Introduction to Information Theory

Part 3

Assignment#1 Results

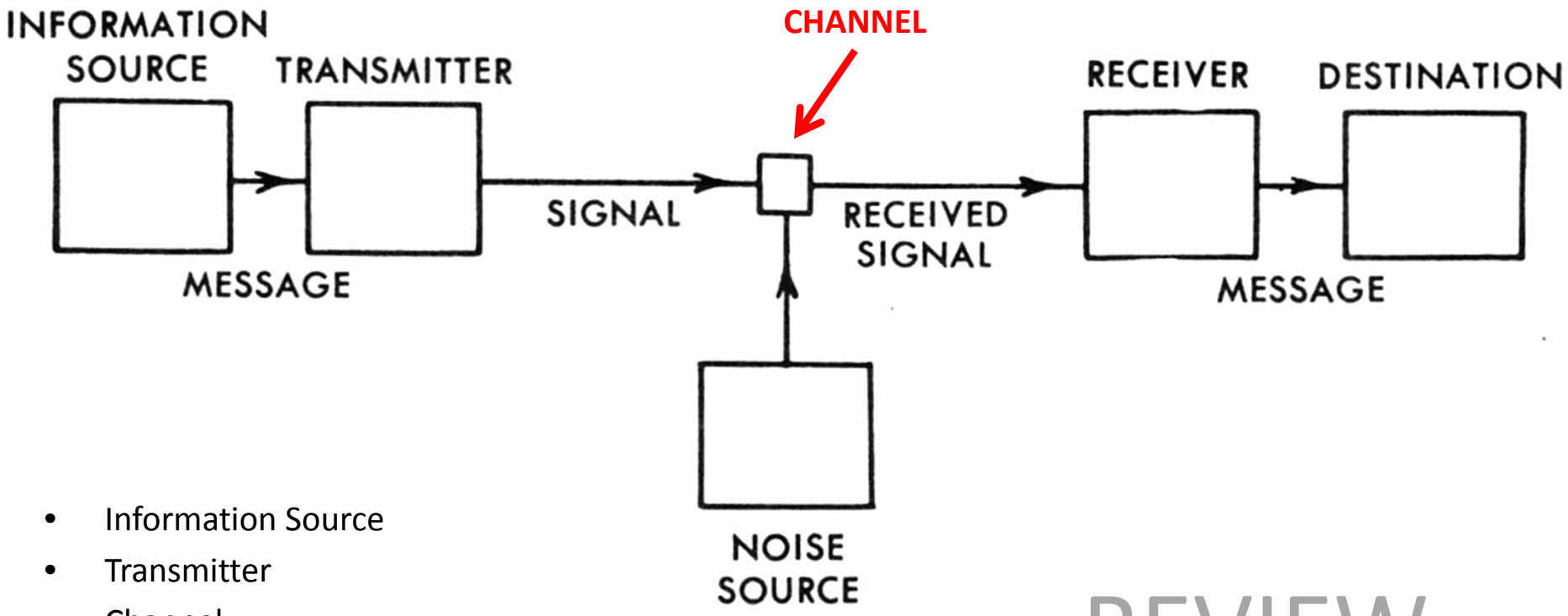
- List text(s) used, total # letters, computed entropy of text. Compare results.

- What is the computed average word length of 3-letter codes on

$$S = \{A, B\}, P = \{0.75, 0.25\}$$

  Compare results.

A General Communication System



- Information Source
- Transmitter
- Channel
- Receiver
- Destination

REVIEW…

Information Source

- Definition of Information:

$$I(p) = log\left(\frac{1}{p}\right) = -log(p)$$

- **Information (I)** is associated with known events/messages
- **Entropy (H)** is the average information w.r.to all possible outcomes.

Given, $P = \{p_1, p_2, \dots, p_3\}$

$$H(P) = \sum_i p_i \log(\frac{1}{p_i})$$

REVIEW...

Source Coding: Basics

- **Block code**: When all codes have the same length. For example, ASCII (8-bits)
- **Average Word Length**:

$$L = \sum_{i=1}^{m} p_i l_i$$

More generally,

$$L_n = \frac{1}{n} \sum_{i=1}^{m} p_i l_i$$

- A code is **efficient** if it has the smallest average word length. (Turns out entropy is the benchmark…)

REVIEW…

Average Code Length & Entropy

- Average length bounds: $H \leq L < H + 1$
- Grouping n symbols together:

$$H(S^n) \leq L < H(S^n) + 1$$

$$nH(S) \leq L < nH(S) + 1$$

$$H(S) \leq \frac{L}{n} < H(S) + \frac{1}{n}$$

$$\lim_{n \to \infty} \frac{L_n}{n} = \mathrm{H}$$

REVIEW…

Shannon's First Theorem

- By coding sequences of independent symbols (in S^n), it is possible to construct codes such that

$$\lim_{n\to\infty} \frac{L_n}{n} = \mathrm{H}$$

The price paid for such improvement is increased coding complexity (due to increased n) and increased delay in coding.

REVIEW...

Entropy & Coding: Central Ideas

- Use short codes for highly likely events. This shortens the average length of coded messages.
- Code several events at a time. Provides greater flexibility in code design.
- Shannon's Source Coding Theorem
- Algorithms
- Applications

REVIEW…

Source Coding

- **Efficient** codes
- **Singular** codes
- **Nonsingular** codes
- **Instantaneous** codes
- **Universal** Codes

  Codes that do not require knowledge of probability distribution but act in the limit as if they did have the knowledge.

REVIEW...

Huffman Codes

- Nonsingular
- Instantaneous
- Efficient
- Non-unique
- Powers of a source lead closer to H
- Requires knowledge of symbol probabilities

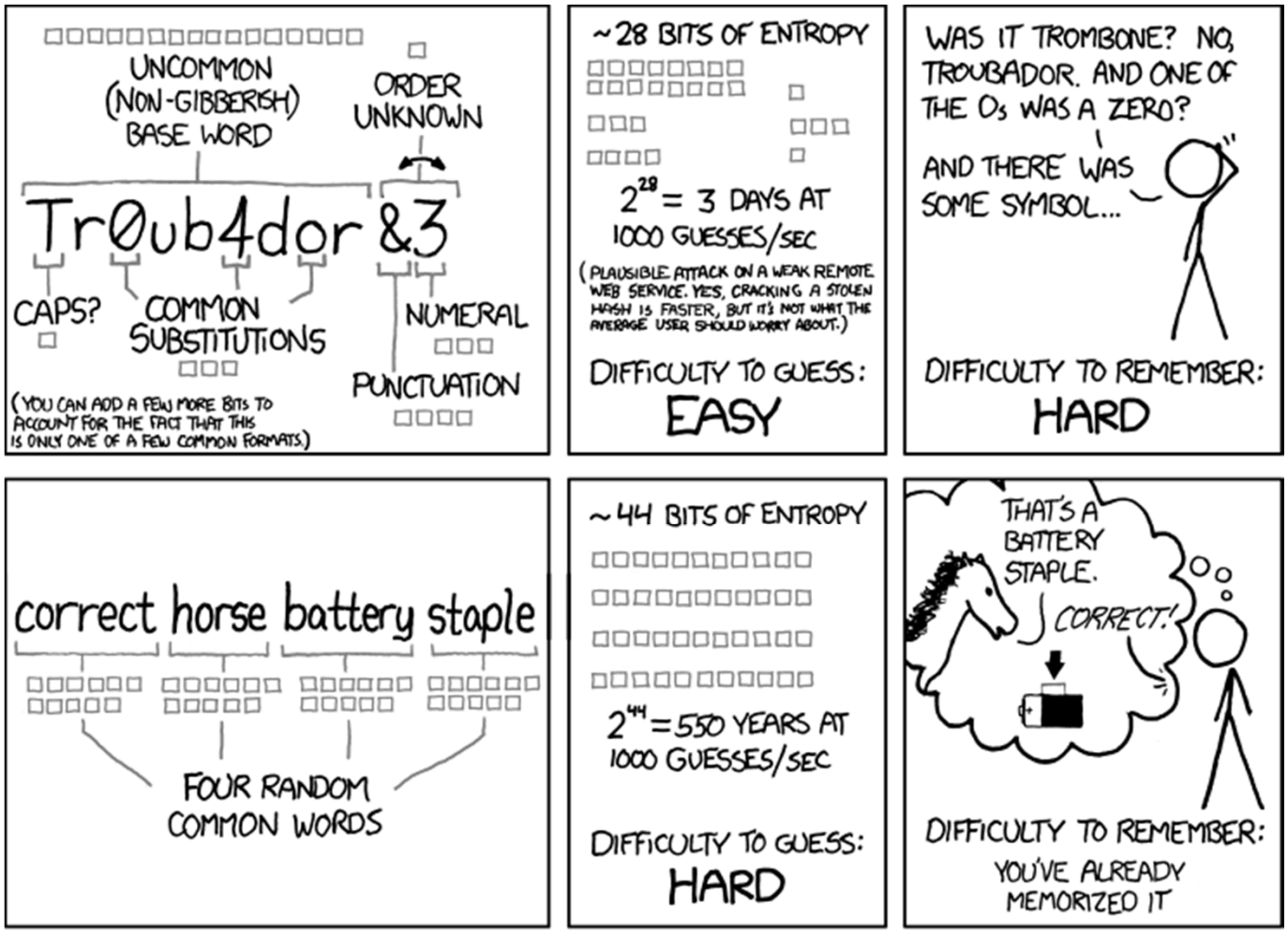
REVIEW…

Entropy & Coding: Central Ideas

- Use short codes for highly likely events. This shortens the average length of coded messages.
- Code several events at a time. Provides greater flexibility in code design.
- Shannon's Source Coding Theorem
- Algorithms: Huffman Encoding, …
- Applications: Compression…

REVIEW…

xkcd(#936): Password Strength



Some Observations

- Successive symbols from a source are not always independent. E.g. in english,

  h is more likely to occur following a *t*.

- This **intersymbol dependency** must be accounted for in an accurate measure of entropy.

Lossless Compression: English Text

- How much can we compress a given text?
- What is the accurate measure of entropy of English texts?
- Zeroth-Order entropy

$$H_0 \leq \log\left(\frac{1}{27}\right)$$
$$\leq 4.755\ bits/letter$$

- First-Order Entropy:

$$H_1 = 3.3$$

- Second-Order Entropy:

$$H_2 = 3.1$$

Zipf's Law, 1949

$$P_n \sim 1/n^a$$

P_n is the frequency of occurrence of the n^{th} ranked item and a is close to 1.

- The most frequent word will occur approximately twice as often as the second most frequent word, three times as often as the third most frequent word, etc.

- For example, in a corpus of over 1 million words:

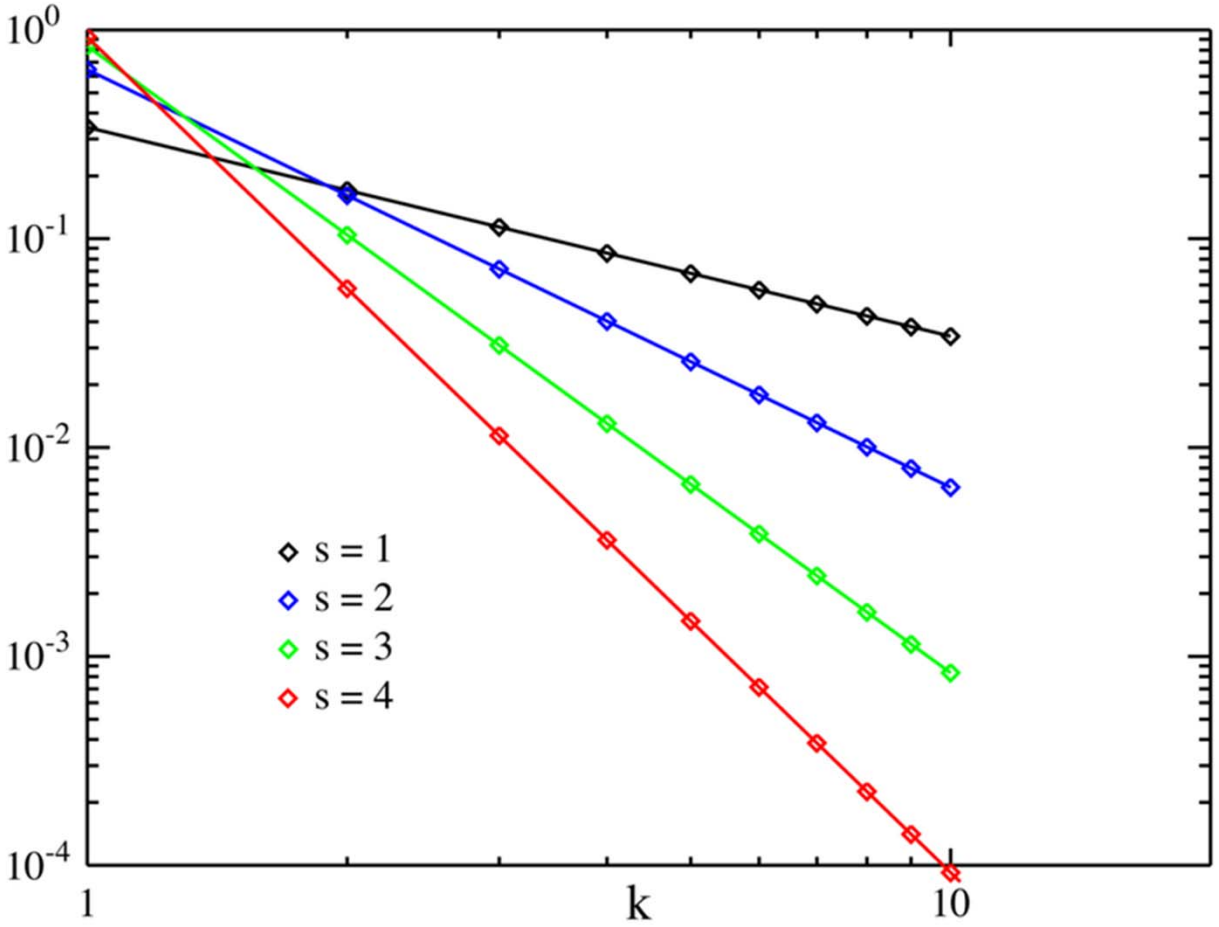
the	69,971	7%
of	36.411	3.5%
and	28,852	2.9%

- For English text:

$$p_m = \frac{A}{m}$$

where m is the word's rank, p_m is the probability of the m^{th} rank word,
A is a constant that depends on the number of active words in the language.

Zipf's Law



Estimating Entropy of English Text

- For English text, W with M words:

$$p_m = \frac{A}{m}$$

where m is the word's rank, p_m is the probability of the m^{th} rank word,
A is a constant that depends on the number of active words in the language.

- If $A = 0.1$, so that

$$\sum_{m=1}^{M} p_m = 1$$

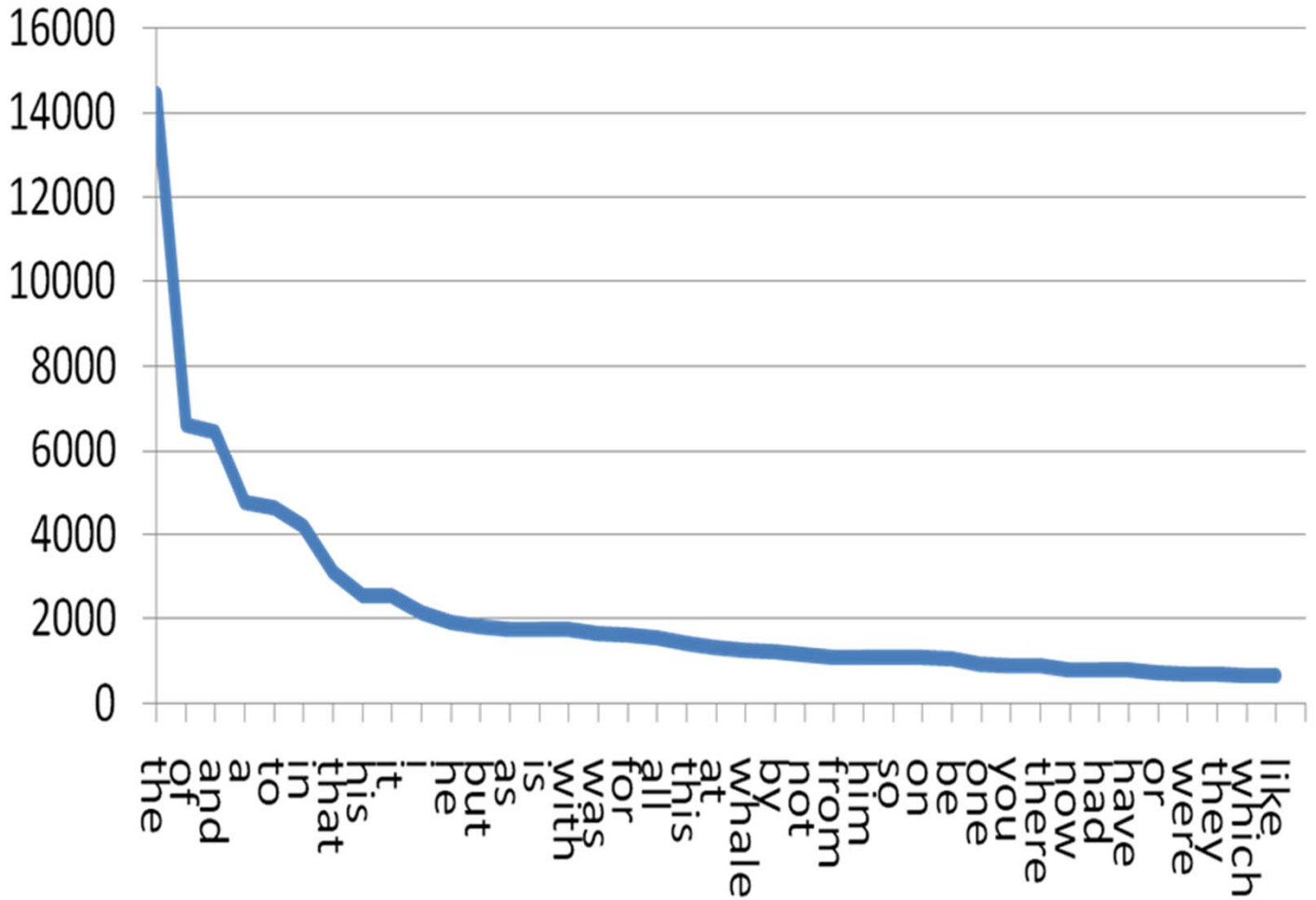
We need, $M = 12{,}366$.

$$H_W = \sum_{m=1}^{m=12{,}366} \frac{.01}{m} \log\left(\frac{m}{.01}\right) = 9.72\ bits/word$$

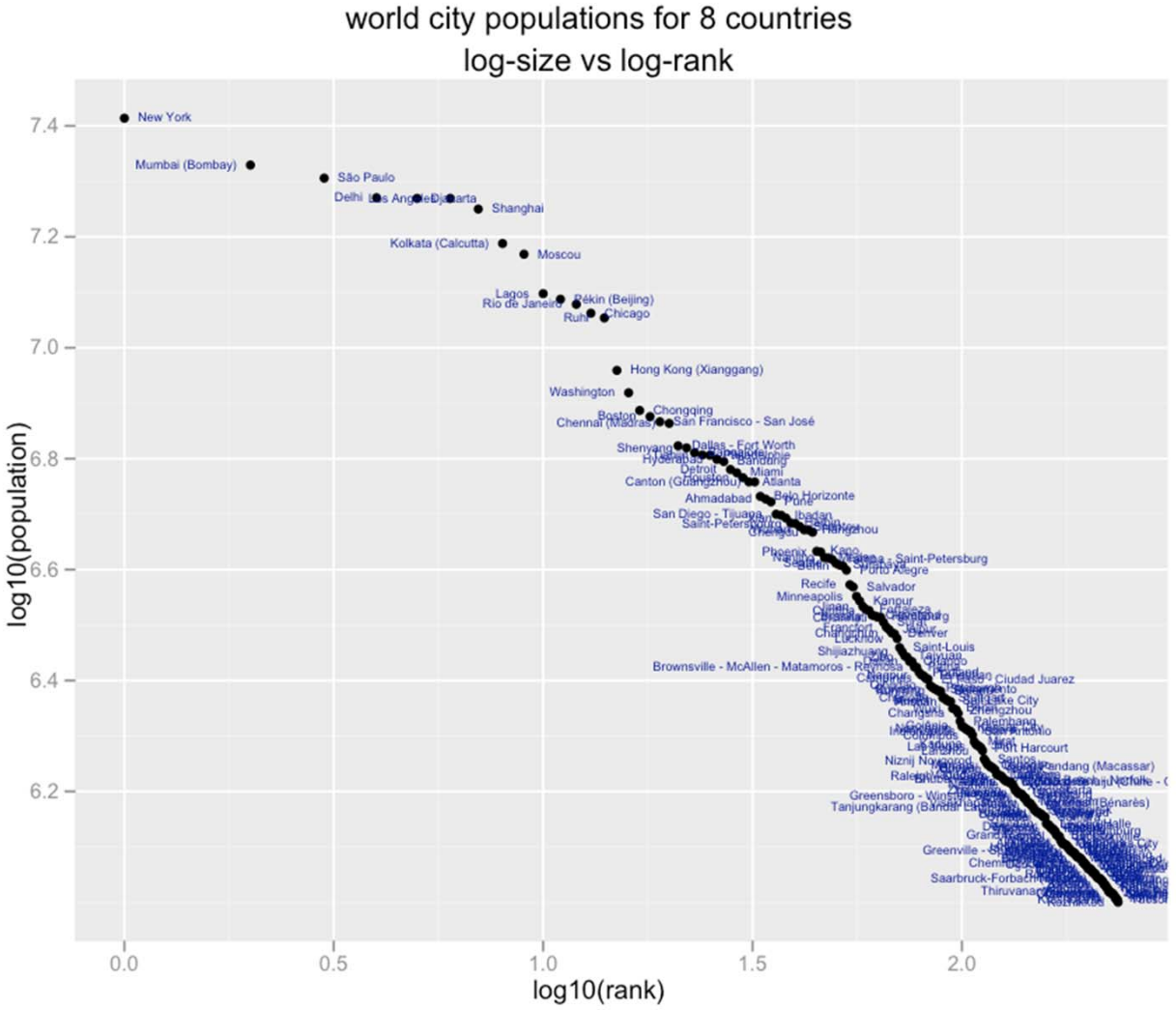
If $\bar{w}$=4.5 letters/word

$$H = \frac{9.72}{4.5} = 2.16 \text{bits/word}$$

The “Long Tail” of Moby Dick



Zipf's Law and Populations



Shannon Redundancy

$$R = 1 - \frac{H}{logM}$$

Where H is the per letter entropy, M is the size of the source alphabet. Thus redundancy of English is

$$1 - \frac{2.16}{\log 27} = 54.6\%$$

With an entropy of 1.5 we get 67% redundancy.

I.e. Huffman coding (even with an entropy of 3.3 or 3.1) will not get close to the theoretical limit.

Can we achieve compression rates close to 33%???

Lempel-Ziv Coding

- Sequences of text repeat patterns (words, phrases, etc)
- Construct a dictionary of common patterns
- Send references to patterns as triples (x, y, z)

Lempel-Ziv Coding (LZ77)

Message				Search Buffer	Look-Ahead Buffer
					THIS-THESIS-IS-THE-THESIS.

Lempel-Ziv Coding (LZ77)

Message				Search Buffer	Look-Ahead Buffer
					THIS-THESIS-IS-THE-THESIS.
0	0	T	T		THIS-THESIS-IS-THE-THESIS.

Lempel-Ziv Coding (LZ77)

Message				Search Buffer	Look-Ahead Buffer
					THIS-THESIS-IS-THE-THESIS.
0	0	T	T		THIS-THESIS-IS-THE-THESIS.
0	0	H	H		THIS-THESIS-IS-THE-THESIS.

Lempel-Ziv Coding (LZ77)

Message				Search Buffer	Look-Ahead Buffer
					THIS-THESIS-IS-THE-THESIS.
0	0	T	T		THIS-THESIS-IS-THE-THESIS.
0	0	H	H		THIS-THESIS-IS-THE-THESIS.
0	0	I	I		THIS-THESIS-IS-THE-THESIS.

Lempel-Ziv Coding (LZ77)

Message				Search Buffer	Look-Ahead Buffer
					THIS-THESIS-IS-THE-THESIS.
0	0	T	T		THIS-THESIS-IS-THE-THESIS.
0	0	H	H		THIS-THESIS-IS-THE-THESIS.
0	0	I	I		THIS-THESIS-IS-THE-THESIS.
0	0	S	S		THIS-THESIS-IS-THE-THESIS.

Lempel-Ziv Coding (LZ77)

Message				Search Buffer	Look-Ahead Buffer
					THIS-THESIS-IS-THE-THESIS.
0	0	T	T		THIS-THESIS-IS-THE-THESIS.
0	0	H	H		THIS-THESIS-IS-THE-THESIS.
0	0	I	I		THIS-THESIS-IS-THE-THESIS.
0	0	S	S		THIS-THESIS-IS-THE-THESIS.
0	0	-	-		THIS-THESIS-IS-THE-THESIS.

Lempel-Ziv Coding (LZ77)

Message				Search Buffer	Look-Ahead Buffer
					THIS-THESIS-IS-THE-THESIS.
0	0	T	T		THIS-THESIS-IS-THE-THESIS.
0	0	H	H		THIS-THESIS-IS-THE-THESIS.
0	0	I	I		THIS-THESIS-IS-THE-THESIS.
0	0	S	S		THIS-THESIS-IS-THE-THESIS.
0	0	-	-		THIS-THESIS-IS-THE-THESIS.
5	2	E	E	THIS-THE	SIS-IS-THE-THESIS.

Lempel-Ziv Coding (LZ77)

Message				Search Buffer	Look-Ahead Buffer
					THIS-THESIS-IS-THE-THESIS.
0	0	T	T		THIS-THESIS-IS-THE-THESIS.
0	0	H	H	T	HIS-THESIS-IS-THE-THESIS.
0	0	I	I	TH	IS-THESIS-IS-THE-THESIS.
0	0	S	S	THI	S-THESIS-IS-THE-THESIS.
0	0	-	-	THIS	-THESIS-IS-THE-THESIS.
5	2	E	THE	THIS-	THESIS-IS-THE-THESIS.
5	1	I	SI	THIS-THE	SIS-IS-THE-THESIS.

Lempel-Ziv Coding (LZ77)

Message				Search Buffer	Look-Ahead Buffer
					THIS-THESIS-IS-THE-THESIS.
0	0	T	T		THIS-THESIS-IS-THE-THESIS.
0	0	H	H	T	HIS-THESIS-IS-THE-THESIS.
0	0	I	I	TH	IS-THESIS-IS-THE-THESIS.
0	0	S	S	THI	S-THESIS-IS-THE-THESIS.
0	0	-	-	THIS	-THESIS-IS-THE-THESIS.
5	2	E	THE	THIS-	THESIS-IS-THE-THESIS.
5	1	I	SI	THIS-THE	SIS-IS-THE-THESIS.
7	2	I	S-I	THIS-THESI	S-IS-THE-THESIS.

Lempel-Ziv Coding (LZ77)

Message				Search Buffer	Look-Ahead Buffer
					THIS-THESIS-IS-THE-THESIS.
0	0	T	T		THIS-THESIS-IS-THE-THESIS.
0	0	H	H	T	HIS-THESIS-IS-THE-THESIS.
0	0	I	I	TH	IS-THESIS-IS-THE-THESIS.
0	0	S	S	THI	S-THESIS-IS-THE-THESIS.
0	0	-	-	THIS	-THESIS-IS-THE-THESIS.
5	2	E	THE	THIS-	THESIS-IS-THE-THESIS.
5	1	I	SI	THIS-THE	SIS-IS-THE-THESIS.
7	2	I	S-I	THIS-THESI	S-IS-THE-THESIS.
10	5	-	S-THE-	THIS-THESIS-I	S-THE-THESIS.
				THIS-THESIS-IS-THE-	THESIS.

Lempel-Ziv Coding (LZ77)

Message				Search Buffer	Look-Ahead Buffer
					THIS-THESIS-IS-THE-THESIS.
0	0	T	T		THIS-THESIS-IS-THE-THESIS.
0	0	H	H	T	HIS-THESIS-IS-THE-THESIS.
0	0	I	I	TH	IS-THESIS-IS-THE-THESIS.
0	0	S	S	THI	S-THESIS-IS-THE-THESIS.
0	0	-	-	THIS	-THESIS-IS-THE-THESIS.
5	2	E	THE	THIS-	THESIS-IS-THE-THESIS.
5	1	I	SI	THIS-THE	SIS-IS-THE-THESIS.
7	2	I	S-I	THIS-THESI	S-IS-THE-THESIS.
10	5	-	S-THE-	THIS-THESIS-I	S-THE-THESIS.
14	6	.	THESIS.	THIS-THESIS-IS-THE-	THESIS.
				THIS-THESIS-IS-THE-THESIS.	

Lempel-Ziv Coding

- Sequences of text repeat patterns (words, phrases, etc)
- Construct a dictionary of common patterns
- Send references to patterns as triples (x, y, z)
  e.g. $(5, 3, F)$
  go back 5 received chars
  take the next 3 from there
  add F to the end
- Size of Search Buffer and Look-Ahead Buffer is finite.
- Used by ZIP, PKSip, Lharc, PNG, gzip, ARJ
- Extended to LZ78 (uses dictionary), LZW (+Terry Welch)
- Achieves optimal rate of transmission in the long run w/o using probability dist.

Decode

Message

0	0	I	
0	0	-	
0	0	M	
3	1	S	
1	1	-	
5	5	L	
5	3	Y	

Decode

Message

0	0	I	I
0	0	-	
0	0	M	
3	1	S	
1	1	-	
5	5	L	
5	3	Y	

Decode

Message

0	0	I	I
0	0	-	I -
0	0	M	
3	1	S	
1	1	-	
5	5	L	
5	3	Y	

Decode

			Message
0	0	I	I
0	0	-	I -
0	0	M	I - M
3	1	S	
1	1	-	
5	5	L	
5	3	Y	

Decode

			Message
0	0	I	I
0	0	-	I -
0	0	M	I - M
3	1	S	I - M I S
1	1	-	
5	5	L	
5	3	Y	

Decode

			Message
0	0	I	I
0	0	-	I -
0	0	M	I - M
3	1	S	I - M I S
1	1	-	I - M I S S -
5	5	L	
5	3	Y	

Decode

			Message
0	0	I	I
0	0	-	I -
0	0	M	I - M
3	1	S	I - M I S
1	1	-	I - M I S S -
5	5	L	I - M I S S - M I S S - L
5	3	Y	

Decode

			Message
0	0	I	I
0	0	-	I -
0	0	M	I - M
3	1	S	I - M I S
1	1	-	I - M I S S -
5	5	L	I - M I S S - M I S S - L
5	3	Y	I - M I S S - M I S S - L I S S Y

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