

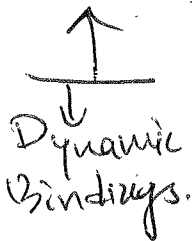
write this

Binding: associating a name to the thing it represents.  
 eg. variable: type, value, memory location, etc.

Binding Times: Time at which a binding is created.

- language definition time
- language implementation time
- Program writing time
- Compile Time
- Link Time
- Load Time
- Run Time

static bindings



~~Referencing Environment~~

Program (c)

```
#include <stdio.h>
int gcd(int a, int b) {
    while (a != b) {
        if (a > b)
            a = a - b;
        else
            b = b - a;
    }
    return a;
} // gcd()
```

```
int main() {
    int x, y;
    // input x, y
    int g = gcd(x, y);
    // print g
    return 0;
} // main()
```

write this

Bindings: Program Writing Time

gcd: function, 2 int params  
 returns int  
 a, b: int, parameters  
 main: function, no params  
 returns int  
 x, y: int  
 g: int

Run Time

// Input x, y  
 a = x, b = y  
 values of a + b  
 value of g

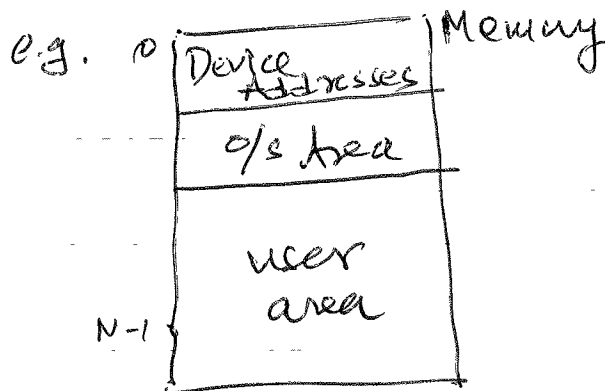
Lifetime: when + how long (i.e. time) during which a name has a binding.

Object Lifetime e.g. `int a;`

name 'a' has a storage (object) associated with it to store an int value.

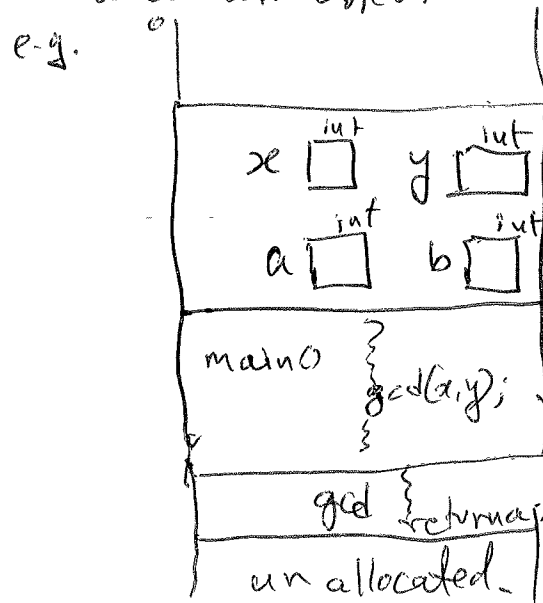
Time between creation + destruction of an object is object's lifetime.

Object lifetime depends on storage allocation scheme used by language implementation.



Static Allocation

when all objects are allocated before runtime.



- all memory locations are fixed at or before runtime.
- same locations are used for x, y for all calls of gcd()
- Lifetime of all objects is same as lifetime of program.

## Consequences:

- Could be potential source of errors

```
gcd(u) {  
  int w;  
  // do something with w  
}
```

upon return + reinvocation, last value of w is retained.

- Does not allow recursion!!

e.g.

```
int gcd(int a, int b) {  
  if (a == b)  
    return a;  
  if (a > b)  
    return gcd(a - b, b);  
  else  
    return gcd(a, b - a);  
}
```

e.g.

```
int fact(int n) {  
  if (n == 0)  
    return 1;  
  else  
    return n * fact(n - 1);  
}
```

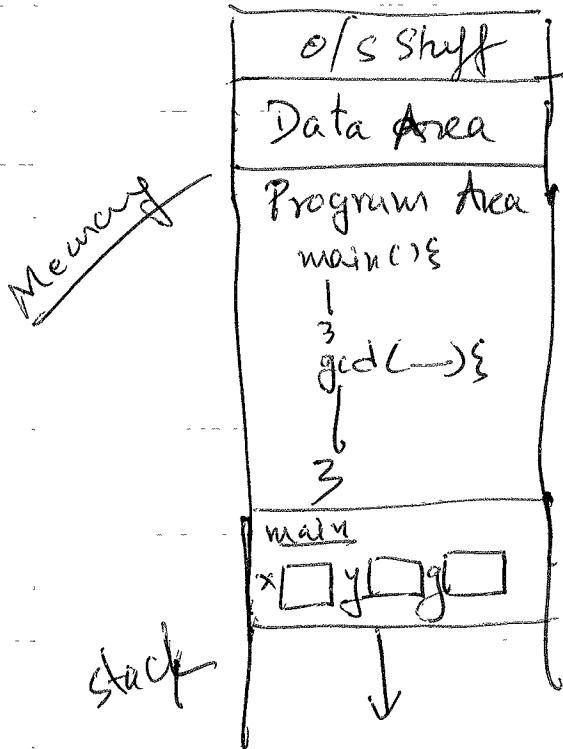
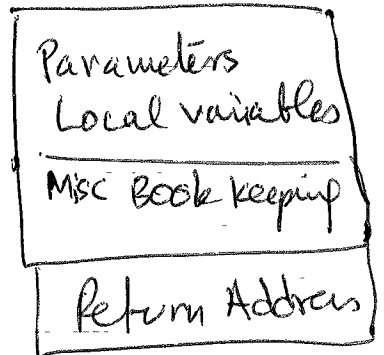
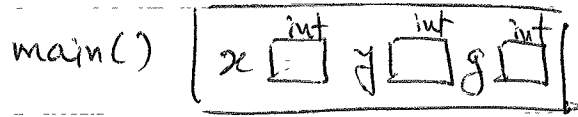
n | 1 2 3 fact(n)  
→ n \* fact(3)  
  n \* fact(2)  
  n \* fact(1)  
  n \* fact(0)  
  n    0

how to save the recursive calls!

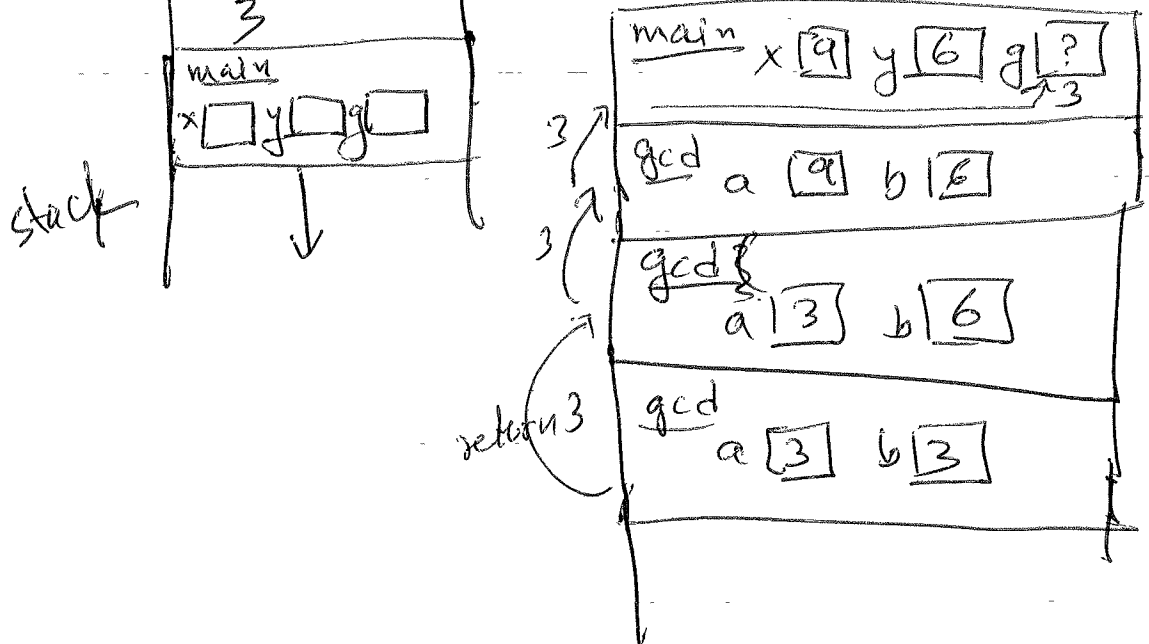
# Stack-based Allocation

Every function call has a Stack Frame.

(aka Activation records)

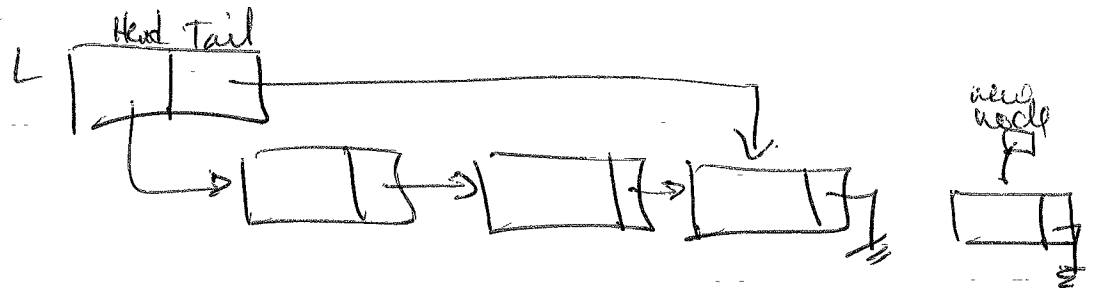


A new stack frame is created during runtime for each function call.  
 e.g. when computing gcd(9,6)



# Heap-based Allocation

For dynamically allocated data



JAVA

Node \* newNode = new Node (-);

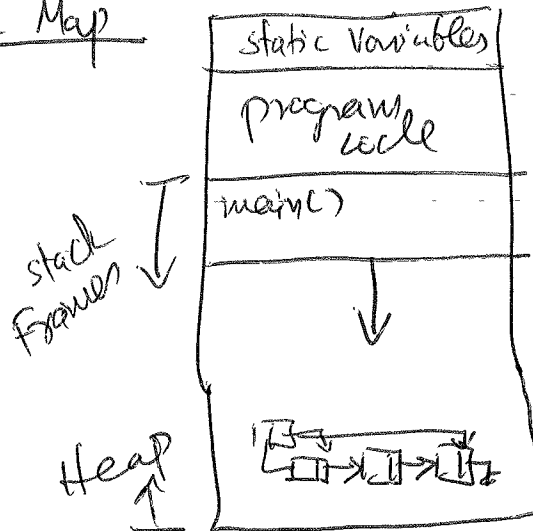
C Node \* newNode = (Node) malloc(sizeof(Node));

Q: Where is newNode allocated?

A: On a Heap of bytes in Memory.

Memory Map

When does stack meet heap stack smash!!



PL decision

- who/when memory allocation is done  
 (statically - no recursion)  
 (dynamically)

- How?

- explicit

Java: new

C: malloc

free

cfree

C++: dispose

new

- implicit

- Python

- What runtime support is needed?

- When allocating any checks done?

No - C

YES - JAVA, PYTHON, etc

- Can used space be reused?

No - C

YES - Java, Python

Needs: Runtime Garbage Collection! (less efficient)