

# Forward & Backward Chaining Inference

November 2021

1

## FOPC – Inference Rules

- Inference rules for Quantifiers ( $\forall$ ,  $\exists$ )
  - Universal Instantiation
  - Existential Instantiation
- Generalized Modus Ponens
  - Unification
- Forward & Backward Chaining
  - Definite Clauses
  - Logic Programming in Prolog
- Resolution
  - Reductio ad Absurdum

2

2

## Generalized Modus Ponens

For atomic sentences  $p_i, p'_i$ , and  $q$

where there is a substitution  $\theta$  such that

$\text{SUBST}(\theta, p'_i) = \text{SUBST}(\theta, p_i)$ , for all  $i$

$p'_1$

$p'_2$

...

$p'_n$

$p'_1 \wedge p'_2 \wedge \dots \wedge p'_n \Rightarrow q$

---

$\text{SUBST}(\theta, q)$

e.g.

King(John)

Greedy(John)

$\forall x [\text{King}(x) \wedge \text{Greedy}(x) \Rightarrow \text{Evil}(x)]$

$\theta$  is  $\{x/\text{John}\}$

$\text{SUBST}(\theta, \text{Evil}(x))$  is Evil(John)

3

3

## Unification - $\text{UNIFY}(p, q) = \theta$

How to find substitutions that make different logical expressions look identical?

For wffs  $p$  and  $q$  (sentences with universally quantified variables)

$\text{UNIFY}(p, q) = \theta$

e.g.  $p = \text{Knows}(\text{John}, x)$   
 $q = \text{Knows}(\text{John}, \text{Mary})$

$\theta = \{x/\text{Mary}\}$

$p = \text{Knows}(\text{John}, x)$   
 $q = \text{Knows}(y, \text{Bill})$

$\theta = \{y/\text{John}, x/\text{Bill}\}$

$p = \text{Knows}(\text{John}, x)$   
 $q = \text{Knows}(x, \text{Elizabeth})$

$\theta = \{ \}$  i.e. Fail! (no unification)

$p = \text{Knows}(\text{John}, x_1)$   
 $q = \text{Knows}(x_2, \text{Elizabeth})$

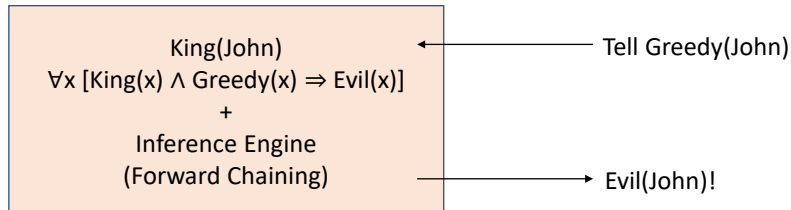
$\theta = \{x_1/\text{Elizabeth}, x_2/\text{John}\}$

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4

## Forward Chaining Inference

- Tell-Ask Systems

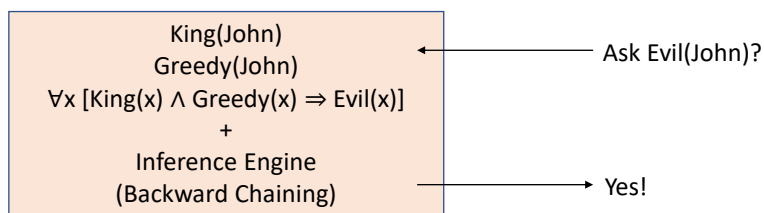


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## Backward Chaining Inference

- Tell-Ask Systems

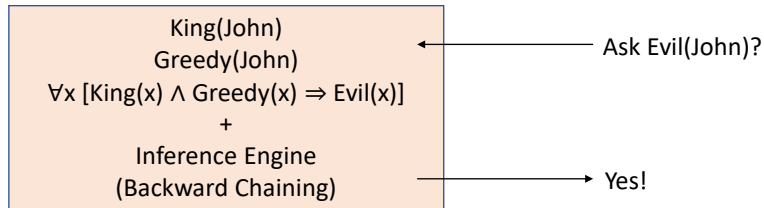


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## Backward Chaining Inference

- Tell-Ask Systems



- **Requires wffs to be in Definite Clause form!**

7

7

## Forward Chaining

```
American(x)  $\wedge$  Weapon(y)  $\wedge$  Sells(x, y, z)  $\wedge$  Hostile(z)  $\Rightarrow$  Criminal(x)
Owns(Nono, M1)
Missile(M1)
Missile(x)  $\wedge$  Owns(Nono, x)  $\Rightarrow$  Sells(CWest, x, Nono)
Missile(x)  $\Rightarrow$  Weapon(x)
Enemy(x, America)  $\Rightarrow$  Hostile(x)
American(Cwest)
Enemy(Nono, America)
```

American(Cwest)

Missile(M1)

Owns(Nono, M1)

Enemy(Nono, America)

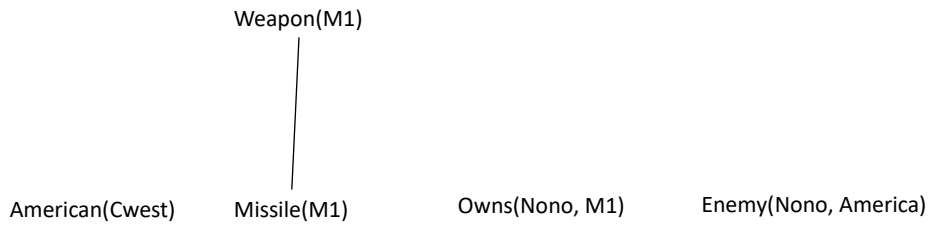
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8

# Forward Chaining

```

American(x) ∧ Weapon(y) ∧ Sells(x, y, z) ∧ Hostile(z) ⇒ Criminal(x)
Owns(Nono, M1)
Missile(M1)
Missile(x) ∧ Owns(Nono, x) ⇒ Sells(CWest, x, Nono)
Missile(x) ⇒ Weapon(x)
Enemy(x, America) ⇒ Hostile(x)
American(Cwest)
American(Cwest)
Enemy(Nono, America)
    
```



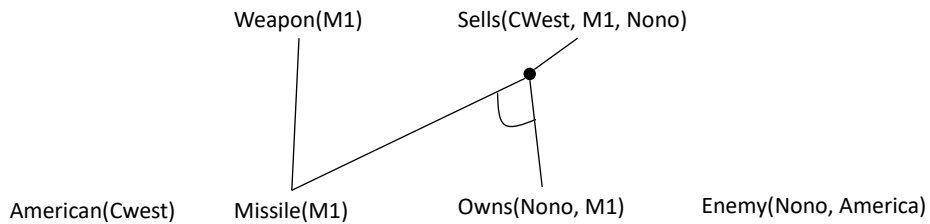
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# Forward Chaining

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American(x) ∧ Weapon(y) ∧ Sells(x, y, z) ∧ Hostile(z) ⇒ Criminal(x)
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American(Cwest)
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```

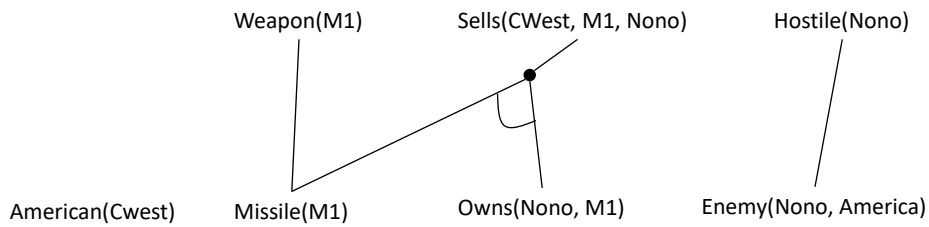


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# Forward Chaining

$American(x) \wedge Weapon(y) \wedge Sells(x, y, z) \wedge Hostile(z) \Rightarrow Criminal(x)$   
 $Owns(Nono, M1)$   
 $Missile(M1)$   
 $Missile(x) \wedge Owns(Nono, x) \Rightarrow Sells(CWest, x, Nono)$   
 $Missile(x) \Rightarrow Weapon(x)$   
 $Enemy(x, America) \Rightarrow Hostile(x)$   
 $American(Cwest)$   
 $Enemy(Nono, America)$

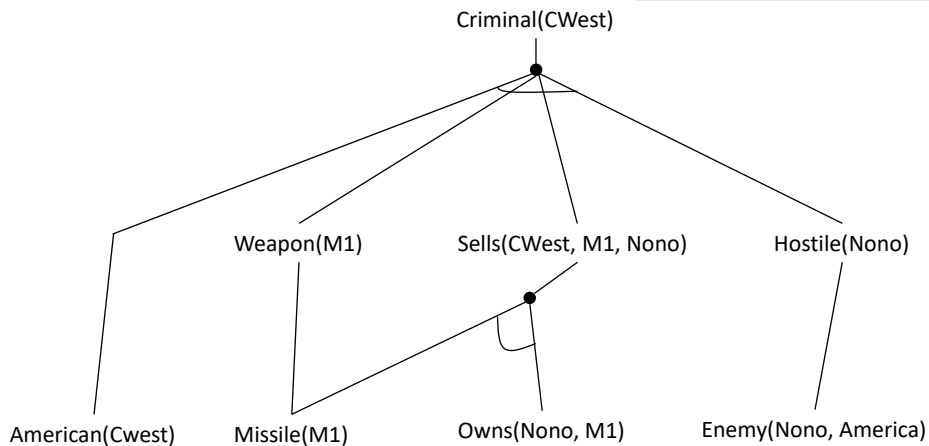


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11

# Forward Chaining

$American(x) \wedge Weapon(y) \wedge Sells(x, y, z) \wedge Hostile(z) \Rightarrow Criminal(x)$   
 $Owns(Nono, M1)$   
 $Missile(M1)$   
 $Missile(x) \wedge Owns(Nono, x) \Rightarrow Sells(CWest, x, Nono)$   
 $Missile(x) \Rightarrow Weapon(x)$   
 $Enemy(x, America) \Rightarrow Hostile(x)$   
 $American(Cwest)$   
 $Enemy(Nono, America)$



12

12

# Backward Chaining

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American(x) ∧ Weapon(y) ∧ Sells(x, y, z) ∧ Hostile(z) ⇒ Criminal(x)
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Enemy(Nono, America)

```

Criminal(CWest)?

13

13

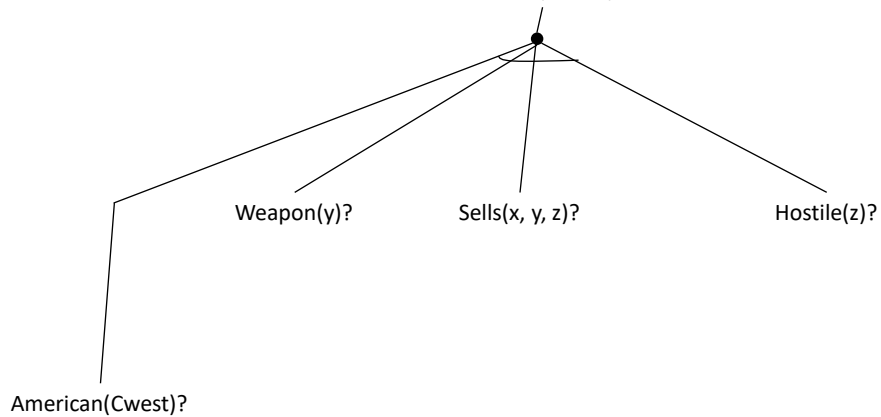
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Enemy(America, x) ⇒ Hostile(x)
American(Cwest)
Enemy(Nono, America)

```

Criminal(CWest)?



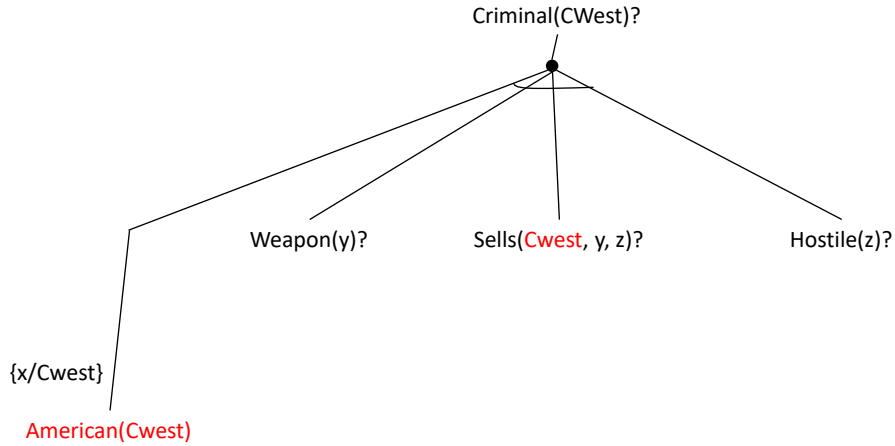
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# Backward Chaining

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American(x) ∧ Weapon(y) ∧ Sells(x, y, z) ∧ Hostile(z) ⇒ Criminal(x)
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```



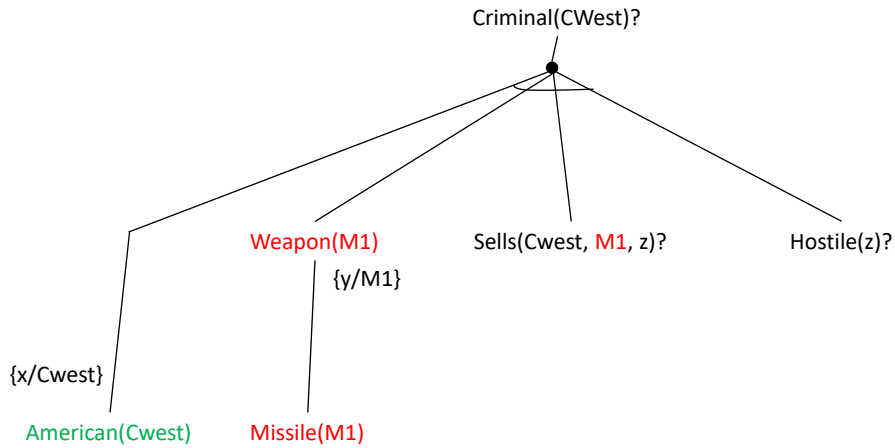
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American(Cwest)
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16

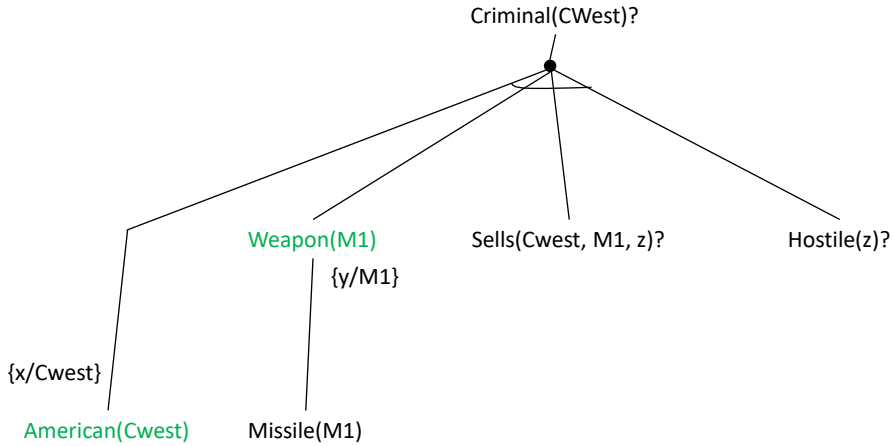
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# Backward Chaining

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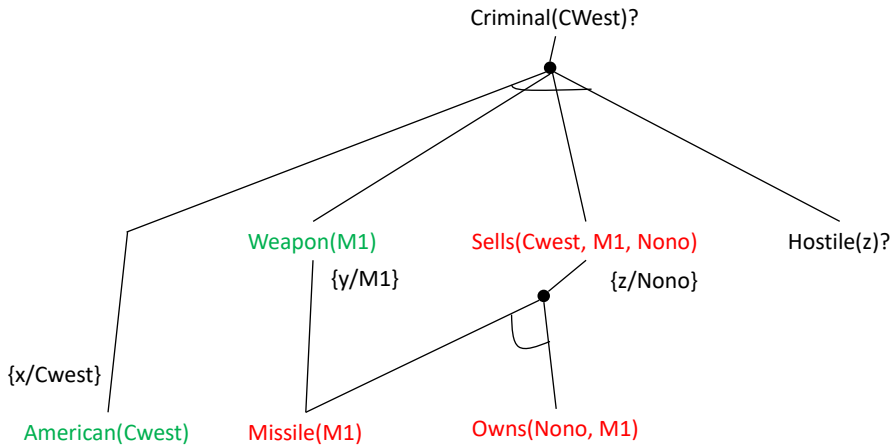
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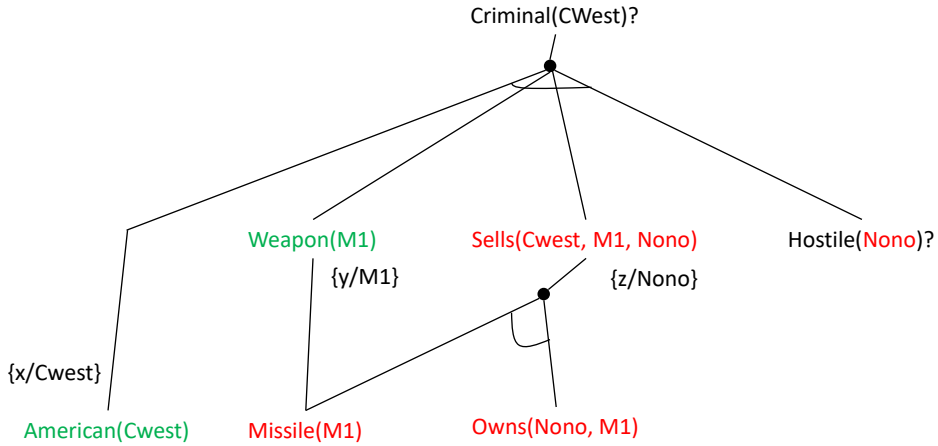
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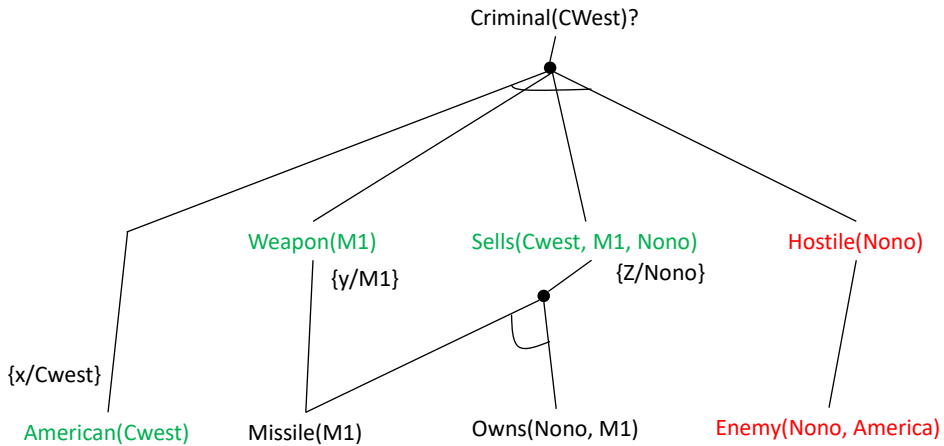
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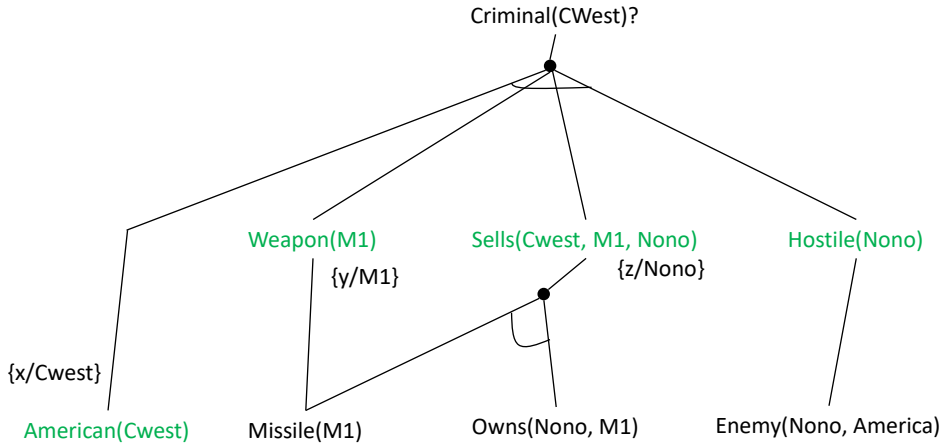


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American(x) ∧ Weapon(y) ∧ Sells(x, y, z) ∧ Hostile(z) ⇒ Criminal(x)
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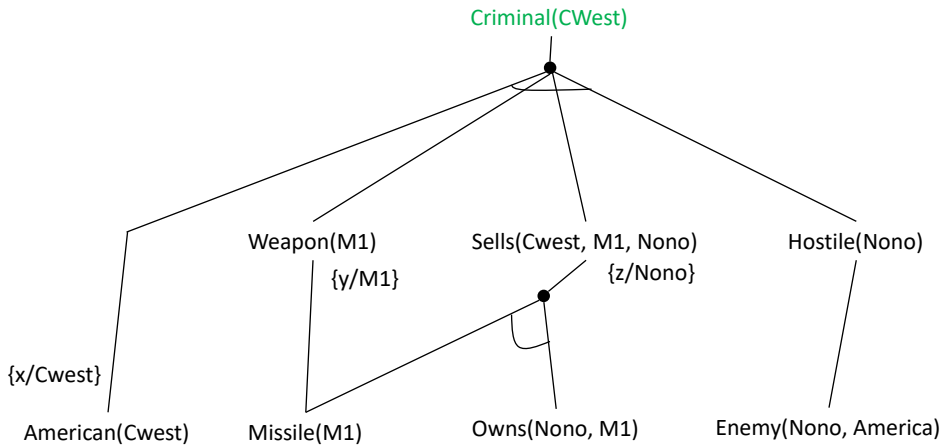


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21

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```



22

22

# Reasoning about states & actions - Planning

Deepak Kumar  
November 2021

23

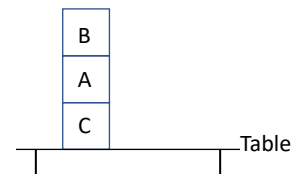
## Blocksworld

- **Objects** A, B, C, Table

- **Relations** On<sup>2</sup>, Clear<sup>1</sup>

On(C, Table)    On(A, C)    On(B, A)  
Clear(B)        Clear(A)    Clear(C)    Clear(Table)

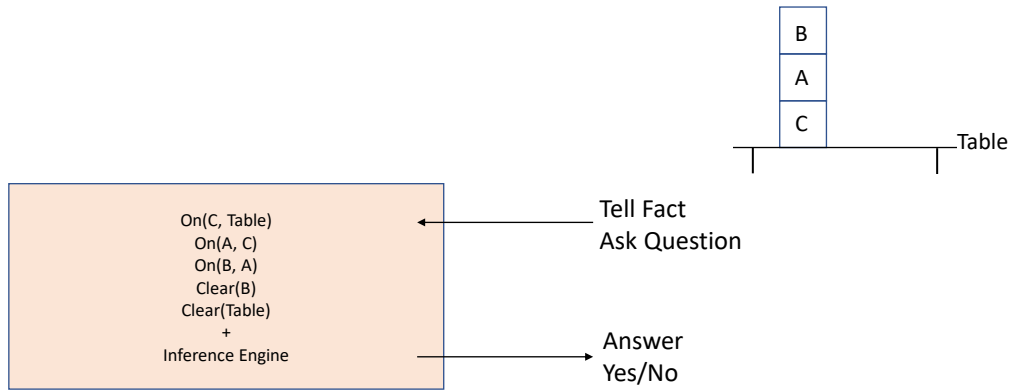
- Example Knowledge Base

$$\Delta = \{\text{On}(C, \text{Table}), \text{On}(A, C), \text{On}(B, A), \text{Clear}(B), \text{Clear}(\text{Table})\}$$


24

24

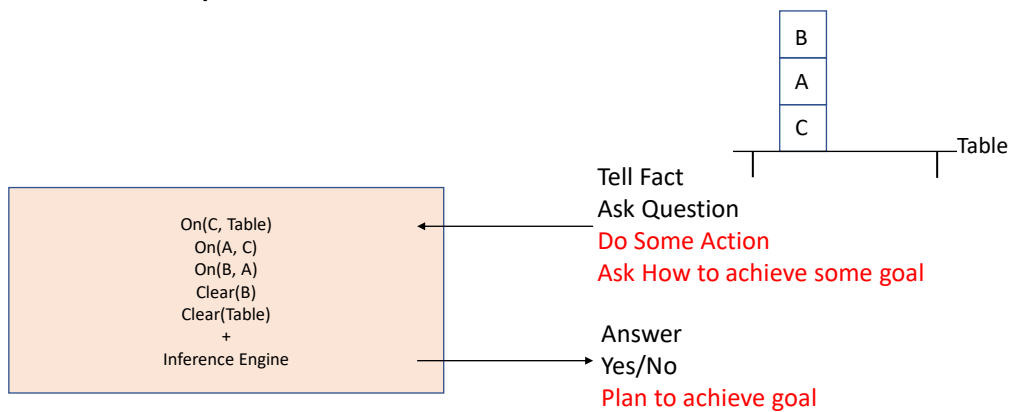
## Tell-Ask Systems



25

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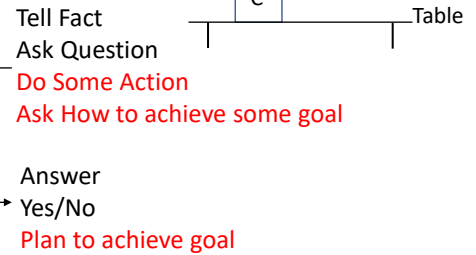
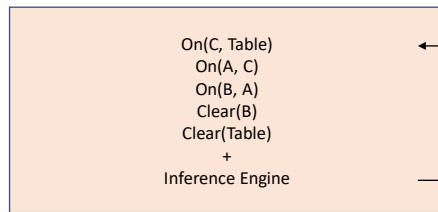
## Tell-Ask-Do Systems



26

26

## Tell-Ask-Do Systems



Can specify a goal as a wff:

Achieve  $\text{On}(A, B) \wedge \text{On}(B, C) \wedge \text{On}(C, \text{Table})$

Achieve  $\text{On}(A, B)$

Achieve  $\exists x \text{On}(x, B)$

Etc.

27

27

## Reasoning about states & actions

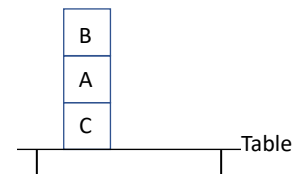
Can specify a goal as a wff:

Achieve  $\text{On}(A, B) \wedge \text{On}(B, C) \wedge \text{On}(C, \text{Table})$

Achieve  $\text{On}(A, B)$

Achieve  $\exists x \text{On}(x, B)$

Etc.



We can test if a goal is satisfied by the KB:

$\Delta \models_R^? \text{Goal}$

But, how to find a set/sequence of actions for achieving a goal satisfied by a wff?

28

## Reasoning about states & actions

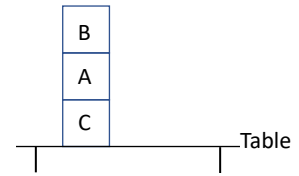
Can specify a goal as a wff:

Achieve  $\text{On}(A, B) \wedge \text{On}(B, C) \wedge \text{On}(C, \text{Table})$

Achieve  $\text{On}(A, B)$

Achieve  $\exists x \text{On}(x, B)$

Etc.



We can test if a goal is satisfied by the KB:

$\Delta \models_R^? \text{Goal}$

But, how to find a set/sequence of actions for achieving a goal satisfied by a wff?

Search a space of logical expressions/state descriptions...

29

## Situation Calculus

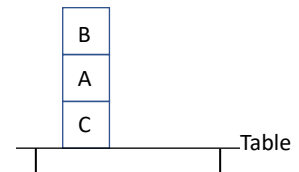
- McCarthy & Hayes, 1969
- Went out of fashion, but made a come back – Reiter 2001, Lakemeyer, 2014.

- If  $S_0$  is the state:

$S_0 = \{\text{On}(B, A), \text{On}(A, C), \text{On}(C, \text{Table}), \dots\}$

- Add the state to the relation:

$\{\text{On}(A, B, S_0), \text{On}(A, C, S_0), \text{On}(C, \text{Table}, S_0), \dots\}$



30

## Situation Calculus

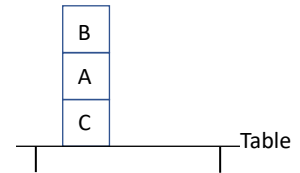
- McCarthy & Hayes, 1969
- Went out of fashion, but made a come back – Reiter 2001, Lakemeyer, 2014.
- If so is the state:

$$S_0 = \{\text{On}(B, A), \text{On}(A, C), \text{On}(C, \text{Table}), \dots\}$$

- Add the state to the relation:

$$\{\text{On}(A, B, S_0), \text{On}(A, C, S_0), \text{On}(C, \text{Table}, S_0) \dots\}$$

These are called Fluents.



31

## Situation Calculus

- Add the state to the relation:

$$\{\text{On}(A, B, S_0), \text{On}(A, C, S_0), \text{On}(C, \text{Table}, S_0) \dots\}$$

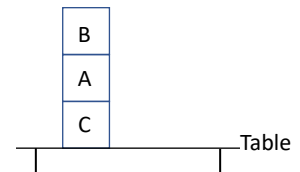
- Now, we can make statements about states:

$$\forall x \forall y \forall s [\text{On}(x, y, s) \wedge \neg(y = \text{Table}) \Rightarrow \neg \text{Clear}(y, s)]$$

$$\forall s \text{Clear}(\text{Table}, s)$$

- We can now prove:

$$\begin{aligned} &\neg \text{Clear}(A, S_0) \\ &\text{Clear}(\text{Table}, S_0) \end{aligned}$$



32



## Representing Actions – Action Schema

- **Action Schema:**  $\text{move}(x, y, z)$  - move  $x$ , from  $y$ , to  $z$
- **Function do():** if  $\alpha$  is an action and  $\sigma$  is a state then  
 $\text{do}(\alpha, \sigma) \rightarrow \sigma'$   
 $\sigma'$  is the state resulting from doing  $\alpha$  in  $\sigma$
- **Effects of actions** represented by wffs:  
 $\forall x \forall y \forall s \forall z [\text{On}(x, y, s) \wedge \text{Clear}(x, s) \wedge \text{Clear}(z, s) \wedge x \neq z \Rightarrow \text{On}(x, z, \text{do}(\text{move}(x, y, z), s))]$   
 $\forall x \forall y \forall s \forall z [\text{On}(x, y, s) \wedge \text{Clear}(x, s) \wedge \text{Clear}(z, s) \wedge x \neq z \Rightarrow \neg \text{On}(x, y, \text{do}(\text{move}(x, y, z), s))]$   
 These are positive(+) and negative(-) axioms.
- Need these for every fluent-action pair!

33

## Axioms for Clear and move

- $$\forall x \forall y \forall s \forall z [\text{On}(x, y, s) \wedge \text{Clear}(x, s) \wedge \text{Clear}(z, s) \wedge (x \neq z) \wedge (y \neq z) \Rightarrow \text{Clear}(y, \text{do}(\text{move}(x, y, z), s))]$$
- $$\forall x \forall y \forall s \forall z [\text{On}(x, y, s) \wedge \text{Clear}(x, s) \wedge \text{Clear}(z, s) \wedge (x \neq z) \wedge (z \neq \text{Table}) \Rightarrow \neg \text{Clear}(z, \text{do}(\text{move}(x, y, z), s))]$$

34

## Axioms for Clear and move

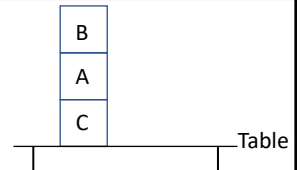
$$\forall x \forall y \forall s \forall z [\text{On}(x, y, s) \wedge \text{Clear}(x, s) \wedge \text{Clear}(z, s) \wedge (x \neq z) \wedge (y \neq z) \\ \Rightarrow \text{Clear}(y, \text{do}(\text{move}(x, y, z), s))]$$

$$\forall x \forall y \forall s \forall z [\text{On}(x, y, s) \wedge \text{Clear}(x, s) \wedge \text{Clear}(z, s) \wedge (x \neq z) \wedge (z \neq \text{Table}) \\ \Rightarrow \neg \text{Clear}(z, \text{do}(\text{move}(x, y, z), s))]$$

Preconditions & Effects

35

## Axioms for Clear and move



$$\forall x \forall y \forall s \forall z [\text{On}(x, y, s) \wedge \text{Clear}(x, s) \wedge \text{Clear}(z, s) \wedge (x \neq z) \wedge (y \neq z) \\ \Rightarrow \text{Clear}(y, \text{do}(\text{move}(x, y, z), s))]$$

$$\forall x \forall y \forall s \forall z [\text{On}(x, y, s) \wedge \text{Clear}(x, s) \wedge \text{Clear}(z, s) \wedge (x \neq z) \wedge (z \neq \text{Table}) \\ \Rightarrow \neg \text{Clear}(z, \text{do}(\text{move}(x, y, z), s))]$$

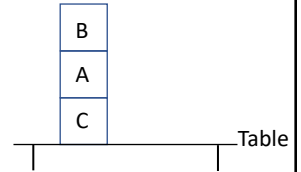
$\text{do}(\text{move}(B, A, \text{Table}), S_0)$

We get

- $\text{On}(B, \text{Table}, \text{do}(\text{move}(B, A, \text{Table}), S_0))$
- $\neg \text{On}(B, A, \text{do}(\text{move}(B, A, \text{Table}), S_0))$
- $\text{Clear}(A, \text{do}(\text{move}(B, A, \text{Table}), S_0))$

36

## Axioms for Clear and move



$$\forall x \forall y \forall z [\text{On}(x, y, s) \wedge \text{Clear}(x, s) \wedge \text{Clear}(z, s) \wedge (x \neq z) \wedge (y \neq z) \Rightarrow \text{Clear}(y, \text{do}(\text{move}(x, y, z), s))]$$

$$\forall x \forall y \forall z [\text{On}(x, y, s) \wedge \text{Clear}(x, s) \wedge \text{Clear}(z, s) \wedge (x \neq z) \wedge (z \neq \text{Table}) \Rightarrow \neg \text{Clear}(z, \text{do}(\text{move}(x, y, z), s))]$$

For  $\text{do}(\text{move}(B, A, \text{Table}), S_0)$

We get

$$\begin{aligned} &\text{On}(B, \text{Table}, \text{do}(\text{move}(B, A, \text{Table}), S_0)) \\ &\neg \text{On}(B, A, \text{do}(\text{move}(B, A, \text{Table}), S_0)) \\ &\text{Clear}(A, \text{do}(\text{move}(B, A, \text{Table}), S_0)) \end{aligned}$$

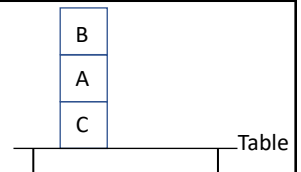
If  $S_1 = \text{do}(\text{move}(B, A, \text{Table}), S_0)$ , we have

$$\begin{aligned} &\text{On}(B, \text{Table}, S_1) \\ &\neg \text{On}(B, A, S_1) \\ &\text{Clear}(A, S_1) \end{aligned}$$

Plus, all wffs for  $S_0$  are still true!

37

## And...more axioms



- Not only do we have to say what changes, we have to specify what doesn't change...everything that doesn't change!

After  $\text{move}(B, A, \text{Table})$

C is still on the Table in  $S_0$  and  $S_1$

B is still clear in  $S_0$  and  $S_1$

...etc...

- Need more axioms...

38

## Frame Axioms

- Axioms for Move-On

$$\forall x \forall y \forall s \forall u \forall v \forall z [On(x, y, s) \wedge (x \neq u) \Rightarrow On(x, y, do(move(u, v, z), s))]$$

$$\forall x \forall y \forall s \forall u \forall v \forall z [-On(x, y, s) \wedge [(x \neq u) \vee (y \neq z)] \Rightarrow -On(x, y, do(move(u, v, z), s))]$$

- Axioms for Move-Clear

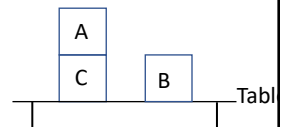
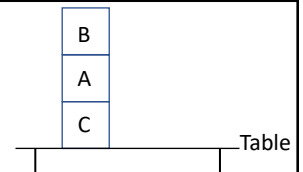
$$\forall x \forall y \forall s \forall z [Clear(u, s) \wedge (u \neq z) \Rightarrow Clear(u, do(move(x, y, z), s))]$$

$$\forall x \forall y \forall s \forall z [-Clear(u, s) \wedge (u \neq y) \Rightarrow -Clear(u, do(move(x, y, z), s))]$$

39

## Planning using Situation Calculus

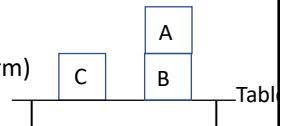
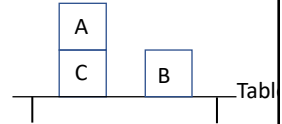
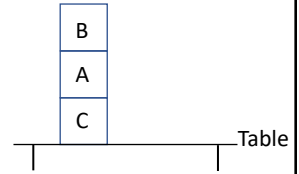
- To get a plan that achieves a goal:  $\Upsilon(s)$
- We need to prove:  $\exists s \Upsilon(s)$
- We'll get the plan by using Green's Trick:  $\exists s \Upsilon(s) \vee Answer(s)$
- Goal:  $\exists s On(B, Table, s)$
- Prove:  $\neg [\exists s On(B, Table, s)] \vee Answer(s)$
- i.e. Add the clause:  $\neg On(B, Floor, s) \vee Answer(s)$  to KB (in clause form)
- $Answer(do(move(B, A, Table), S_0))$



40

## Planning using Situation Calculus

- To get a plan that achieves a goal:  $\Upsilon(s)$
- We need to prove:  $\exists s \Upsilon(s)$
- We'll get the plan by using Green's Trick:  $\exists s \Upsilon(s) \vee \text{Answer}(s)$
- Goal:  $\exists s \text{On}(B, \text{Table}, s)$
- Example 1: Prove:  $\neg [\exists s \text{On}(B, \text{Table}, s)] \vee \text{Answer}(s)$   
i.e. Add the clause:  $\neg \text{On}(B, \text{Table}, s) \vee \text{Answer}(s)$  to KB (in clause form)  
 $\text{Answer}(\text{do}(\text{move}(B, A, \text{Table}), S_0))$
- Example 2: Prove:  $\neg [\exists s \text{On}(A, B, s) \wedge \text{On}(B, \text{Table}, s)] \vee \text{Answer}(s)$   
i.e. Add the clause:  $\neg \text{On}(A, B, s) \vee \neg \text{On}(B, \text{Table}, s) \vee \text{Answer}(s)$  to KB (in clause form)  
 $\text{Answer}(\text{do}(\text{move}(A, C, B), \text{do}(\text{move}(B, A, \text{Table}), S_0)))$



41

41

## Situation Calculus Planning - Summary

- **Action Schema:**  $\text{move}(x, y, z)$  - move  $x$ , from  $y$ , to  $z$
- **Function do():**  $\text{do}(\alpha, \sigma) \rightarrow \sigma'$
- **Effects Axioms**  
 $\forall x \forall y \forall s \forall z [\text{On}(x, y, s) \wedge \text{Clear}(x, s) \wedge \text{Clear}(z, s) \wedge x \neq z \Rightarrow \text{On}(x, z, \text{do}(\text{move}(x, y, z), s))]$   
 $\forall x \forall y \forall s \forall z [\text{On}(x, y, s) \wedge \text{Clear}(x, s) \wedge \text{Clear}(z, s) \wedge x \neq z \Rightarrow \neg \text{On}(x, y, \text{do}(\text{move}(x, y, z), s))]$   
 Need these for every fluent-action pair!
- Also need Frame Axioms
- Prove, using Green's Trick to extract plan:  $\exists s \Upsilon(s) \vee \text{Answer}(s)$
- Using resolution on:  $\neg[\exists s \Upsilon(s) \vee \text{Answer}(s)]$

42

## Problems with Situation Calculus Planning

- Too many Frame Axioms
- Proof effort is too large for even simple problems
- There are more problems:
  - Frame Problem
  - Qualifications Problem
  - Ramification Problem
  - Etc.

43

44