A bit of humor: Input methods

Review

- 2's complement of *a*
- $2^n - a$
- 1's complement of *a*
- Algorithm
  - n-bit binary representation of *a*
  - negate all bits
  - add 1
- 2's complement of 27
- how many bits? – 8
  - 00011011_2
  - 11100100_2
  - 11100101_2
- Why does it work?
- How can you tell that a number is negative?

Translating from English

- Translate the statements:
  - "All hummingbirds are richly colored"
  - "No large birds live on honey"
  - "Birds that do not live on honey are dull in color"
  - "Hummingbirds are small"
- Assign our predicates
  - Let P(x) be "x is a hummingbird"
  - Let Q(x) be "x is large"
  - Let R(x) be "x lives on honey"
  - Let S(x) be "x is richly colored"
- Let our domain be all birds

Translating from English

- "All hummingbirds are richly colored"
  - $\forall x\ P(x) \rightarrow S(x)$
- "No large birds live on honey"
  - $\sim(\exists x\ Q(x) \wedge R(x)) \equiv \forall x \sim Q(x) \vee \sim R(x)$
  - $\forall x\ Q(x) \rightarrow \sim R(x) \equiv \forall x \sim Q(x) \vee \sim R(x)$
- "Birds that do not live on honey are dull in color"
  - $\forall x \sim R(x) \rightarrow \sim S(x)$
- "Hummingbirds are small"
  - $\forall x\ P(x) \rightarrow \sim Q(x)$

Restricted Quantifiers

- Arbitrary domain
- Universal
  - $\forall x\ P(x) \rightarrow Q(x)$ versus $\forall x\ P(x) \wedge Q(x)$
- Existential
  - $\exists x\ P(x) \wedge Q(x)$ versus $\exists x\ P(x) \rightarrow Q(x)$
- There exists a red dragon
- $\exists x$ dragon(x) $\rightarrow$ red(x)
- What if x is human? or duck?

Multiple Quantifiers

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Dianna Xu

Multiple quantifiers

- Let our domain be $\mathcal{R}$
- ∀x∃y P(x, y)
  - "For all x, there exists a y such that P(x,y)"
  - Example: ∀x∃y x+y == 0
- ∃x∀y P(x,y)
  - There exists an x such that for all y P(x,y) is true"
  - Example: ∃x∀y x*y == 0

Order of quantifiers

- ∃x∀y and ∀x∃y are not equivalent!
- P(x,y) = (x+y == 0)
  - ∃x∀y P(x,y) is false
  - ∀x∃y P(x,y) is true

Binding variables

- Let P(x,y) be x > y
- Consider: ∀x P(x,y)
  - This is not a proposition!
  - What is y?
    - If it's 5, then ∀x P(x,y) is false
    - If it's x-1, then ∀x P(x,y) is true
- y is a free variable - not "bound" by a quantifier

Binding variables 2

- (∃x P(x)) ∨ Q(x)
  - The x in Q(x) is not bound; thus not a proposition
- (∃x P(x)) ∨ (∀x Q(x))
  - Both x values are bound; thus it is a proposition
- (∃x P(x) ∧ Q(x)) ∨ (∀y R(y))
  - All variables are bound; thus it is a proposition
- (∃x P(x) ∧ Q(y)) ∨ (∀y R(y))
  - The y in Q(y) is not bound; this not a proposition

Translating between English and quantifiers

- The product of two negative integers is positive
  - ∀x∀y (x<0) ∧ (y<0) → (xy > 0)
- The average of two positive integers is positive
  - ∀x∀y (x>0) ∧ (y>0) → ((x+y)/2 > 0)
- The difference of two negative integers is not necessarily negative
  - ∃x∃y (x<0) ∧ (y<0) ∧ (x-y≥0)
- The absolute value of the sum of two integers does not exceed the sum of the absolute values of these integers
  - ∀x∀y |x+y| ≤ |x| + |y|

Translating between English and quantifiers

- ∃x∀y x+y = y
  - There exists an additive identity for all real numbers
- ∀x∀y ((x≥0) ∧ (y<0)) → (x-y > 0)
  - A non-negative number minus a negative number is greater than zero
- ∃x∃y ((x≤0) ∧ (y≤0)) ∧ (x-y > 0)
  - The difference between two non-positive numbers is not necessarily non-positive (i.e. can be positive)
- ∀x∀y ((x≠0) ∧ (y≠0)) ↔ (xy ≠ 0)
  - The product of two numbers is non-zero if and only if both factors are non-zero

Negating multiple quantifiers

- Recall negation rules for single quantifiers:
  - $\sim(\forall x\ P(x)) \equiv \exists x \sim P(x)$
  - $\sim(\exists x\ P(x)) \equiv \forall x \sim P(x)$
  - Essentially, you change the quantifiers, and negate what it's quantifying
- Examples:
  - $\sim(\forall x \exists y\ P(x,y))$
    - $\equiv \exists x \sim(\exists y\ P(x,y))$
    - $\equiv \exists x \forall y \sim P(x,y)$
  - $\sim(\forall x \exists y \forall z\ P(x,y,z))$
    - $\equiv \exists x \sim(\exists y \forall z\ P(x,y,z))$
    - $\equiv \exists x \forall y \sim(\forall z\ P(x,y,z))$
    - $\equiv \exists x \forall y \exists z \sim P(x,y,z)$

Negating multiple quantifiers 2

- Consider $\sim(\forall x \exists y\ P(x,y)) \equiv \exists x \forall y \sim P(x,y)$
  - The left side is saying "for all x, there exists a y such that P is true"
  - To disprove it (negate it), you need to show that "there exists an x such that for all y, P is false"
- Consider $\sim(\exists x \forall y\ P(x,y)) \equiv \forall x \exists y \sim P(x,y)$
  - The left side is saying "there exists an x such that for all y, P is true"
  - To disprove it (negate it), you need to show that "for all x, there exists a y such that P is false"

Negation examples

- Rewrite these statements so that the negations only appear within the predicates

a) $\sim(\exists y \exists x\ P(x,y))$
   1. $\forall y \sim(\exists x\ P(x,y))$
   2. $\forall y \forall x \sim P(x,y)$

b) $\sim(\forall x \exists y\ P(x,y))$
   1. $\exists x \sim(\exists y\ P(x,y))$
   2. $\exists x \forall y \sim P(x,y)$

c) $\sim(\exists y\ Q(y) \land \forall x \sim R(x,y))$
   1. $\forall y \sim(Q(y) \land \forall x \sim R(x,y))$
   2. $\forall y \sim Q(y) \lor \sim(\forall x \sim R(x,y))$
   3. $\forall y \sim Q(y) \lor \exists x\ R(x,y)$

Negation examples

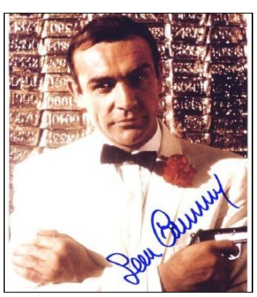
- Negate the following:

a) $\forall x \exists y \forall z\ T(x,y,z)$
   1. $\sim(\forall x \exists y \forall z\ T(x,y,z))$
   2. $\exists x \sim(\exists y \forall z\ T(x,y,z))$
   3. $\exists x \forall y \sim(\forall z\ T(x,y,z))$
   4. $\exists x \forall y \exists z \sim T(x,y,z)$

b) $\forall x \exists y\ P(x,y) \lor \forall x \exists y\ Q(x,y)$
   1. $\sim(\forall x \exists y\ P(x,y) \lor \forall x \exists y\ Q(x,y))$
   2. $\sim(\forall x \exists y\ P(x,y)) \land \sim(\forall x \exists y\ Q(x,y))$
   3. $\exists x \sim(\exists y\ P(x,y)) \land \exists x \sim(\exists y\ Q(x,y))$
   4. $\exists x \forall y \sim P(x,y) \land \exists x \forall y \sim Q(x,y)$

Negation

- There is a secret agent who appeals to all women
- Negation?
- For every secret agent there is a woman he doesn't appeal to.
- Common mistake: There is a secret agent who doesn't appeal to all women



Prolog

- A programming language using logic!
- Entering facts (propositions):

```
instructor(xu, cs231).
enrolled(alice, cs231).
enrolled(bob, cs231).
enrolled(claire, cs231).
```

- Extracting data

```
?- enrolled (alice, cs231).
```

Result:

```
yes
```

Prolog 2

- Extracting data
```
?- enrolled(X, cs231).
```
Result:
```
alice
bob
Claire
```
- Entering predicates:
```
teaches(P,S) :- instructor(P,C), enrolled(S,C).
```
- Extracting data
```
?- teaches(X, alice).
```
Result:
```
Xu
```


Arguments with Quantified Statements

CS 231

Dianna Xu

Vacuous Truth

- Presently, all men on the moon are happy.
- ∀x OnTheMoon(x) → Happy(x)
- There is no man on the moon presently.
- ∀x OnTheMoonPresently(x) → Happy(x)
- The statement is vacuously true.
- Presently, all men on the moon are dinosaurs.

Universal Instantiation

- ∀x ϵ D, P(x)
- x_0 ϵ D
- P(x_0)
- Example:
- All men are mortal.
- Socrates is a man.
- Socrates is mortal.

Existential Generalization

- P(x_0)
- x_0 ϵ D
- ∃x ϵ D, P(x)

Universal Modus Ponens

$$\begin{array}{ll} p & P(a) \\ \underline{p \rightarrow q} & \underline{\forall x, P(x) \rightarrow Q(x)} \\ \therefore q & \therefore Q(a) \end{array}$$

Universal Modus Tollens

$$\begin{array}{ll} \sim q & \sim Q(a) \\ \underline{p \rightarrow q} & \underline{\forall x, P(x) \rightarrow Q(x)} \\ \therefore \sim p & \therefore \sim P(a) \end{array}$$

2/2/16

Universal Transitivity

$$\begin{array}{l} \forall x, P(x) \rightarrow Q(x) \\ \underline{\forall x, Q(x) \rightarrow R(x)} \\ \therefore \forall x, P(x) \rightarrow R(x) \end{array}$$

2/2/16

Example of proof

- Given the hypotheses:
  - "Linda, a student in this class, owns a red convertible."
  - "Everybody who owns a red convertible has gotten at least one speeding ticket"
- Can you conclude: "Somebody in this class has gotten a speeding ticket"?

C(Linda)
R(Linda)

∀x (R(x)→T(x))

∃x (C(x)∧T(x))

2/2/16

Example of proof

1.	∀x (R(x)→T(x))	3rd hypothesis
2.	R(Linda) → T(Linda)	Universal instantiation using step 1
3.	R(Linda)	2nd hypothesis
4.	T(Linda)	Modes ponens using steps 2 & 3
5.	C(Linda)	1st hypothesis
6.	C(Linda) ∧ T(Linda)	Conjunction using steps 4 & 5
7.	∃x (C(x)∧T(x))	Existential generalization using step 6

Thus, we have shown that "Somebody in this class has gotten a speeding ticket"

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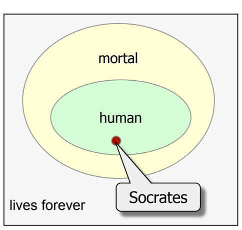
Diagrams for Validity



- To check the validity of an argument
- NOT a proof!

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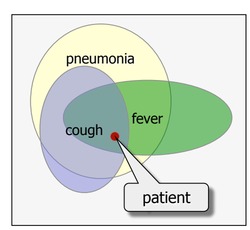
Abduction

- A form of logical inference that goes from observation to a hypothesis that accounts for the reliable data
- The lawn is wet → It rained last night



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Common Errors

- Converse error

$$\frac{\begin{array}{l} Q(a) \\ \forall x, P(x) \rightarrow Q(x) \end{array}}{\therefore P(a)}$$

- Inverse error

$$\frac{\begin{array}{l} \sim P(a) \\ \forall x, P(x) \rightarrow Q(x) \end{array}}{\therefore \sim Q(a)}$$

2/2/16

Example

- Anyone who grows a money tree is rich
- Bill Gates is rich
- Bill Gates grows a money tree

- Bill Gates does not grow a money tree
- Bill Gates is not rich

2/2/16

- Every Great American City Has At Least One College. Worcester Has Ten.
  - — Highway billboard in Worcester, MA

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