

## Mathematical Induction

CS 231  
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## How do you climb infinite stairs?

- Not a rhetorical question!
- First, you get to the base platform of the staircase
- Then repeat:
  - From your current position, move one step up

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## What is induction?

- A method of proof:  $\forall n, n \geq a, P(n)$
- Three parts:
  - Base case(s): show it is true for one element
    - $P(a)$  (get to the stair's base platform)
  - Inductive hypothesis: assume it is true for any given element
    - Assume  $P(k), k \geq a$  (assume you are on a stair)
  - Show that it is true for the next highest element
    - $P(k+1)$  (show you can move to the next stair)

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## Why does induction work?

- Establish that the truth of a proposition follows from smaller instances of the same proposition:  $P(k) \rightarrow P(k+1)$
- Establish the truth of the smallest instance:  $P(a)$
- In induction, the truth percolates up through the layers to prove the whole proposition



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## Induction example

- Show that the sum of the first  $n$  odd integers is  $n^2$ 
  - Example: If  $n = 5, 1+3+5+7+9 = 25 = 5^2$
  - Formally, show:  $\forall n P(n)$  where  $P(n) = \sum_{i=1}^n 2i-1 = n^2$
- Base case: Show that  $P(1)$  is true

$$P(1) = \sum_{i=1}^1 2(i)-1 = 1^2$$

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## Induction example, continued

- Inductive hypothesis: assume true for  $k$ 
  - Thus, we assume that  $P(k)$  is true, or that
- Note: we don't yet know if this is true or not!
- Inductive step: show true for  $k+1$ 
  - We want to show that:

$$\sum_{i=1}^k 2i-1 = k^2$$

$$\sum_{i=1}^{k+1} 2i-1 = (k+1)^2$$

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## Induction example, continued

- Recall the inductive hypothesis:  $\sum_{i=1}^k 2^i - 1 = k^2$

- Proof of inductive step:  $\sum_{i=1}^{k+1} 2^i - 1 = (k+1)^2$

$$2(k+1) - 1 + \sum_{i=1}^k 2^i - 1$$

$$= 2(k+1) - 1 + k^2$$

$$= k^2 + 2k + 1$$

$$= (k+1)^2$$

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## Induction example

- Show that the sum of the first  $n$  powers of 2 is  $2^{n+1} - 1$ , where  $n$  starts at 0.

– Example: If  $n = 4$ :

$$- 1 + 2 + 2^2 + 2^3 + 2^4 = 1 + 2 + 4 + 8 + 16 = 31 = 2^5 - 1$$

– Formally, show:  $\forall n P(n)$  where  $P(n) = \sum_{i=0}^n 2^i = 2^{n+1} - 1$

- Base case: Show that  $P(0)$  is true

$$P(0) = \sum_{i=0}^0 2^i = 1 = 2^1 - 1$$

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## Induction example, continued

- Inductive hypothesis: assume true for arbitrary  $k$

– Thus, we assume that  $P(k)$  is true, or that

$$\sum_{i=0}^k 2^i = 2^{k+1} - 1$$

- Inductive step: show true for  $k+1$

– Want to show that:  $\sum_{i=0}^{k+1} 2^i = 2^{k+2} - 1$

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## Induction example, continued

- Recall the inductive hypothesis:  $\sum_{i=0}^k 2^i = 2^{k+1} - 1$

- Proof of inductive step:

$$\sum_{i=0}^{k+1} 2^i = 1 + 2 + \dots + 2^k + 2^{k+1} = \sum_{i=0}^k 2^i + 2^{k+1}$$

$$= 2^{k+1} - 1 + 2^{k+1} = 2 \times 2^{k+1} - 1 = 2^{k+2} - 1 \quad \blacksquare$$

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## What did we show

- Base case:  $P(0)$
- If  $P(k)$  is true, then  $P(k+1)$  is true
  - i.e.,  $P(k) \rightarrow P(k+1)$
- We know it's true for  $P(0)$
- Because of  $P(k) \rightarrow P(k+1)$ , if it's true for  $P(0)$ , then it's true for  $P(1)$
- Because of  $P(k) \rightarrow P(k+1)$ , if it's true for  $P(1)$ , then it's true for  $P(2)$
- Because of  $P(k) \rightarrow P(k+1)$ , if it's true for  $P(2)$ , then it's true for  $P(3)$
- Because of  $P(k) \rightarrow P(k+1)$ , if it's true for  $P(3)$ , then it's true for  $P(4)$
- And onwards to infinity
- Thus, it is true for all possible values of  $n$

$$\left[ P(0) \wedge \forall k (P(k) \rightarrow P(k+1)) \right] \rightarrow \forall n P(n)$$

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## How to do inductive proofs

- Show the base case
- Establish the inductive hypothesis**
- Manipulate the inductive step so that you can substitute in part of the inductive hypothesis**
- Prove the inductive step

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### Another induction example

- Show that  $n! < n^n, \forall n > 1$
- Base case:  $n = 2$   
 $2! < 2^2$   
 $2 < 4$
- Inductive hypothesis: assume  $k! < k^k$
- Inductive step: show that  $(k+1)! < (k+1)^{k+1}$

$(k+1)! = (k+1)k! < (k+1)k^k < (k+1)(k+1)^k = (k+1)^{k+1}$

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### Another induction example

- Show that  $6 | n^3 - n, \forall n \in \mathbb{Z}, n \geq 0$
- Base Case:  $P(0) : 6 | 0 = 0^3 - 0$
- Inductive hypothesis:  $6 | k^3 - k$
- Inductive step:  
 $(k+1)^3 - (k+1) = k^3 + 3k^2 + 3k + 1 - k - 1$   
 $= k^3 - k + 3k^2 + 3k = k^3 - k + 3(k^2 + k)$   
 $= k^3 - k + 3k(k+1)$  ■

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### Applications of Induction

- Algebraic (in)equalities are not the only suitable applications of induction.
- How to apply the inductive step is less obvious in non-algebraic applications of induction
  - the manipulation needed to apply the hypothesis is not algebraic

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### Square Cutting

- Prove that given two or more squares, one can always cut them and reform them into a large square.

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### Trominoes

- Polyomino – generalization of domino
- Tromino –
- Prove that if any one square is removed from a  $2^n \times 2^n$  checkerboard ( $n \geq 1, n \in \mathbb{Z}$ ), the remaining squares can be completely covered by L-shaped trominoes.

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### Trominoes

- $P(1)$ :  $2 \times 2$  board:
- Assume a  $2^k \times 2^k$  checkerboard can be covered except for any one square
- $P(k+1)$ :  $2^{k+1} \times 2^{k+1}$  checkerboard
  - divide into 4 quadrants
  - each quadrant is of size  $2^k \times 2^k$ .

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### All Horses are the Same Color

- If there is only one horse, it's of one color
- Assume within any set of  $k$  horses, there is only one color
- Consider  $k+1$  horse, and divide into sets of  $\{1, 2, 3, \dots, k\}$  and  $\{2, 3, 4, \dots, k, k+1\}$ .
  - Each is a set of  $k$  horses and can be of only one color.
  - Since there is overlap among the sets, there is only one color for all  $k+1$  horses.

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### All Men are Bald

- A man with 0 (or 1) hair is clearly bald
- Assume a man with  $k$  hairs is bald
- One more hair on a bald head does not cure baldness, thus a man with  $k+1$  hair is also bald.

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