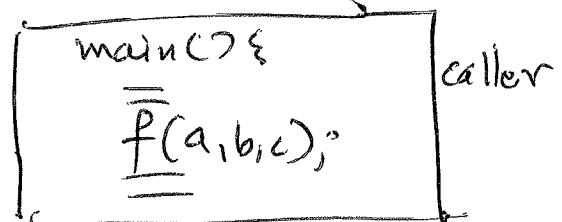


Subroutines

	<u>in</u>	<u>in out</u>	<u>out</u>
void	<u>f</u> (<u>int</u> [<u>3</u>] <u>x</u> , <u>int</u> [<u>3</u>] <u>y</u> , <u>int</u> [<u>3</u>] <u>z</u>) {		
	for(<u>int</u> <u>i</u> = <u>0</u> ; <u>i</u> < <u>x</u> .length; <u>i</u> ++) {		
		 <u>y</u>[<u>i</u>] = <u>x</u>[<u>i</u>] + <u>y</u>[<u>i</u>];	
		 <u>z</u>[<u>i</u>] = <u>y</u>[<u>i</u>];	
		}	
		}	
		// f()	



On invocation

- save the CPU state
- Set up referencing environment and stack frame.
- Transfer parameters
- Execute function code

On Return

- Transfer Parameters (if needed) and return value(s)
- Restore CPU state pop stack frame
- Continue caller execution

Role of Parameters

- receive some data as input to subroutine
- return some results in parameters
- both

return value \downarrow in

e.g. $y = f(x);$ $\underline{\text{in}}$ $\underline{\text{input}}$
 $\text{swap}(a, b);$ $\underline{\text{input}}$

Common Parameter Passing Schemes

1. Pass by value
2. Pass by result
3. Pass by value-result
4. Pass by reference
5. Pass by name.

BIG QUESTION
which parameter
passing scheme is used
by the PLs I know??

1. Pass-by-value in

Invocation

- allocate space for formal parameters on stack frame
- Evaluate actual parameters
- Copy value of actual parameters into formal parameters
- Execute function

Return

- Transfer return value (if any)
- Pop stack frame.

2. Pass-by-result out

Invocation

- Allocate space for formal parameters on stack frame
- Execute function

Return

- Copy values of formal parameters into actual parameters
- Transfer return value (if any)
- Pop stack frame.

3. Pass-by-value-result in-out

Invocation

- Same as pass by value

Return

- Same as pass by result.

4. Pass by reference

(in-out)

Invocation

- Copy the reference of actual parameters into formal parameters.
- Execute function.

Return

- Pop stack frame.

5. Pass by Name

Invocation

- Textually substitute the names (or expressions) for formal parameters.
- Execute the function.

Return

- Pop the stack frame.

Example

C/Java

```
void swap (int a, int b) {
    int temp = a;
    a = b;
    b = temp;
}
```

//swap()

main()

```
int x = 5, y = 7;
swap(x, y);
```

1. pass-by-value

2. pass by result

3. pass by value-result

4. pass by reference

5. pass by name.

Issues

① Aliasing can lead to issues.

e.g. ~~comp~~

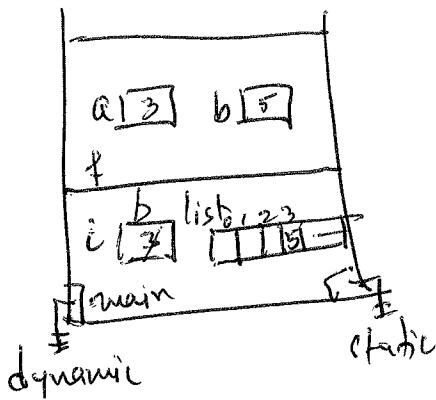
int i = 3;

```
void f(int a, int b) {  
    |  
    i = b;  
}| //f()
```

```
int list[10];  
list[i] = 5;  
f(i, list[i]);
```

Value-Result

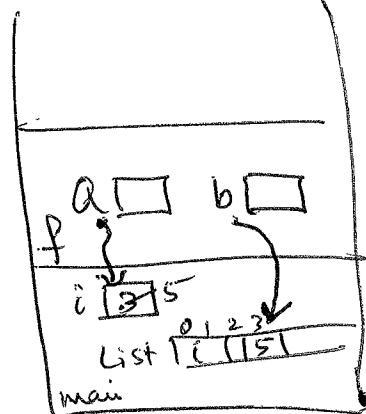
~~Value~~



on return a=3 is copied into i,
so change in i=5 is gone.

Reference

a + b are aliases!

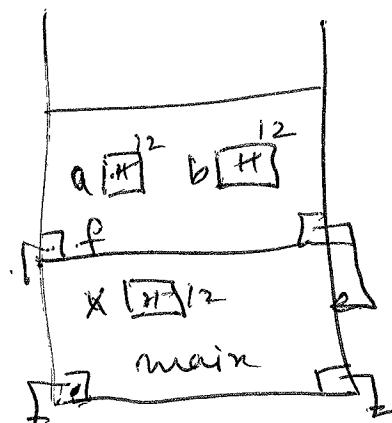


e.g.

```
void f (int a, int b) {  
    a++;  
    b++;  
}
```

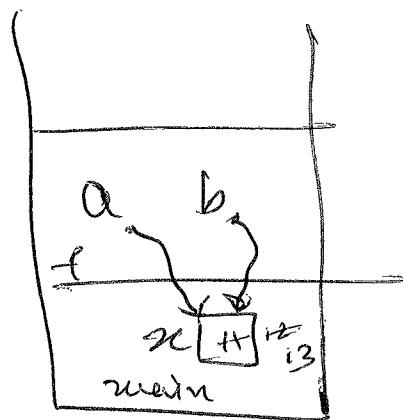
```
int x = 11;  
f(x, x);  
x = ??
```

Value - Result



$$x = 12$$

Reference



$$x = 13$$

~~in swap~~
by
swap (i, a[3])
call by name.

Parameter Passing Schemes in PLs

C : pass-by-value,

can do pass-by-reference using pointers.

Arrays are "by reference" since $a[0] \equiv @a[0]$

C++ : pass by value, but also has reference types

C# : pass by value, also has pass-by reference

```
void swap (ref int a, ref int b) {  
    3
```

call: swap (ref x, ref y);

Python: - pass-by-value assignment

- a combination of pass-by-value
+ pass by reference

- immutable objects (numbers, strings, tuples)
are pass by value

- mutable objects (lists, dictionaries),
ref to object is assigned to formal parameter
i.e. passed by reference