Mergesort and Quicksort

Sorting algorithms

- Insertion, selection and bubble sort have quadratic worst-case performance
- Mergesort and Quicksort have time O(n lg n)

Merge Sort

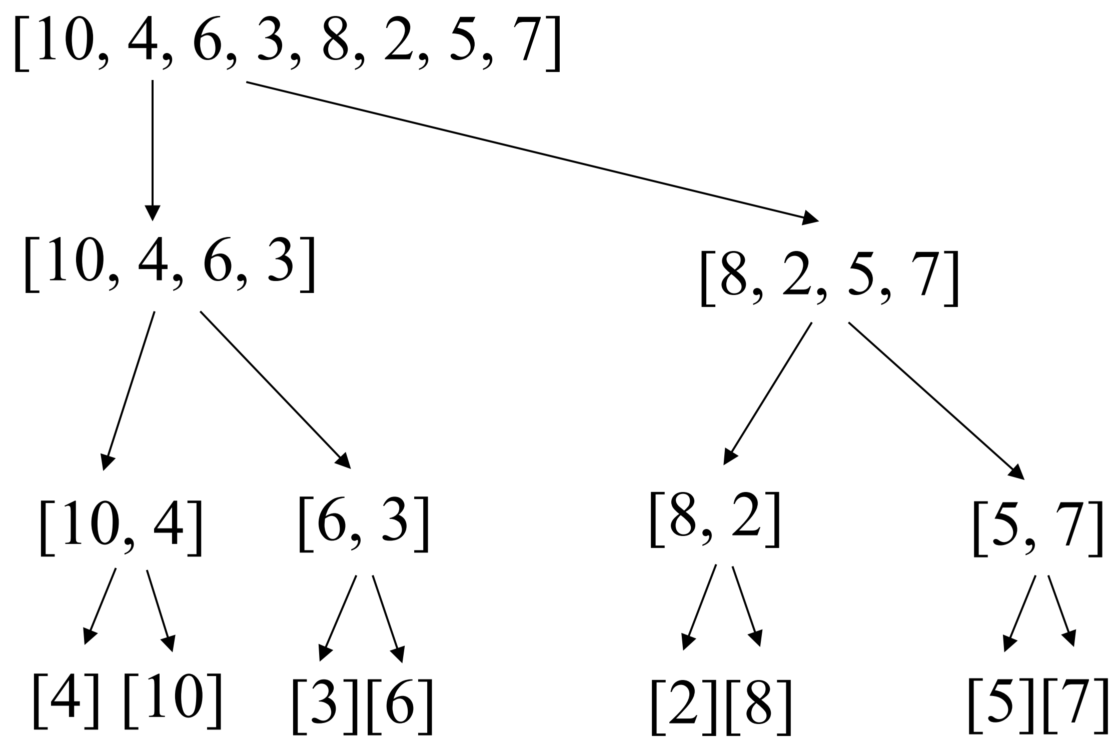
- Apply divide-and-conquer to sorting problem
- Problem: Given *n* elements, sort elements into non-decreasing order
- Divide-and-Conquer:
  - If n=1 terminate (every one-element list is already sorted)
  - If n>1, partition elements into two or more sub-collections; sort each; combine into a single sorted list
- How do we partition?

Partitioning

- Let' s try to achieve balanced partitioning
- A gets *n/2* elements, B gets rest half
- Sort A and B recursively
- Combine sorted A and B using a process called *merge*, which combines two sorted lists into one

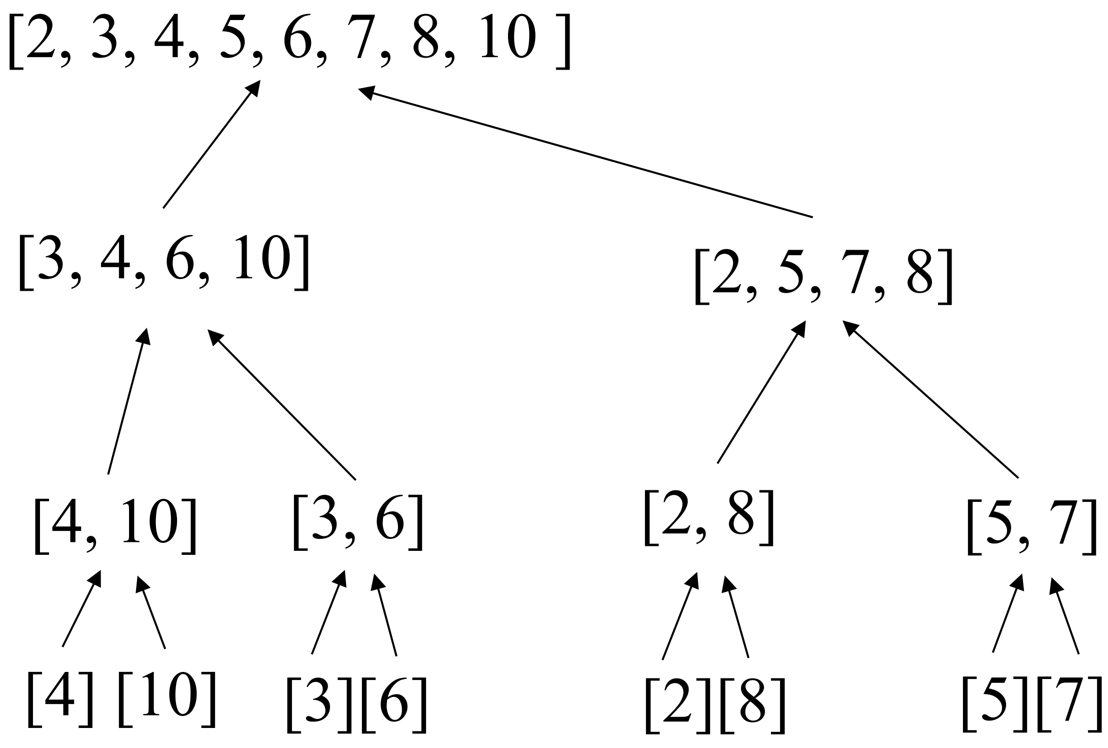
Example

- Partition into lists of size n/2



Example Cont' d

- Merge



Evaluation

- Recurrence equation:
- Assume n is a power of 2

$$T(n) = \begin{cases} c_1 & \text{if } n=1 \\ 2T(n/2) + c_2 n & \text{if } n>1,\ n=2^k \end{cases}$$

Solution

By Substitution:

$T(n) = 2T(n/2) + c_2 n$
$T(n/2) = 2T(n/4) + c_2 n/2$

$T(n) = 4T(n/4) + 2\ c_2 n$
$T(n) = 8T(n/8) + 3\ c_2 n$

$T(n) = 2^i T(n/2^i) + ic_2 n$

Assuming $n = 2^k$, expansion halts when we get $T(1)$ on right side; this happens when $i=k$ $T(n) = 2^k T(1) + kc_2 n$

Since $2^k=n$, we know $k = \lg n$; since $T(1) = c_1$, we get

$T(n) = c_1 n + c_2\ n \lg n$;

thus an upper bound for $T_{mergeSort}(n)$ is $O(n \lg n)$

Quicksort Algorithm

Given an array of *n* elements (e.g., integers):

- If array only contains one element, return
- Else
  - pick one element to use as *pivot.*
  - Partition elements into two sub-arrays:
    - Elements less than or equal to pivot
    - Elements greater than pivot
  - Quicksort two sub-arrays
  - Return results

Example

We are given array of n integers to sort:

40	20	10	80	60	50	7	30	100

Pick Pivot Element

There are a number of ways to pick the pivot element. In this example, we will use the first element in the array:

40	20	10	80	60	50	7	30	100

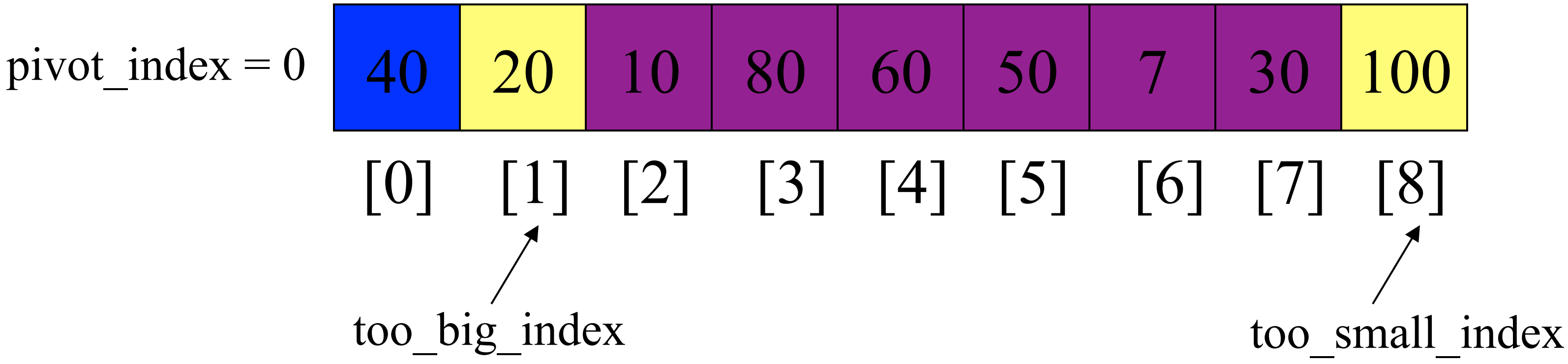
Partitioning Array

Given a pivot, partition the elements of the array such that the resulting array consists of:

1. One sub-array that contains elements >= pivot
2. Another sub-array that contains elements < pivot

The sub-arrays are stored in the original data array.

Partitioning loops through, swapping elements below/above pivot.


pivot_index = 0
40
20
10
80
60
50
7
30
100
[0]
[1]
[2]
[3]
[4]
[5]
[6]
[7]
[8]
too_big_index
too_small_index

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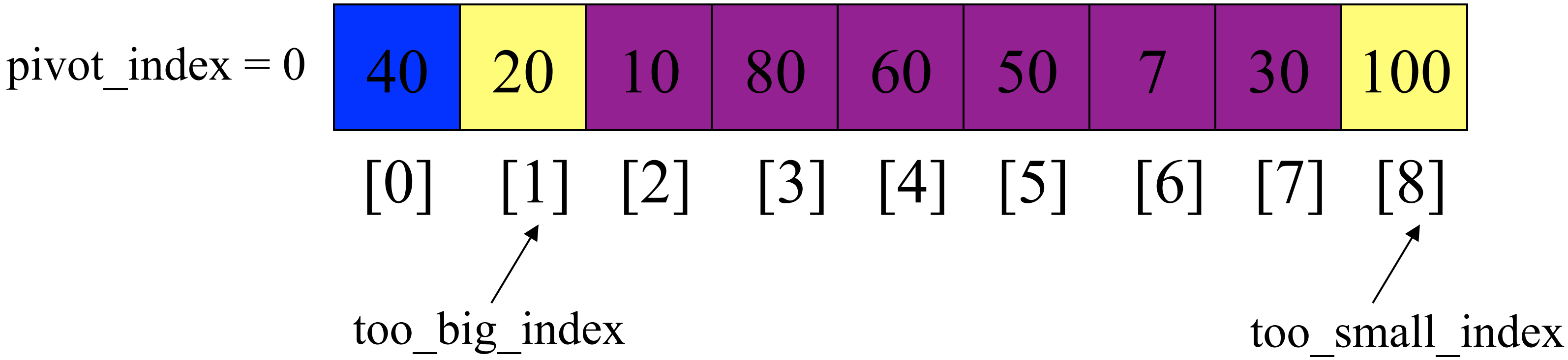

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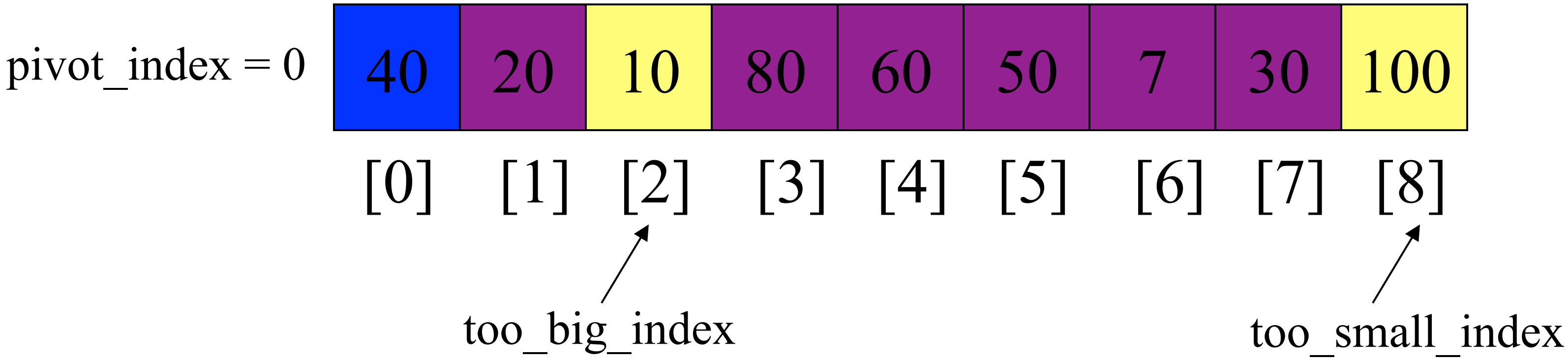

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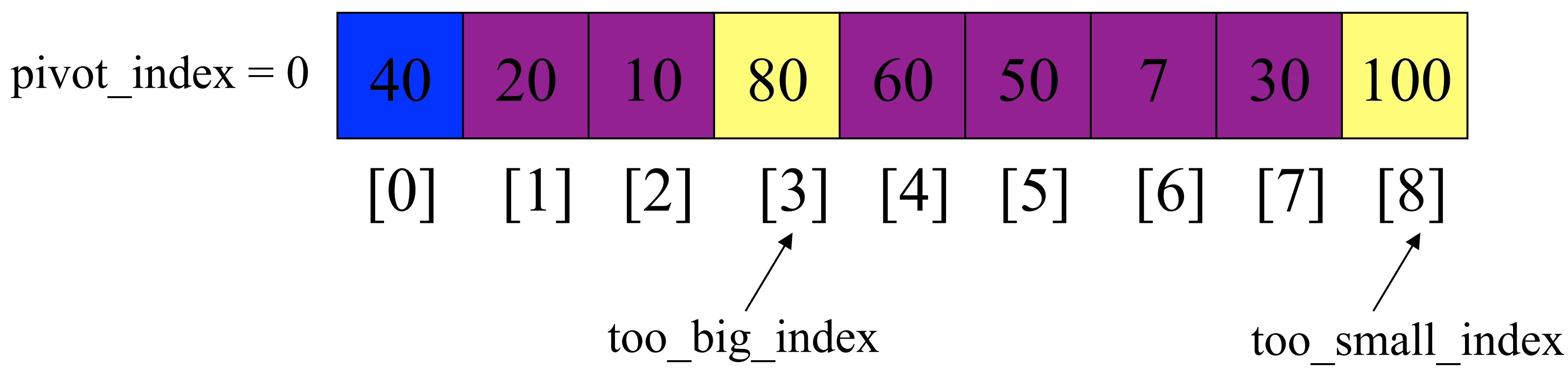

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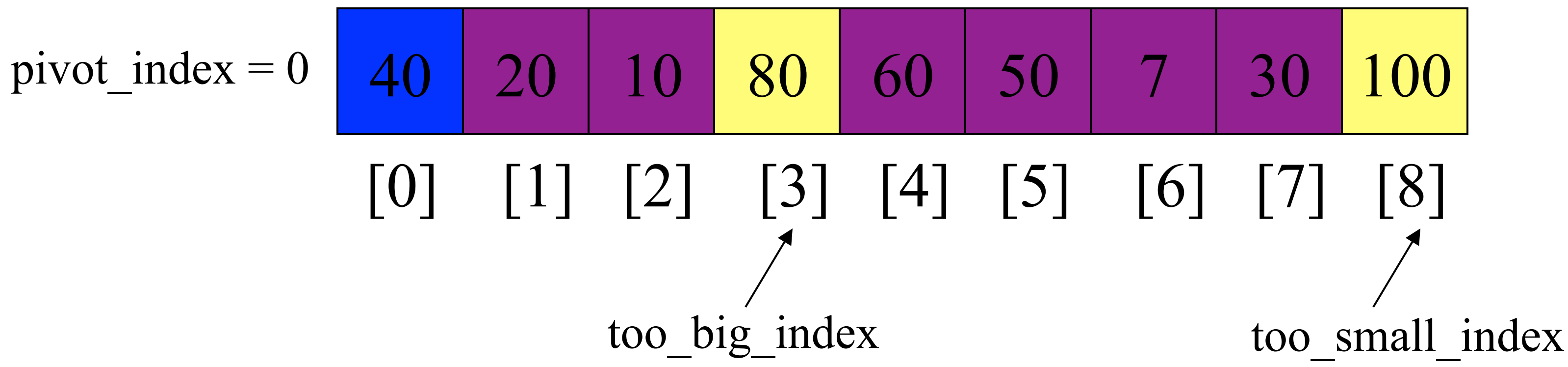

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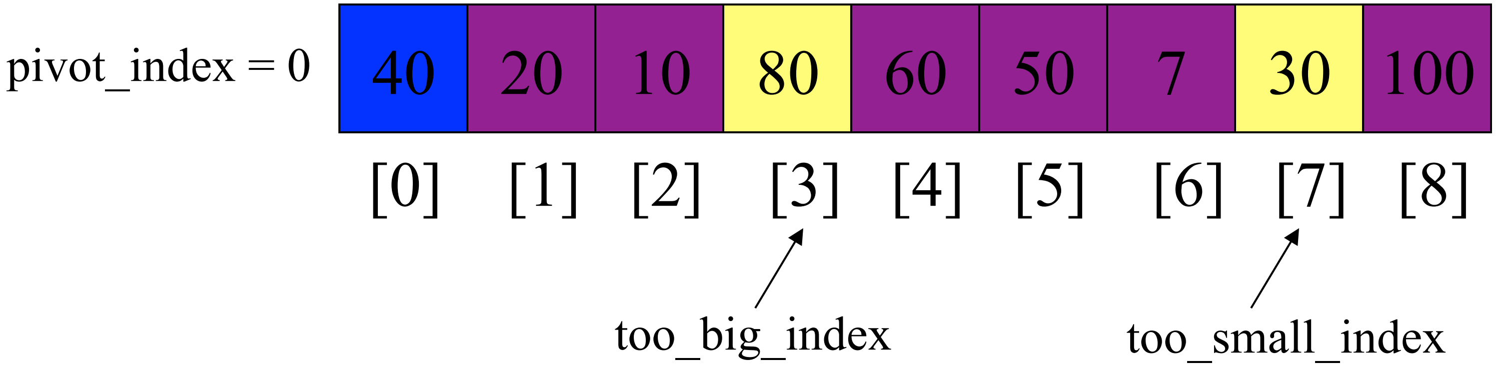

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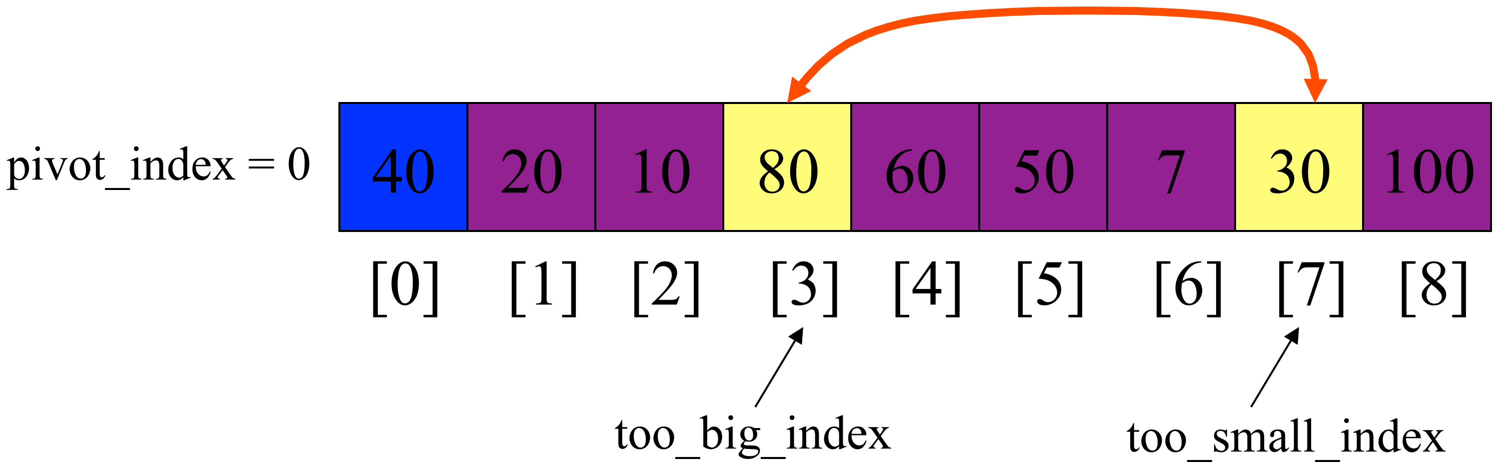

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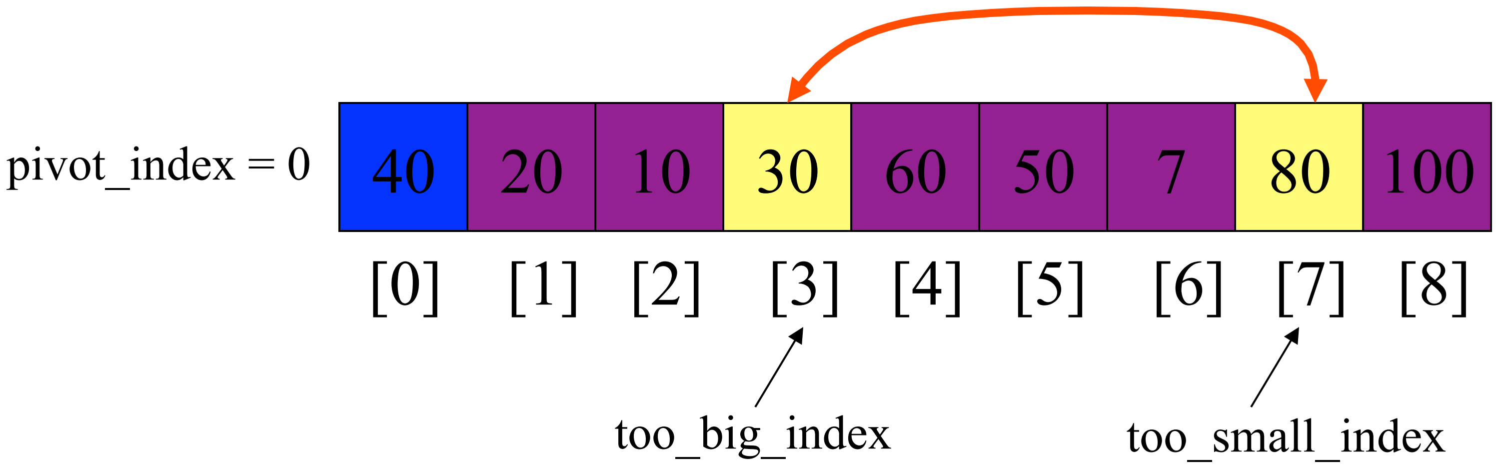

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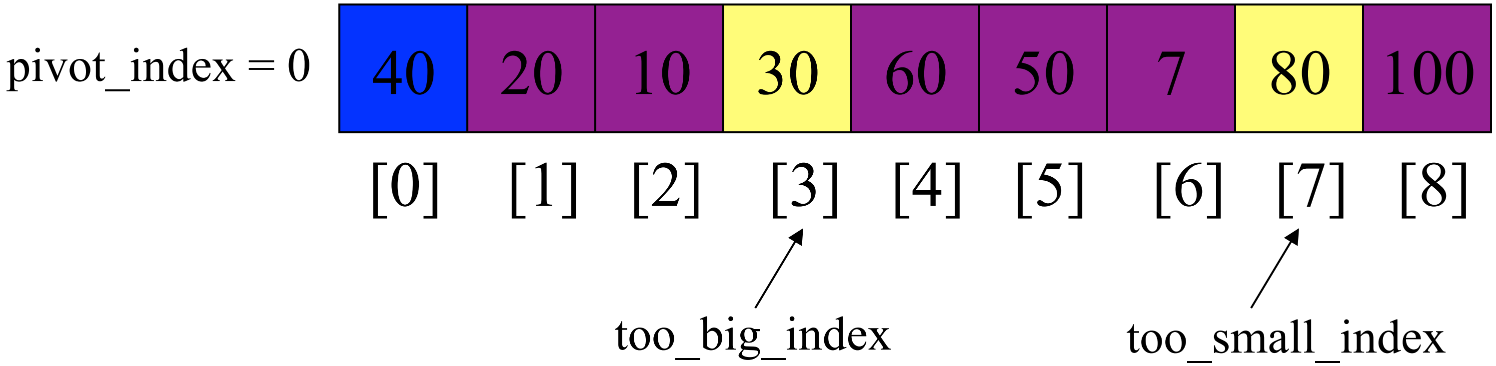

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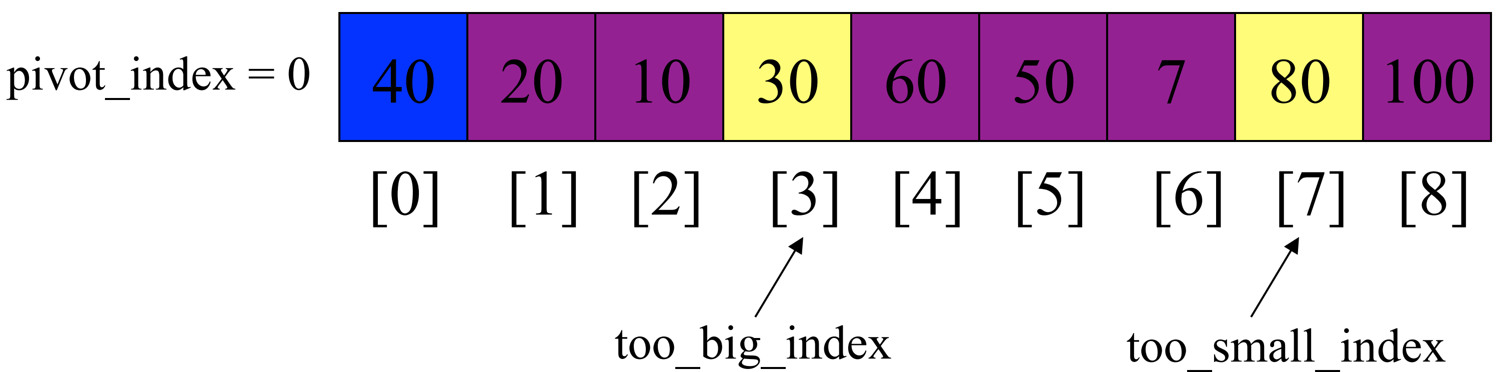

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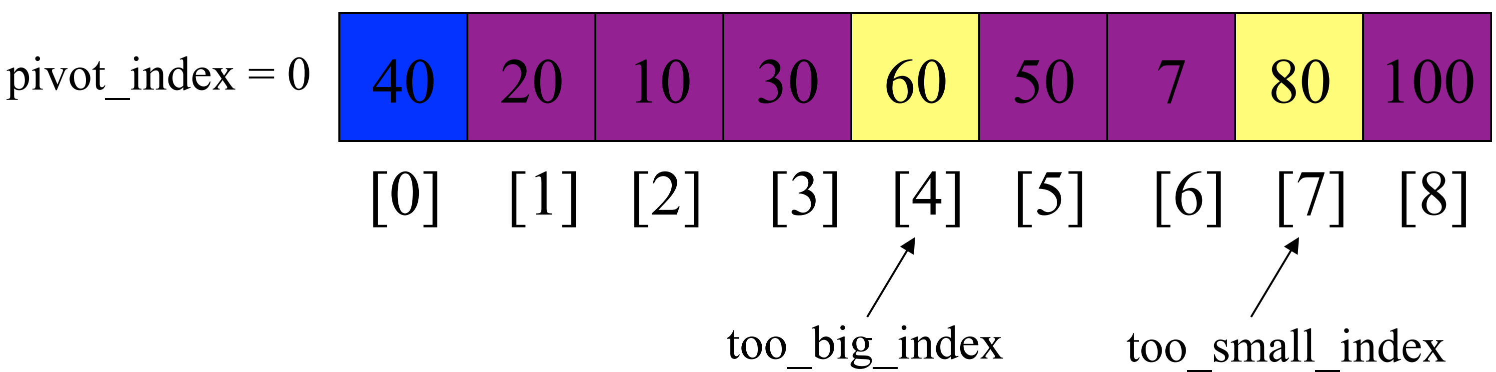

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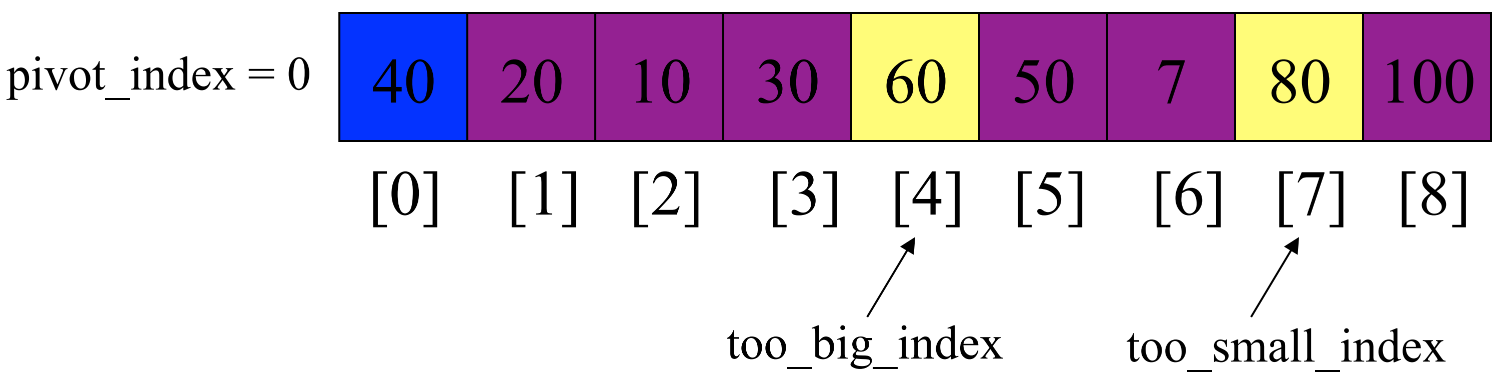

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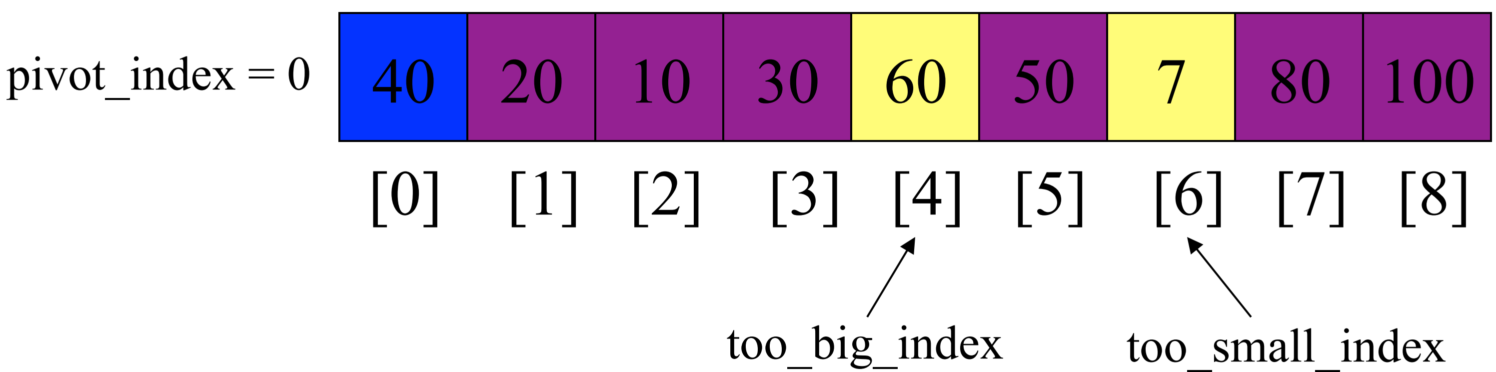

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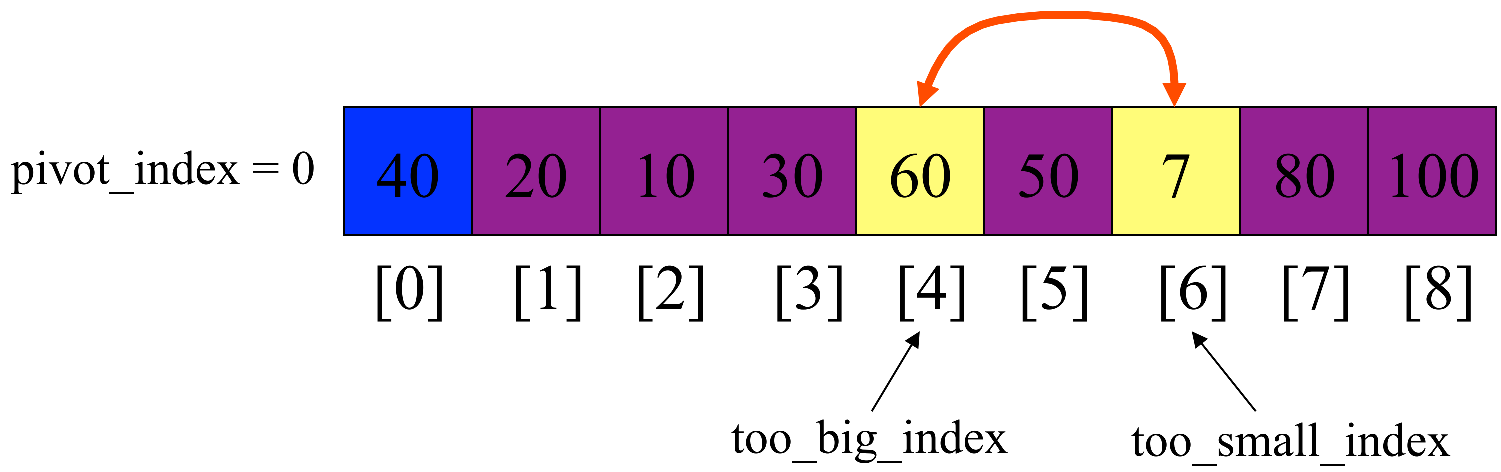

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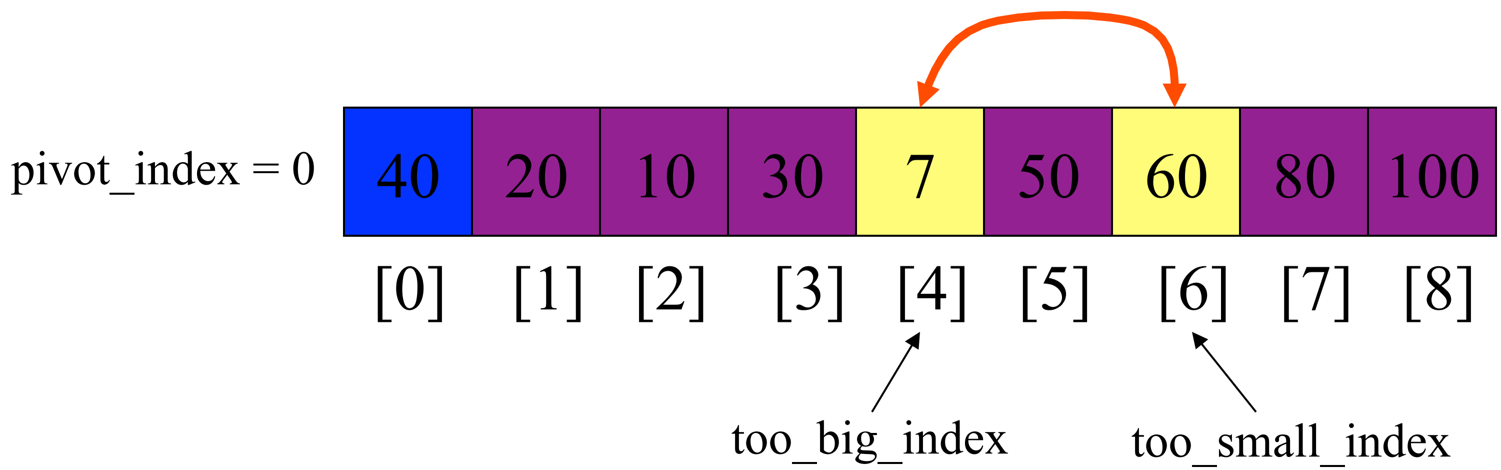

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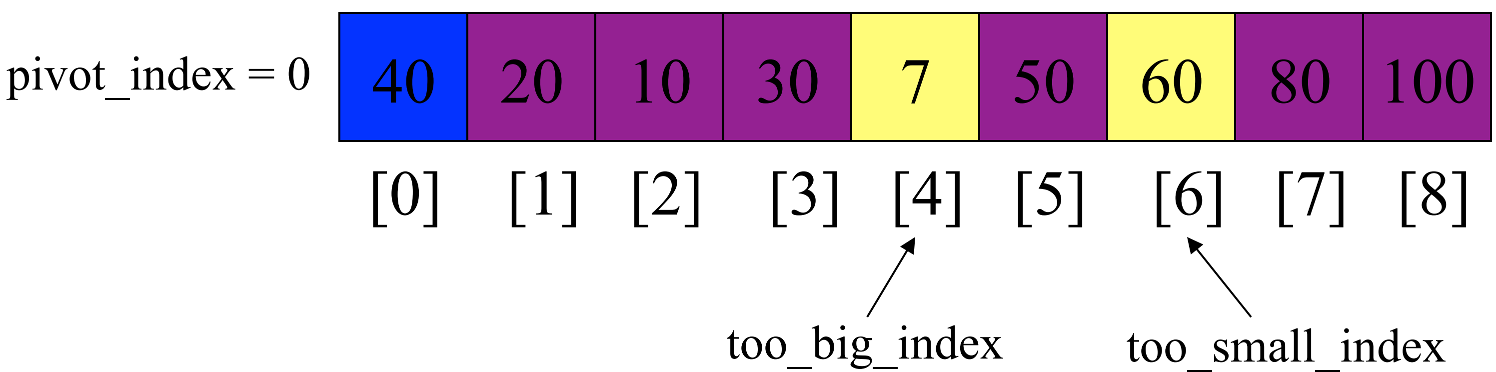

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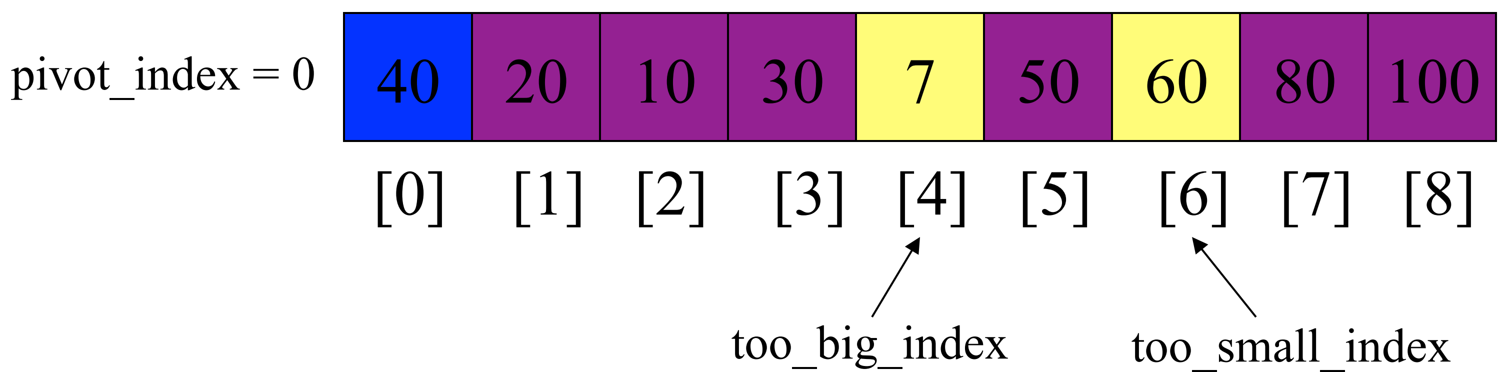

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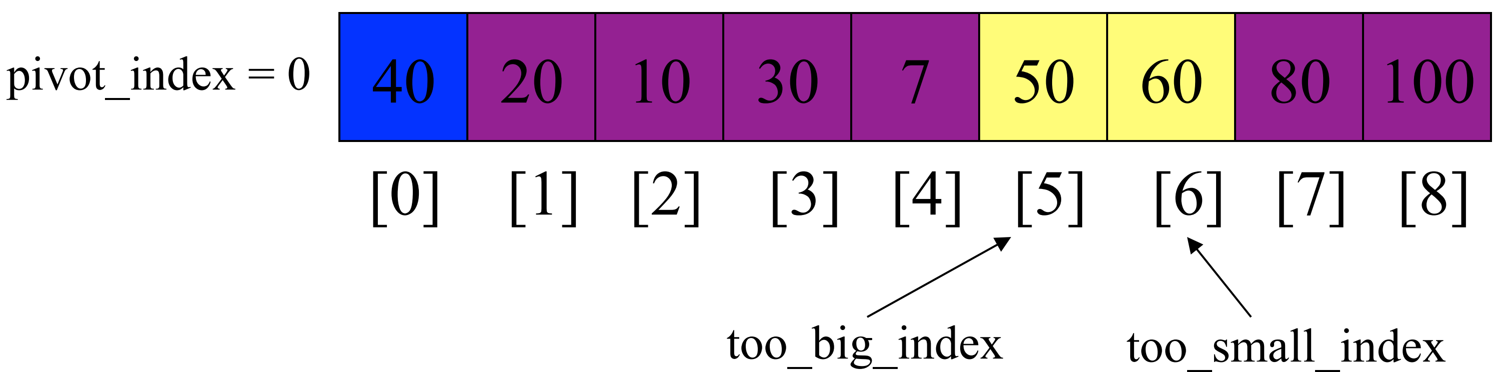

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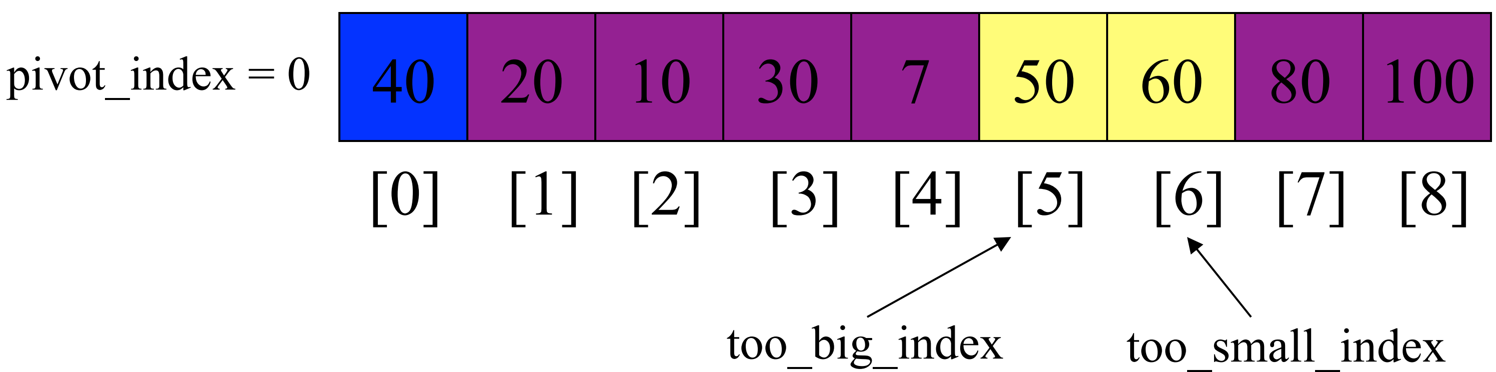

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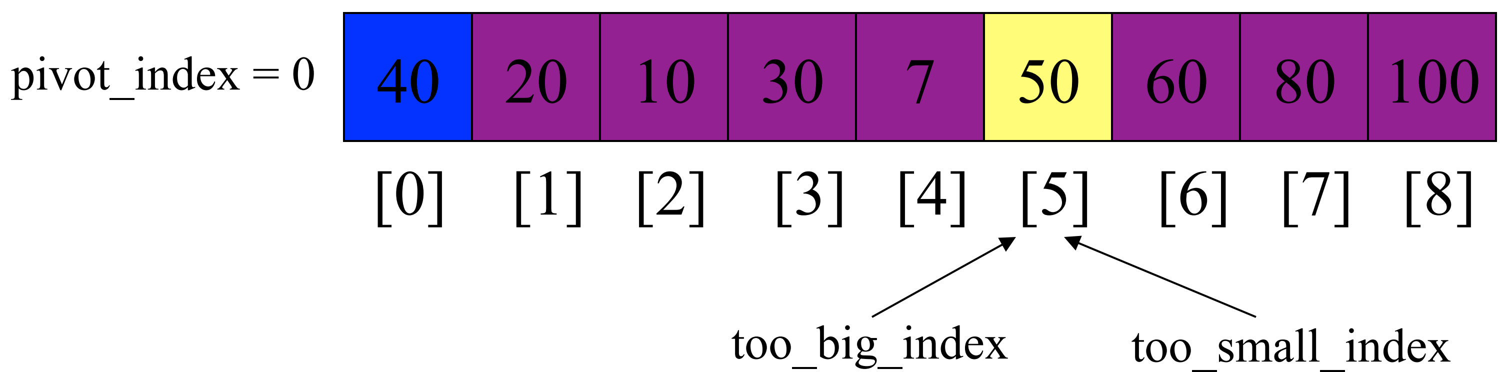

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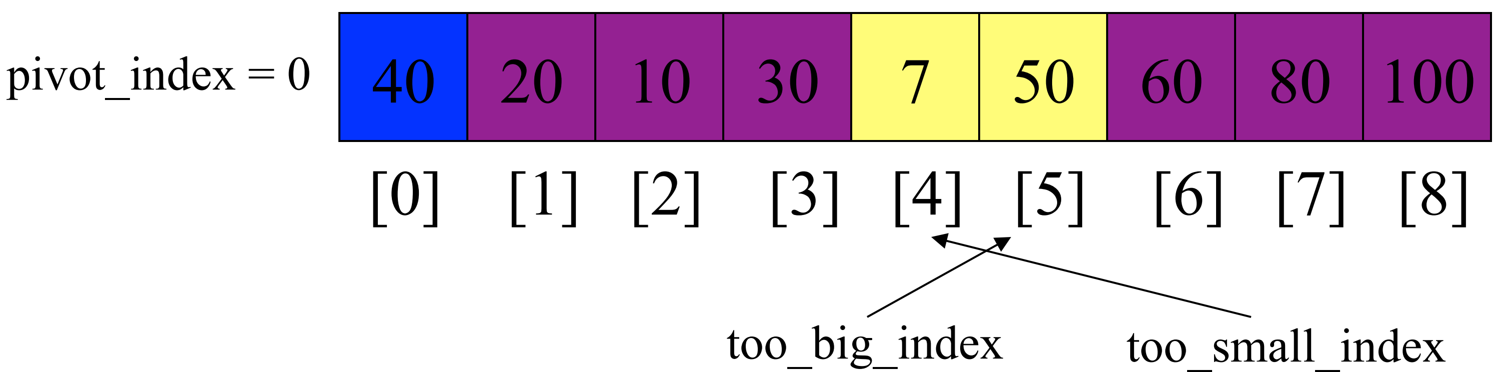

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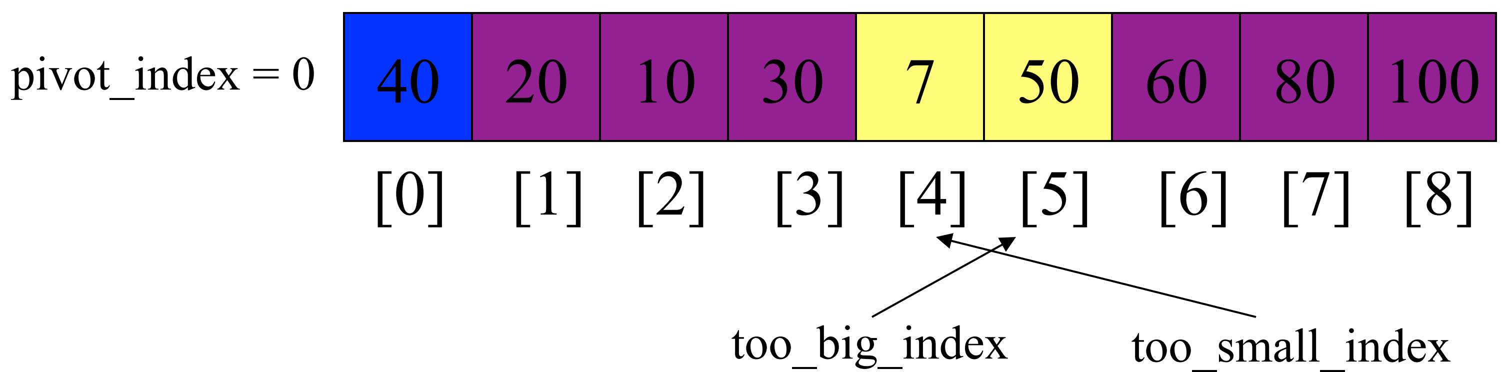

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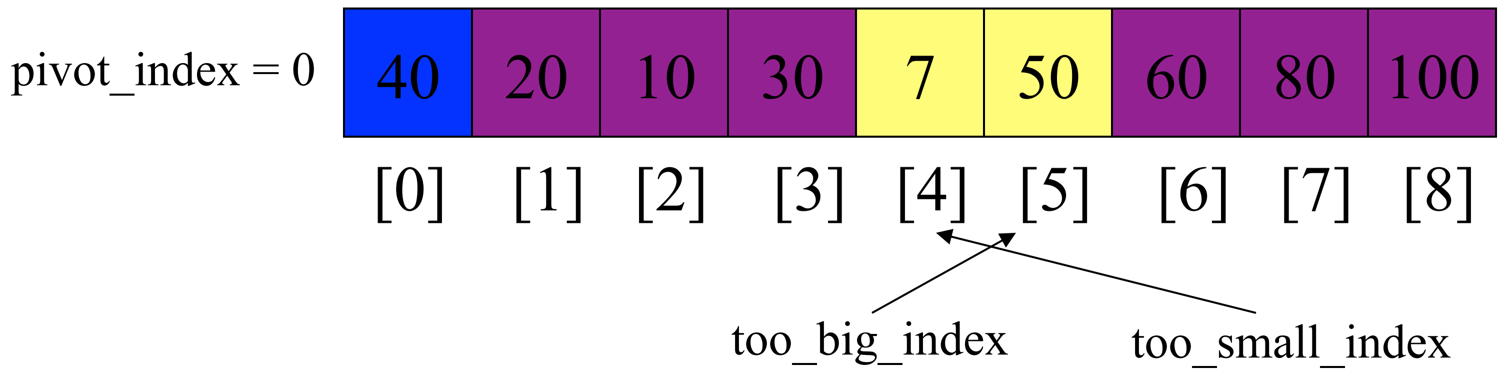

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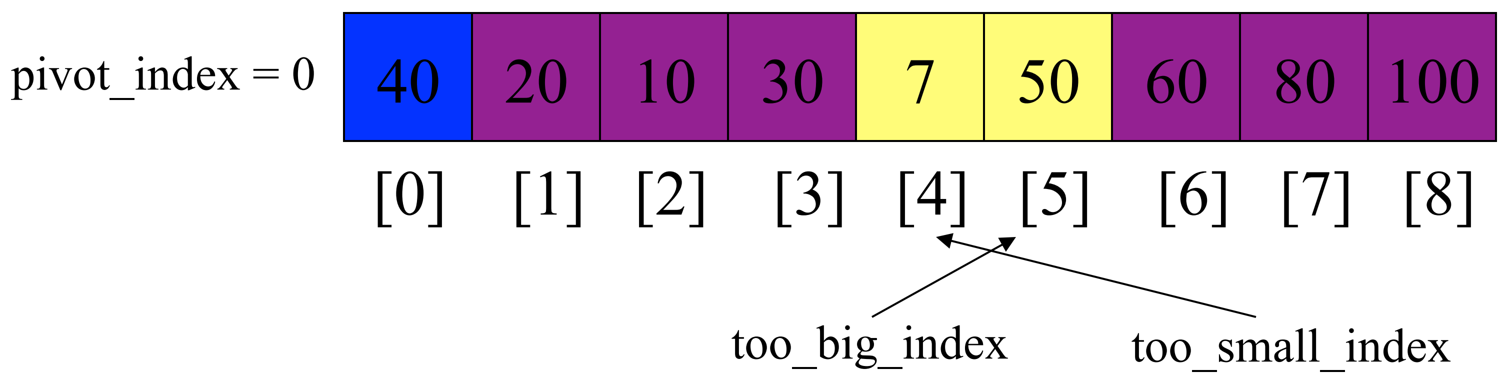

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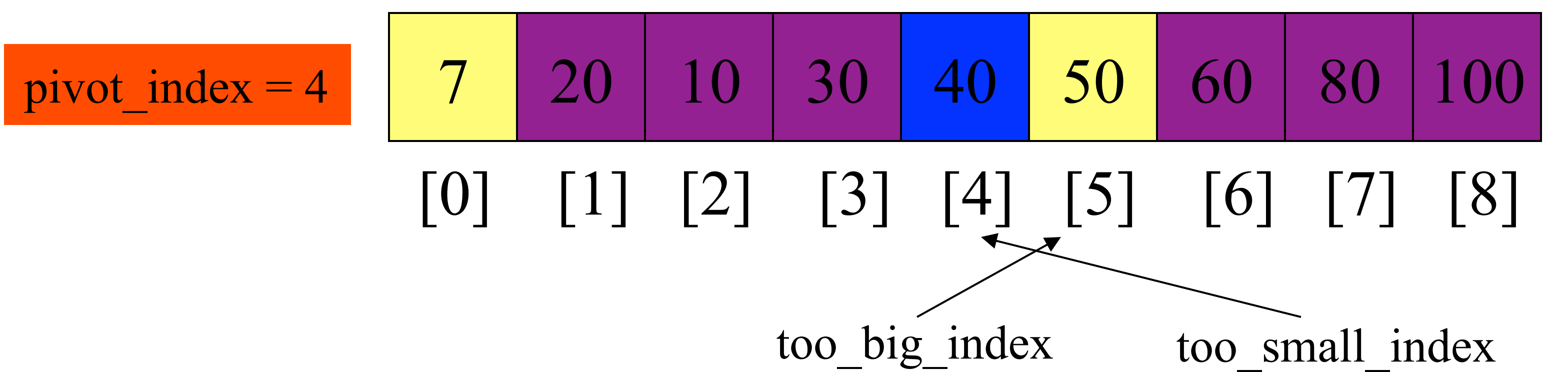
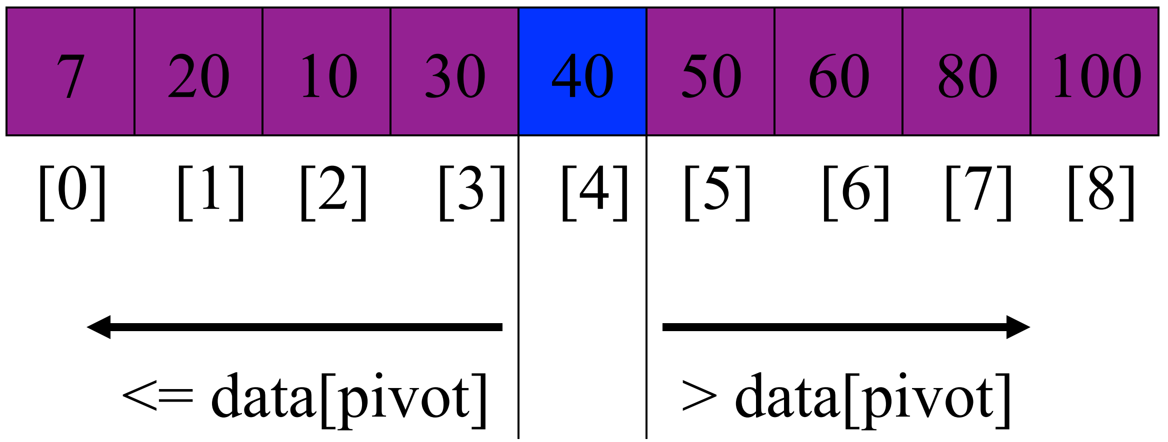
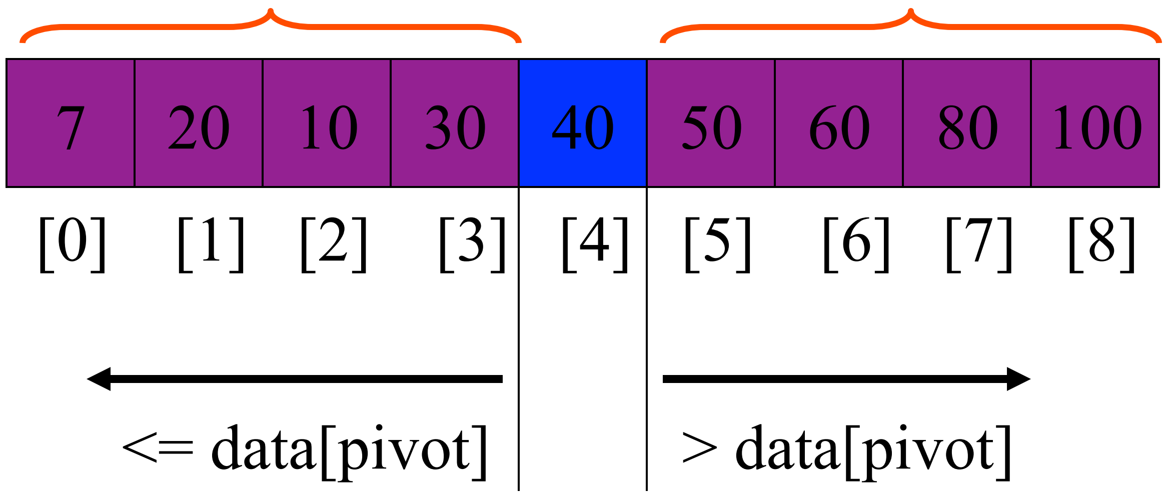


Partition Result



Recursion: Quicksort Sub-arrays



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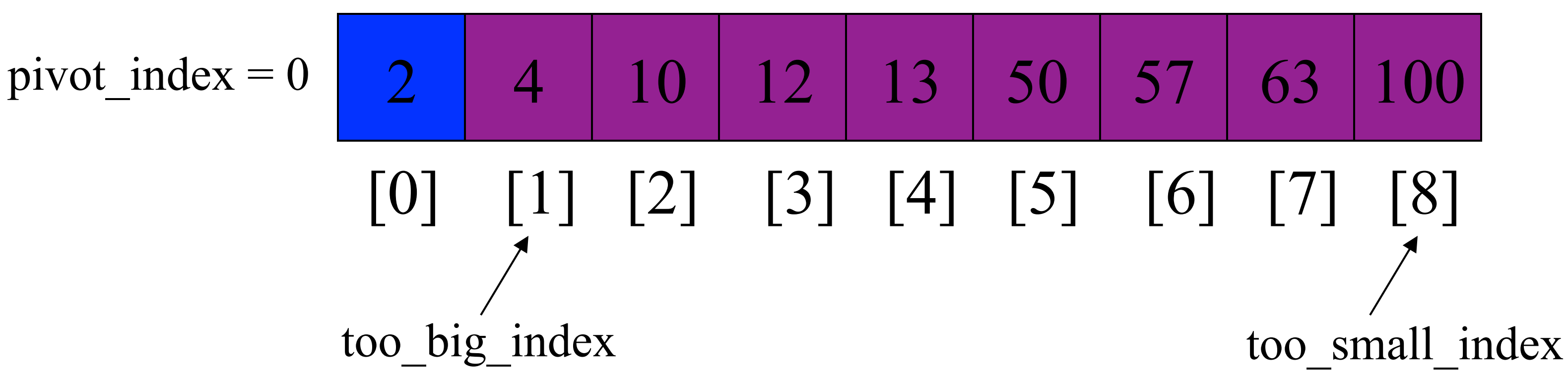
- Worst case running time?

Quicksort: Worst Case

- Assume first element is chosen as pivot.
- Assume we get array that is already in order:



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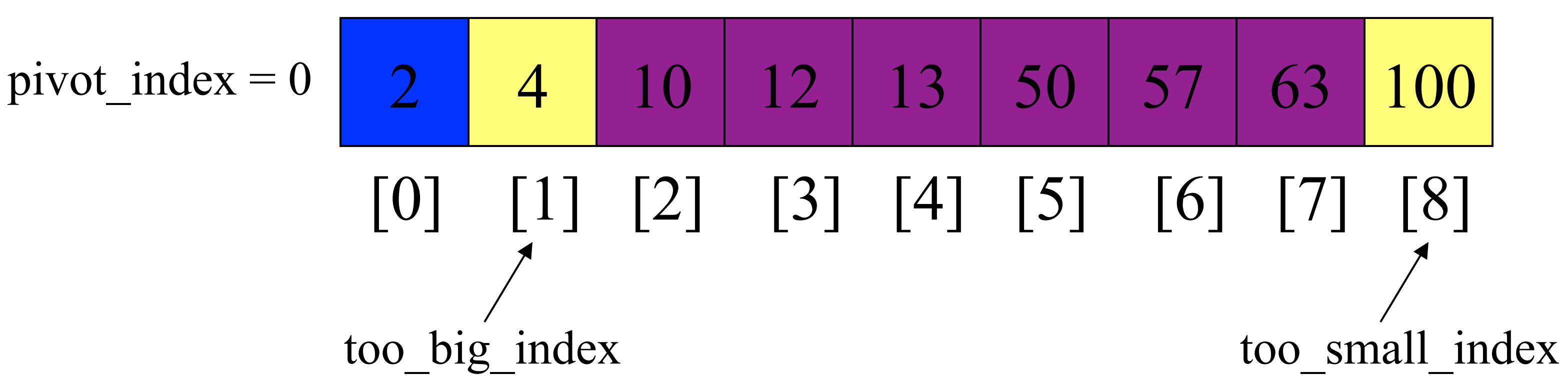

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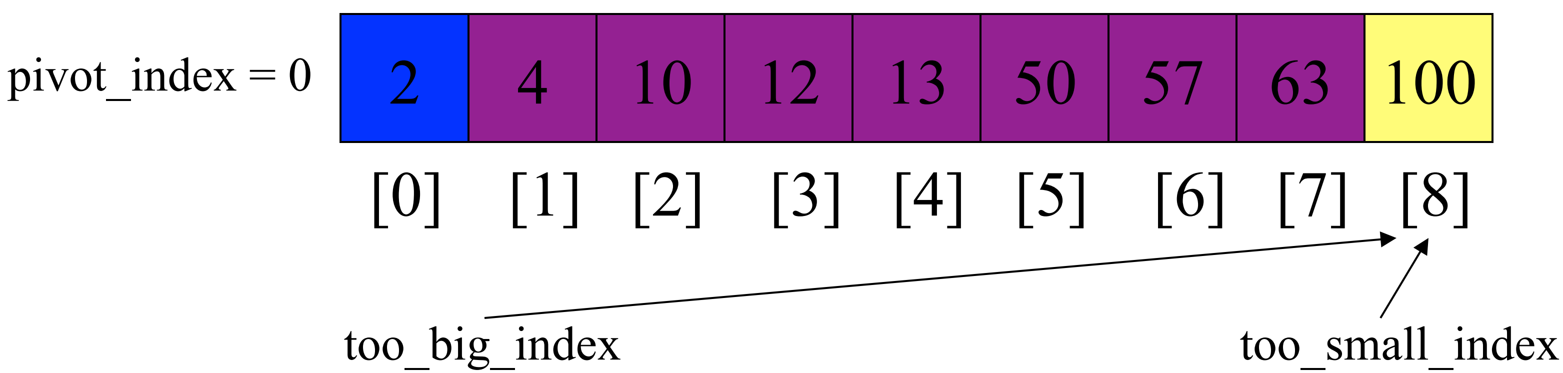

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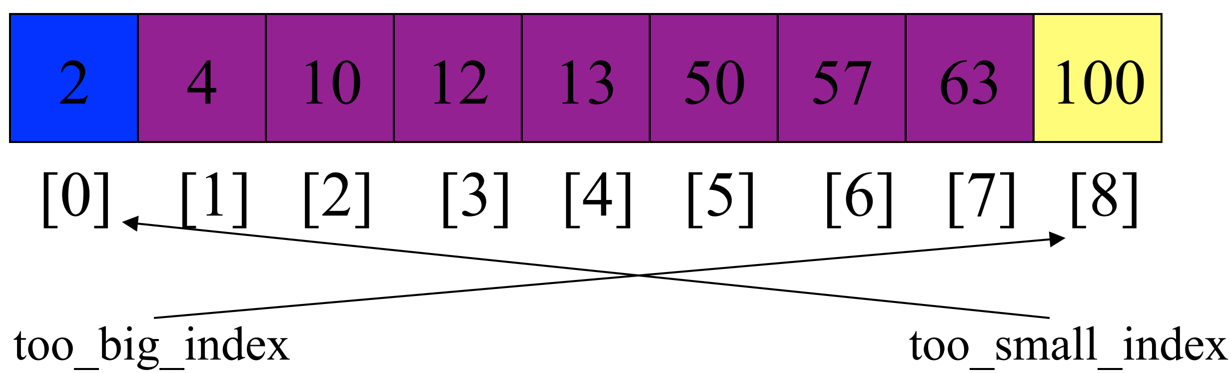

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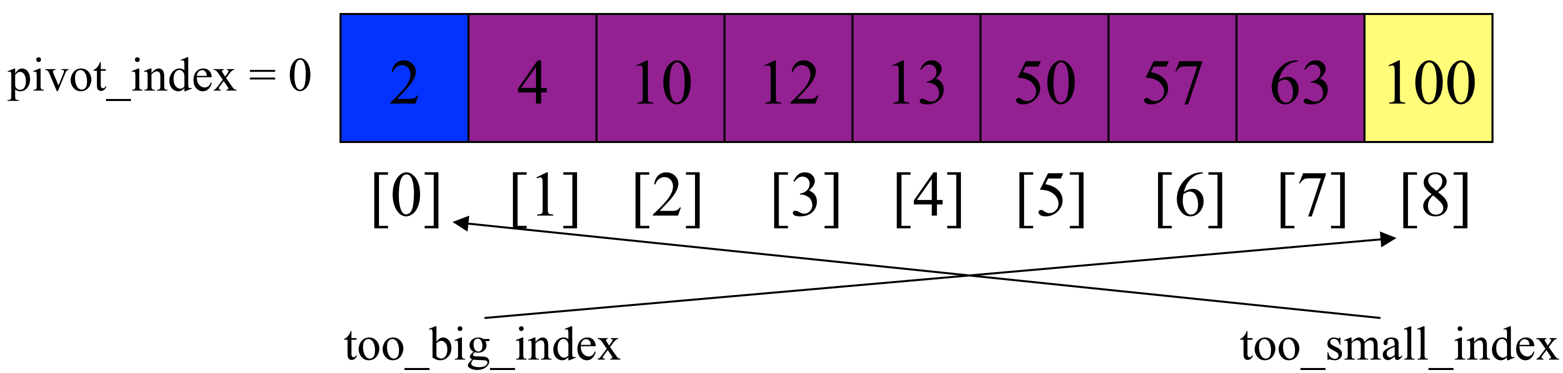

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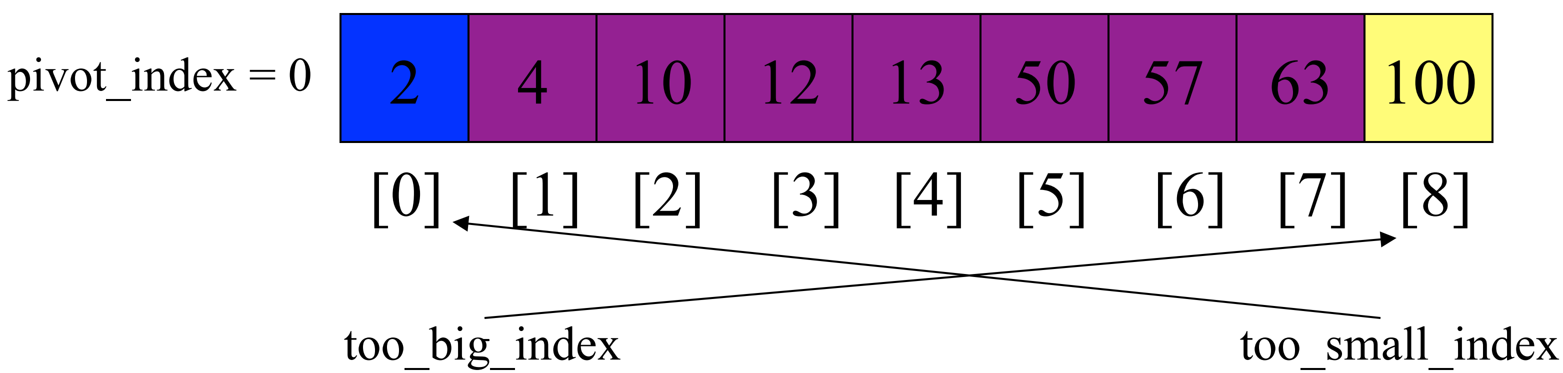

1. While data[too_big_index] <= data[pivot]
   ++too_big_index
2. While data[too_small_index] > data[pivot]
   --too_small_index
3. If too_big_index < too_small_index
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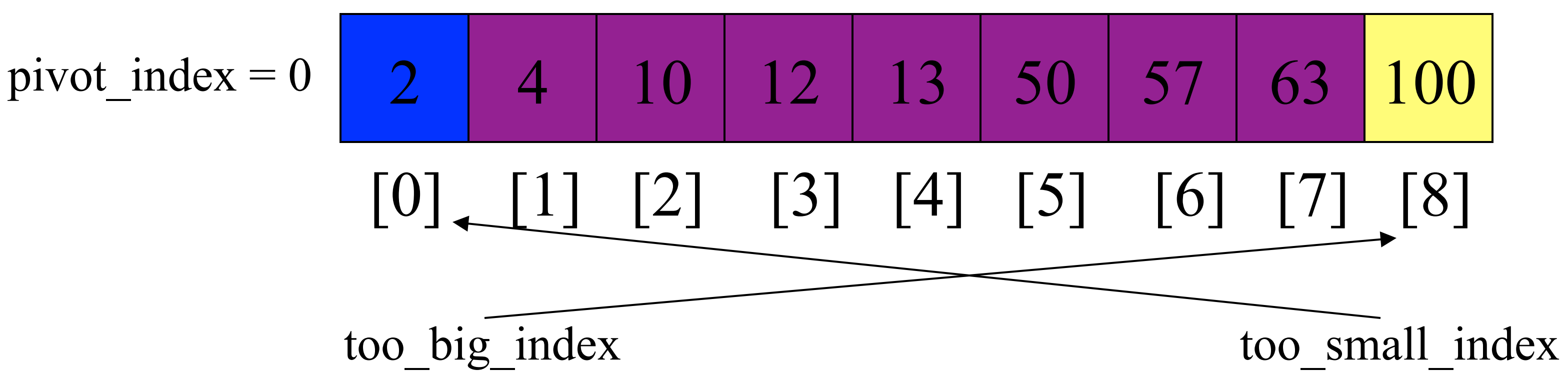

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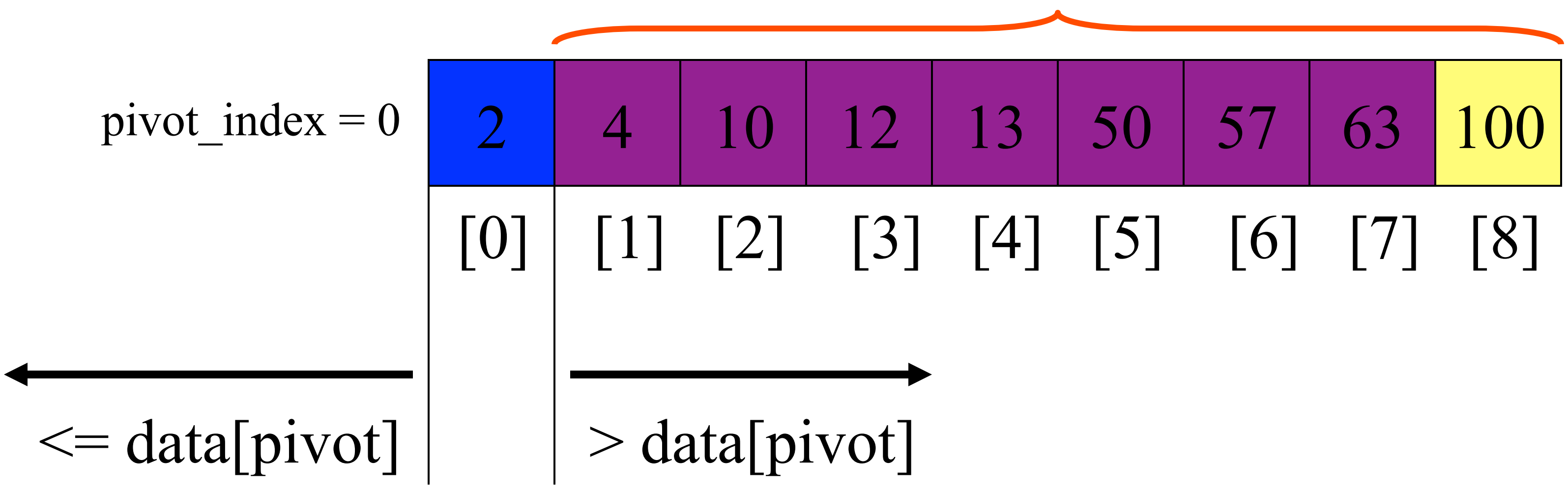

Quicksort Analysis

- Assume that keys are random, uniformly distributed.
- Best case running time: O(n lg n)
- Worst case running time?
  - Recursion:
    1. Partition splits array in two sub-arrays:
       - one sub-array of size 0
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- What can we do to avoid worst case?

Improved Pivot Selection

Pick median value of three elements from data array:

data[0], data[n/2], and data[n-1].

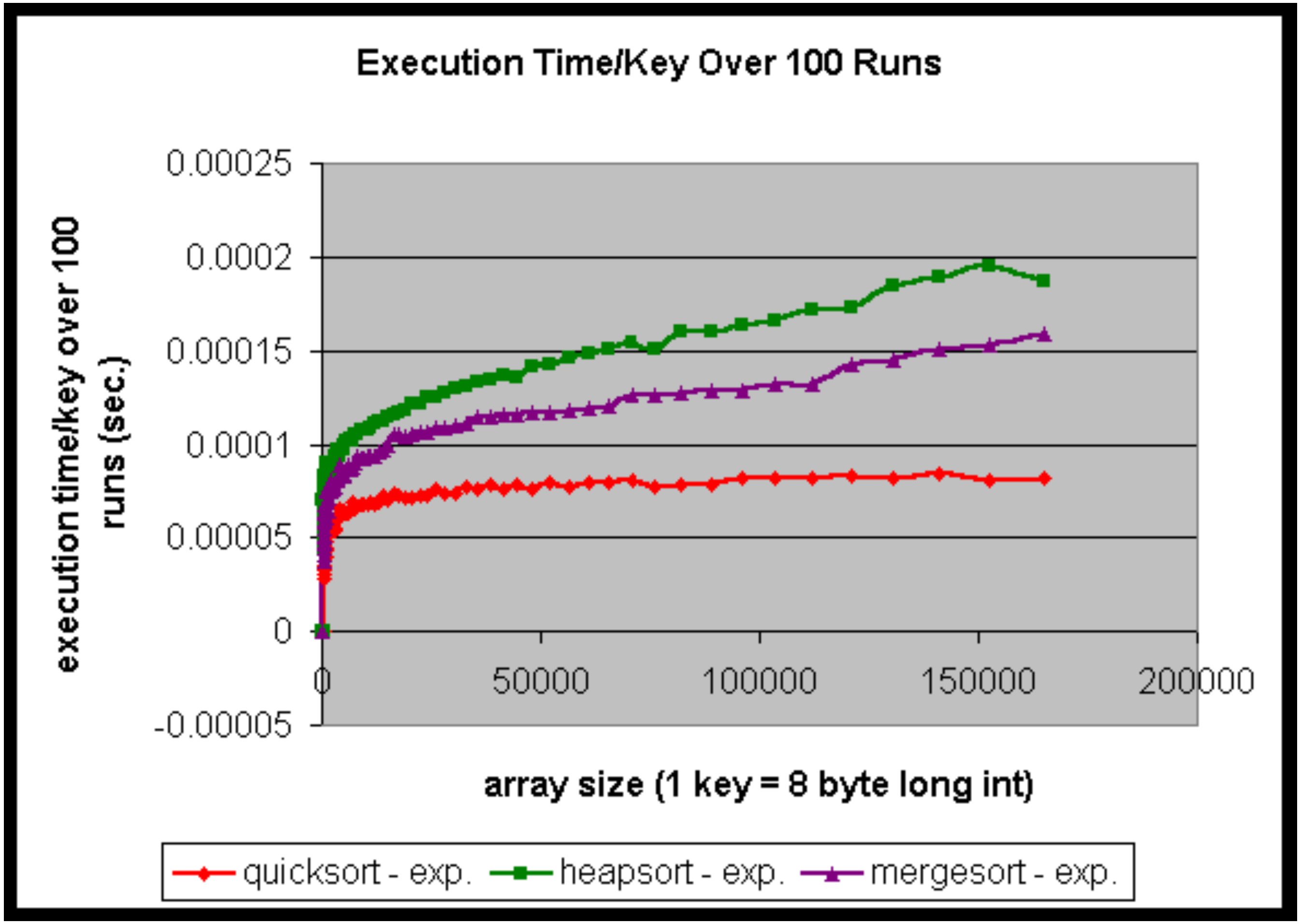
Use this median value as pivot.

Randomize array initially

Improving Performance of Quicksort

- Improved selection of pivot.
- For sub-arrays of size 3 or less, apply brute force search:
  - Sub-array of size 1: trivial
  - Sub-array of size 2:
    - if(data[first] > data[second]) swap them
  - Sub-array of size 3: left as an exercise.

Empirical Comparison



Sorting Algorithm Animations

SHARE

Problem Size: 20 · 30 · 40 · 50 Magnification: 1x · 2x · 3x

Algorithm: Insertion · Selection · Bubble · Shell · Merge · Heap · Quick · Quick3

Initial Condition: Random · Nearly Sorted · Reversed · Few Unique

http://www.sorting-algorithms.com/