









## Height

- If *T* is any binary tree of height *h* with *t* leaves (*h*>0), then  $t \le 2^h$
- Or  $\log_2 t \le h$
- Proof by strong induction on *h* 
  - P(0): T has only a root → 1 leaf:  $1 \le 2^{\circ}$
  - Assume P(i),  $0 \le i \le k$ : All binary trees with height less than or equal to k has at most  $2^k$  leaves

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-P(k+1): T is a binary tree of height k+1

## Induction Cont.

- $-k \ge 0 \rightarrow k+1 \ge 1$ , root has at least one child
- Root has exactly one child c:
  - The subtree rooted at c,  $T_c$  is of height k
  - By the inductive hypothesis,  $T_{\rm c}$  has at most  $2^k$  leaves
- Root has two children  $c_1$  and  $c_2$ :
  - One of the subtrees (say  $T_{c1}$ ) is of height k and  $T_{c2}$  is of any height between 0 and k
  - By the inductive hypothesis, both have at most 2<sup>k</sup> leaves, which gives a total of at most 2<sup>k+1</sup>.

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