

## Type checking

```
v1 : t1  
v2 : t2  
def f(v:t2){  
|  
}  
}
```

Q.1. Can we assign:

$$v1 = v2 ?$$

Q.2.  $* \dots = \dots . v1 + v2 \dots ?$

Q.3  $f(v1)$

Answers to above questions depend on whether a PL is strongly typed or weakly typed + on the type compatibility rules of the PL

### Strongly Typed PLs

- when they are strict about types.

↗ i.e. we cannot do any of the above...

### Weakly Typed PLs

when PLs are not strict. e.g. Python, JS

### When to do type checking?

#### Static Typing (statically typed PLs)

- when all type checking is done at compile time

e.g. Java, SML, C, C++

#### Dynamic Typing (dynamically typed PLs)

- when type checking is done during run-time

## Type Equivalence

- defines when two types in a PL/program are equivalent.

e.g. are float + int equivalent?

float a;

int b;

a = b; ??

## Two types of Type Equivalence

### Name Equivalence

expressions

two types are equivalent

if they have the same name.

e.g. above a = b is not allowed as float + int are different type names

### Structural Equivalence

Two types are structurally equivalent if they are made up of the same parts.

e.g. C      typedef float celsius;  
              typedef float fahr;

celsius c = 100.0;  
fahr f;

f = c; OK in C

since C uses structural equivalence.

## Java

```
public class Celsius {  
    }  
    private float temp;
```

3 // class Celsius

```
public class Fahr {  
    }  
    private float temp;
```

3 // class Fahr.

Celsius c = new Celsius(100.0);  
Fahr f;

f = c; X not allowed  
since Java uses  
name equivalence.

In fact in Java, we can then write

In Celsius:

```
public float toFahr() {  
    }
```

In Fahr

```
public float toCelsius() {  
    }
```

use

```
f = c.toCelsius();  
c = f.toFahr();
```

Some PLs have both!

e.g. Ada.

```
type celsius = real;  
type fahr = real;
```

In Ada, we use structural equivalence for these types.  
BUT, Ada also allows derived types

```
type celsius = new real;  
type fahr = new real;
```

Now celsius + fahr are new + different types even though they are structurally equivalent.

Another example w/ functions

C/Java

```
int max(int a, int b) {  
    ...  
    can be we do: float  
                    not x, y = ...  
                    max(x, y)?
```

Need to define something more practical than name/structural equivalence — Type Compatibility Rules

e.g. Two types are compatible if:

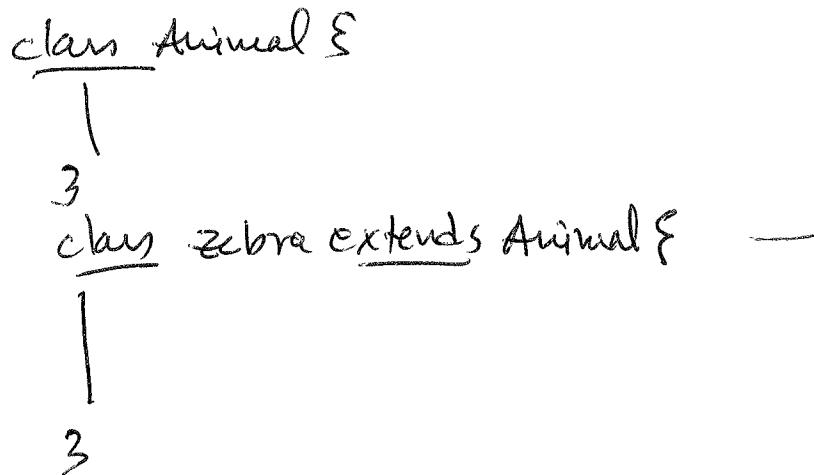
1. The two types are equivalent
2. One is a subtype of the other
3. Both are arrays with the same types of elements.

## Java type compatibility

1. Identical types are compatible

i.e. int and int  
float and float

2. Subtypes + Supertypes



zebra is a subtype of the supertype Animal

i.e. we can do

Animal a = new zebra(); ✓

But not Zebra z = new Animal(); X

3. Interfaces

any class that implements an interface is a subtype of that interface

4. Type Casting

- explicitly convert a variable to another

e.g. float f;

int i;

i = (int)f;

but cannot do

String s;

i = (int)s;

etc