

Type checking

$v_1 : t_1$

$v_2 : t_2$

def $f(v : t_2) \{$

|

}

Q.1. Can we assign:

$v_1 = v_2$?

Q.2. $* \dots = \dots \cdot v_1 + v_2 \dots$?

Q.3. $f(v_1)$

Answers to above questions depend on whether a PL is strongly typed or weakly typed + on the type compatibility rules of the PL

Strongly Typed PLs

- when they are strict about types.

→ i.e. we cannot do any of the above...

Weakly Typed PLs

when PLs are not strict. e.g. Python, JS

When to do type checking ?

Static Typing (statically typed PLs)

- when all type checking is done at compile time.

* e.g. Java, SML, C, C++

Dynamic Typing (dynamically typed PLs)

- when type-checking is done during run-time.

Type Equivalence

- defines when two types in a PL/program are equivalent.

e.g. - are float + int equivalent?

```
float a;  
int b;
```

a = b; ??

Two types of Type Equivalence

Name Equivalence

expressions

two types are equivalent if they have the same name.

e.g. above a = b is not allowed as float + int are different type names

Structural Equivalence

Two types are structurally equivalent if they are made up of the same parts.

e.g. C

```
typedef float celsius;  
typedef float fahr;
```

```
celsius c = 100.0;  
fahr f;
```

f = c; OK in C

since C uses structural equivalence.

Java

```
public class Celsius {  
    private float temp;  
}
```

} // class Celsius

```
public class Fahr {  
    private float temp;  
}
```

} // class Fahr.

```
Celsius c = new Celsius(100.0);  
Fahr f;
```

```
f = c; X not allowed  
since Java uses  
name equivalence.
```

In fact in Java, we can then write

In Celsius:

```
public float toFahr() {
```

```
    }  
}
```

In Fahr

```
public float toCelsius() {
```

```
    }  
}
```

use

```
f = c.toCelsius();  
c = f.toFahr();
```

Some PLs have both!

e.g. Ada

```
type celsius = real;  
type fahr    = real;
```

In Ada, we use structural equivalence for these types.
BUT, Ada also allows derived types

```
type celsius = new real;  
type fahr    = new real;
```

Now celsius + fahr are new + different types ~~are~~ even though they are structurally equivalent.

Another example w/ functions

C/Java

```
int max(int a, int b) {
```

3

can we do: ~~int~~ float x, y = ...
max(x, y)?

Need to define something more practical than name/structural equivalence — Type Compatibility Rules

e.g. Two types are compatible if:

1. The two types are equivalent
2. One is a subtype of the other
3. Both are arrays with the same types of elements.

Java type compatibility

1. Identical types are compatible
i.e. int and int
float and float

2. subtypes + supertypes

class Animal {

|

3

class zebra extends Animal { —

|

3

zebra is a subtype of the supertype Animal

i.e. we can do

Animal a = new zebra(); ✓

But not zebra z = new Animal(); ✗

3. Interfaces

any class that implements an interface is a
subtype of that interface

4. Type Casting

- explicitly convert a variable to another

e.g.

float f;

int i;

i = (int) f;

but cannot do

String s;

i = (int) s;

etc