

Type Checking: When are two types identical

- $v_1 : t_1$                                $\text{def } f(v_1 : t_2) \{$   
 $v_2 : t_2$                               }  
1.  $v_1 = v_2 ?$   
2.  $\dots = v_1 + v_2 - \dots ?$   
3.  $f(v_1) ?$

- Strongly + Weakly typed language
- Static vs dynamic typing
- Type Equivalence — Name Equivalence
  - Structural Equivalence
- Type Compatibility - a more practical approach  
Two types are compatible if
  - 1. They are equivalent
  - 2. One is a subtype of another
  - 3. Both are arrays w/ same type elements

JAVA Type Compatibility

- 1. Identical types
- 2. Subtypes + Supertype
- 3. Interfaces
- 4. Type Casting

## Arrays

In some languages the index range of an array is part of the type

e.g. PASCAL

type A = array [1..10] of integer;

type B = array [0..9] of integer;

~~var~~ ~~A+B~~ var a: A;  
b: B;

are  $a + b$  equivalent?

how about

function max (x: ~~A~~ A)

≡

max (b); ??

only if

function max (x: array [1..10] of integer)

≡

still cannot do  $\max(b)$ !

So, ~~Modula-2~~ used open arrays,

FUNCTION max (x: array of integer)

now we can call  $\max(a);$

$\max(b);$

but indexing in max will be  $x[0..]$ .

What about C + Java?

```
int max (int arr[]) {  
    |  
    3
```

~~arr we don't call~~

we can call `max()` with any int array of any size.

BUT, how do we know how many elements in `arr`?

e.g.

```
int [10] a;  
int [10] b;
```

— `= max(a);`  
— `= max(b);`

how does `max` know sizes of `a[10]` & `b[10]`.

we have to do

```
int max (int arr[], int n) {  
    |  
    3
```

+ call  
— `= max (a, 10);`  
— `= max (b, 100);`

what about Java?

```
public static int max (int arr[]) {  
    } — arr.length is # elements  
    | in arr[]  
    | 3
```

No need to send `n`  
since Java uses the  
length attribute  
Python?? `len(a)`

## Type Conversion / Coercion

when two types may not be structurally equivalent  
but there is a way to convert one from another.

— can be done implicitly or explicitly

e.g. int x;

double y;

x=3;

y = y + x;

or y = x;  $\leftarrow$  x is implicitly converted to double.

however, doing x=y;

does not make sense - why?

PLs provide either casts

i.e. x = (int)y;

x = round(y);

~~or~~ ~~round~~

also, int n = (int)(Math.random() \* 10);

$[0.0 - 1.0] \downarrow$   
 $10.0$

and in Java    System.out.println("A=" + 3);

"A=" 3  $\rightarrow$  printed.

## Python - Duck Typing !!

"If it talks like a duck, walks like a duck  
then it is a duck!"

- No explicit type compatibility rules
- Not ~~type~~ a strongly typed language

Python is dynamically typed.

### Rule

If an object can perform the actions that are expected in a context, then it is compatible.

int is compatible with float

errors happen at run time.

e.g. ~~a = 5~~  
list(a) ← a not a list!

## Type Inference

- automatic detection of types of an expression

Javascript

$x := 5.3$

inferred as float.

Haskell, SML, Javascript, etc.

Can get very complicated [where the PL theory research actually is!]

e.g. ~~def~~ add-one(x) {  
    int result  
    result = ~~x~~ x + 1  
    return result.

Inference

add-one : int  $\rightarrow$  int

Haskell

map(double, a)

double :: int  $\rightarrow$  int

a : [int]

~~map~~ :: map :: (int  $\rightarrow$  int)  $\rightarrow$  [int]  $\rightarrow$  [int]

More abstractly

map(f, l)

map f [ ] = [ ]

map f (first: rest) = f first  $\circ$  map f rest

$\therefore$  map :: ( $\alpha \rightarrow \beta$ )  $\rightarrow$  [ $\alpha$ ]  $\rightarrow$  [ $\beta$ ]

~~Abstraction~~:

## Control Abstraction - Subroutines

Control Abstraction vs Data Abstraction

### Subroutines

- Functions : return value(s)
- Procedures : do not return any value

e.g. Java

```
double area (double radius) {  
    |   return Math.PI * radius * radius;  
} //area()
```

Function

```
void drawCircle (double x, double y, double radius) {  
    |   ==  
    |   == return;  
    |   == return;  
} //drawCircle()
```

## Vocabulary + Design Considerations

- Caller, function call, invocation - `double a = area(r);`
- signature - `double area(double)`
- Arguments/actual parameters, formal parameters
  - $a = \text{area}(r);$   $\begin{matrix} \leftarrow & \text{actual} \\ \text{parameter} & \end{matrix}$
  - show formal parameter in example code.
- return  $\dots \rightarrow \text{return } \langle \text{expression} \rangle;$ 
  - primitive type only?
  - can be an aggregate type?
  - can be a function?
  - can return 0 or 1 value?
  - can return multiple values?

## Parameter Association

• positional  $\text{drawCircle}(x1, y1, r1);$

• named / keyword parameters

$\text{drawCircle}(x=x1, y=y1, radius=r1);$

$\text{drawCircle}(radius=r1, y=y1, x=x1);$

• Mixed association

$\text{drawCircle}(x1, y1, radius=r1);$

• default parameters

def convert(n, base=2):

binary = convert(n)

octal = convert(n, 8)

hexadecimal = convert(n, base=16)

• variable number of arguments

printf(<string>, v1, v2, ...);

C

```

int max (int nargs, ...){  

    int max, a;  

    va-list args;  

    va-start (args, nargs);  

    max = va-arg (args, 1);  

    for (int i=2; i<=nargs; i++)  

        if ((a = va-arg (args, i)) > max)  

            max = a;  

    va-end (args);  

    return max;  

}
//main()  

max(5);  

max(5, 1) → 5  

max(5, 16, 7) → 16

```

Java

```

int max(int a, int ... ns){  

    int m = a;  

    for (int x : ns)  

        if (x > m)  

            m = x;  

    return m;
}

```

Python

```

def max(a, *argv):  

    m = a  

    for x in argv:  

        if x > m:  

            m = x  

    return m

```