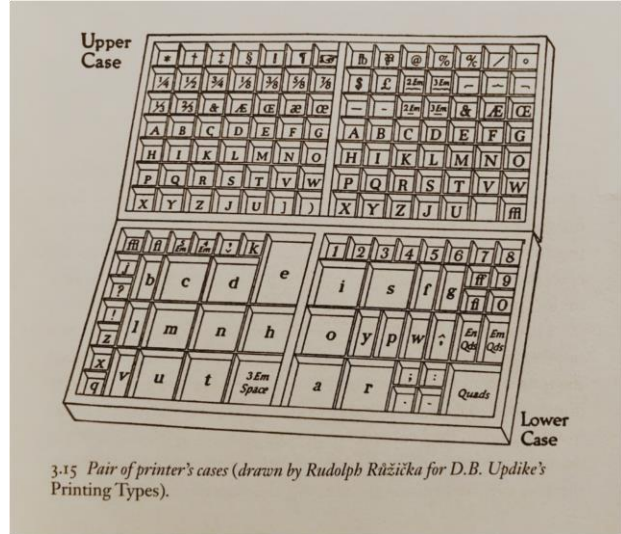
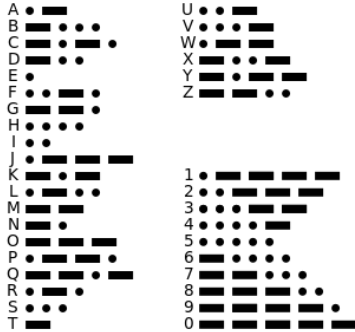
## al Kindi (9<sup>th</sup> Century C.E.)



From: Mathematicians on Stamps: <http://jeff560.tripod.com/stamps.html>

# Morse Code

- Short codes for common letters, longer codes for less common letters
- Sounds: [Click here](#)



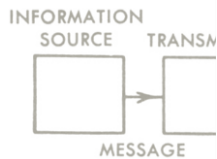
9/11/2019

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# Example: Text Analysis

Letter	Freq.	I
a	0.06428	3.95951
b	0.01147	6.44597
c	0.02413	5.37297
d	0.03188	4.97116
e	0.10210	3.29188
f	0.01842	5.76293
g	0.01543	6.01840
h	0.04313	4.53514
i	0.05767	4.11611
j	0.00082	10.24909
k	0.00514	7.60474
l	0.03338	4.90474
m	0.01959	5.67385
n	0.05761	4.11743
o	0.06179	4.01654
p	0.01571	5.99226
q	0.00084	10.21486
r	0.04973	4.32981
s	0.05199	4.26552
t	0.07327	3.77056
u	0.02201	5.50592
v	0.00800	6.96640
w	0.01439	6.11899
x	0.00162	9.26697
y	0.01387	6.17152
z	0.00077	10.34877
SPC	0.20096	2.31502

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## Information Source

- An information source generates a finite number of messages (or symbols)

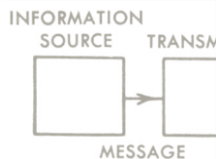
$$\{a_1, a_2, \dots, a_n\}$$

- Source emits the symbols with probabilities

$$P = \{p_1, p_2, \dots, p_n\}$$

- Assume independence: successive symbols do not depend on past symbols.
- What is the average amount of information?

43



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**ANSWER: Entropy!**

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## Definition of Entropy

- Information ( $I$ ) is associated with known events/messages
- Entropy ( $H$ ) is the average information w.r.to all possible outcomes

$$H(P) = \sum_i p_i \log \frac{1}{p_i}$$

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y	0.01387	6.17152
z	0.00077	10.34877
SPC	0.20096	2.31502

$$H(P) = \sum_i p_i \log \frac{1}{p_i} = 4.047$$

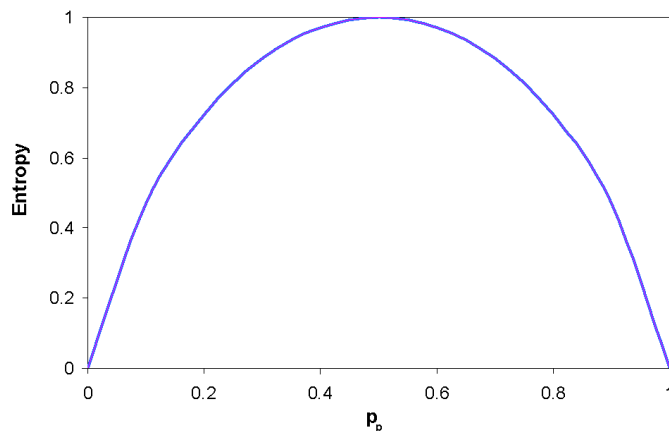
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## Entropy: What about it?

- Does  $H(P)$  have a maximum? Where?
- Is entropy a good name for this stuff? How is it related to entropy in thermodynamics?
- How does entropy help in communication? What else can we do with it?
- Why use the letter  $H$ ? 😊

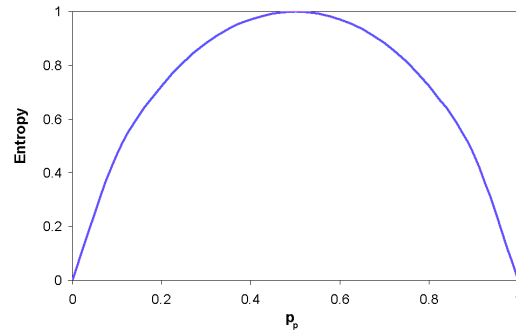
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## Entropy (2 outcomes)



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## Entropy: Properties



$$0 \leq H(P_n) \leq \log(n)$$

Entropy is maximized if  $P$  is uniform.

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## H for Entropy?

"The enthalpy is [often] written  $U$ .  $V$  is the volume, and  $Z$  is the partition function.  $P$  and  $Q$  are the position and momentum of a particle.  $R$  is the gas constant, and of course  $T$  is temperature.  $W$  is the number of ways of configuring our system (the number of states), and we have to keep  $X$  and  $Y$  in case we need more variables. Going back to the first half of the alphabet,  $A$ ,  $F$ , and  $G$  are all different kinds of free energies (the last named for Gibbs).  $B$  is a virial coefficient or a magnetic field.  $I$  will be used as a symbol for information;  $J$  and  $L$  are angular momenta.  $K$  is Kelvin, which is the proper unit of  $T$ .  $M$  is magnetization, and  $N$  is a number, possibly Avogadro's, and  $O$  is too easily confused with  $0$ . This leaves  $S$  . . ."

From Spikes: Exploring the Neural Code,  
by Reike et al, Bradford 1999.

...and  $H$ .

In Spikes they also eliminate  $H$  (e.g., as the Hamiltonian). But, others,  $I$ , along with Shannon, prefer to honor Hartley. Thus,  $H$  for entropy . . .

From: Tom Carter's Lecture Notes on Information Theory and Entropy,  
Prepared for Complex Systems Summer School, Santa Fe, June 2012.  
sustan.csustan.edu/~tom/

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