

## 1 Predicates and Quantifiers

We have seen how to represent properties of objects. For example,  $B(x)$  may represent that  $x$  is a student at Bryn Mawr College. Here  $B$  stands for “is a student at Bryn Mawr College” and  $x$  is a free variable that may be true for some values of  $x$  and false for others. As another example, if we want to express “ $x$  is a friend of  $y$ ”, we could use  $F(x, y)$  where  $F$  stand for “is a friend of” and both  $x$  and  $y$  are free variables. In these two examples,  $B$  and  $F$  are called *predicate symbols* and  $x$  and  $y$  are called *predicate variables*. A *predicate* is a predicate symbol together with suitable predicate variables.

**Definition 1** A predicate is a sentence that contains a finite number of variables and becomes a statement when specific values are substituted for the variables. The domain of a predicate variable is the set of all values that may be substituted in place of the variable.

Sometimes we would like to express whether a predicate  $P(x)$  is true for all values of  $x$  or some values of  $x$ . To express these ideas, we introduce two symbols, called *quantifiers*.

- Universal quantifier  $\forall$ : To say that the predicate  $P(x)$  is true for all possible value of  $x$ , we write  $\forall xP(x)$ , which is read “For all  $x$ ,  $P(x)$ ”.
- Existential quantifier  $\exists$ : To say that there exists or there is at least one value of  $x$  in the universe for which the predicate  $P(x)$  is true, we use  $\exists xP(x)$  and read “there exists an  $x$  such that  $P(x)$ ”. The statement  $\exists xP(x)$  is true means that the truth set is not equal to  $\emptyset$ .

For example, if we want to express “All human beings are mortal”, using the predicate  $M(x)$  to denote “human being  $x$  is mortal and  $H$  being the set of all human beings, then we write  $\forall x \in H, M(x)$ . The statement  $\forall x \in U, P(x)$  is true when the truth set of  $P(x)$  is the whole universe  $U$ .

**Example 1** What do the following formulas mean? Are they true or false?

1.  $\forall x(x^2 \geq 0)$ , where the universe of discourse is  $\mathbb{R}$ , the set of all real numbers.
2.  $\exists x(M(x) \wedge B(x))$ , where the universe of discourse is the set of all people,  $M(x)$  stands for the statement “ $x$  is a man,” and  $B(x)$  means “ $x$  has brown hair.”
3.  $\forall x(M(x) \rightarrow B(x))$ , with the same universe and the same meanings for  $M(x)$  and  $B(x)$ .
4.  $\forall xL(x, y)$ , where the universe is the set of all people, and  $L(x, y)$  means “ $x$  likes  $y$ .”

Note that in the statement  $\forall xL(x, y)$ , variable  $x$  is a bound variable and  $y$  is free. In general, even if  $x$  is a free variable in some statement  $P(x)$ , it is a bound variable in the statements  $\forall xP(x)$  and  $\exists xP(x)$ . Because of this, we say that quantifiers *bind* variables. Recall that a bound variable can always be replaced with a new variable without changing the meaning of the statement. Therefore  $\forall xL(x, y)$  is equivalent to  $\forall zL(z, y)$ . Both mean that everyone likes  $y$ . Words *all*, *every*, *every one*, and *everything* are usually an indication of using the universal quantifier  $\forall$ , and words *someone*, *exist*, and *something* usually indicate that the existantial quantifier  $\exists$  needs to be used.

**Example 2** Analyze the logical forms of the following statements.

1. *Someone didn't do the homework.*
2. *Everything in that store is either overpriced or poorly made.*

3. Susan likes everyone who dislikes Joe.
4.  $A \subseteq B$ .

**Example 3** Analyze the logical forms of the following statements.

1. Some students are married.
2. All parents are married.
3. Nobody likes a sore loser.
4. If a person in the dorm has a friend who has the measles, then everyone in the dorm will have to be quarantined.

**Exercise 1** Analyze the logical forms of the following statements.

1. Nobody's perfect.
2.  $A \cap B \subseteq B \setminus C$ .
3. If  $A \subseteq B$ , then  $A$  and  $C \setminus B$  are disjoint. Two sets are said to be disjoint if, and only if, they have no elements in common. Symbolically,  $S_1$  and  $S_2$  are disjoint iff  $S_1 \cap S_2 = \emptyset$ .

**Example 4** What do the following statements mean? Are they true or false? The universe of discourse in each case is  $\mathbb{R}$ , the set of all real numbers.

1.  $\forall x \exists y (x + y = 3)$
2.  $\exists y \forall x (x + y = 3)$

**Example 5** What do the following statements mean? Are they true or false? The universe of discourse in each case is  $\mathbb{N}$ , the set of all natural numbers.

1.  $\forall x \exists y (x < y)$
2.  $\exists y \forall x (x < y)$
3.  $\exists x \forall y (x < y)$
4.  $\forall y \exists x (x < y)$
5.  $\exists x \exists y (x < y)$
6.  $\forall x \forall y (x < y)$

## 2 Equivalence Involving Quantifiers

Quantifier Negation laws

$\neg\exists xP(x)$	is equivalent to	$\forall x\neg P(x)$
$\neg\forall xP(x)$	is equivalent to	$\exists x\neg P(x)$

**Example 6** Negate these statements and then reexpress the results as equivalent positive statements.

1.  $A \subseteq B$ .
2. Every student has a course he took and he doesn't like.

Necessary and Sufficient Conditions, Only If		
$\forall x, r(x)$ is a <b>sufficient condition</b> for $s(x)$	means	“ $\forall x$ , if $r(x)$ then $s(x)$ ”
$\forall x, r(x)$ is a <b>necessary condition</b> for $s(x)$	means	“ $\forall x$ , if not $r(x)$ then not $s(x)$ ”, or, equivalently, “ $\forall x$ , if $s(x)$ then $r(x)$ ”
$\forall x, r(x)$ <b>only if</b> $s(x)$	means	“ $\forall x$ , if not $s(x)$ then not $r(x)$ ”, or, equivalently, “ $\forall x$ , if $r(x)$ then $s(x)$ ”

**Example 7** Rewrite the following statements as quantified conditional statements. Do not use the word necessary or sufficient.

1. Squareness is a sufficient condition for rectangularity.
2. Being at least 35 years old is a necessary condition for being President of the United States.

We have seen from Example 1 that changing the order of two quantifiers can sometimes change the meaning of a statement. However, if the quantifiers are both  $\forall$  or both  $\exists$ , then the order can always be switched without affecting the meaning of the statement.

**Example 8** Analyze the logical forms of the following statements.

1. All friends like each other.
2. Everyone likes at least two people.
3. James likes exactly one person.

In general, for any set  $A$ , statement  $\forall x \in A P(x)$  means that for every value of  $x$  in the set  $A$ ,  $P(x)$  is true. Statement  $\exists x \in A P(x)$  means that there is some value of  $x$  that is in  $A$  and that also makes  $P(x)$  to be true. Therefore,  $\forall x \in A P(x)$  is equivalent to  $\forall x(x \in A \rightarrow P(x))$ , and  $\exists x \in A P(x)$  is equivalent to  $\exists x(x \in A \wedge P(x))$ .

Formal Logical Notation

$\forall x \text{ in } A, P(x)$	can be written as	$\forall x(x \in A \rightarrow P(x))$
$\exists x \text{ in } A \text{ such that } P(x)$	can be written as	$\exists x(x \in A \wedge P(x))$

Note that if  $A = \emptyset$ , then  $\exists x \in AP(x)$  is false no matter what  $P$  stands for. What about  $\forall x \in AP(x)$ ? Is it true?

$$\begin{aligned} & \forall x \in AP(x) \\ \text{is equivalent to } & \neg\neg\forall x \in AP(x) \quad \text{Double Negation law} \\ \text{is equivalent to } & \neg\exists x \in A\neg P(x) \quad \text{by Quantifier Negation law} \end{aligned}$$

If  $A = \emptyset$ , then  $\exists x \in A\neg P(x)$  must be false, and so  $\neg\exists x \in A\neg P(x)$  must be true. Therefore,  $\forall x \in AP(x)$  must true. Also note that since  $\forall x \in AP(x)$  is equivalent to  $\forall x(x \in A \rightarrow P(x))$ , if  $A = \emptyset$ ,  $x \in A$  is always false, and so the implication is always true. Such a statement is often said *vacuously* true.

As another note: universal quantifier distributes over conjunction. That is,  $\forall x(P(x) \wedge Q(x))$  is equivalent to  $\forall xP(x) \wedge \forall xQ(x)$ . Why? However, existential quantifier does not distributes over conjunction. In other words,  $\exists x(P(x) \wedge Q(x))$  is not equivalent to  $\exists xP(x) \wedge \exists xQ(x)$ . Why?

As an example of the distributive law for the universal quantifier and conjunction, let's consider the equation  $A = B$  where  $A$  and  $B$  are two sets. The equation  $A = B$  means  $\forall x(x \in A \leftrightarrow x \in B)$ , and so

$$\begin{aligned} & \forall x(x \in A \leftrightarrow x \in B) \\ \text{is equivalent to } & \forall x[(x \in A \rightarrow x \in B) \wedge (x \in B \rightarrow x \in A)] \quad \text{Def of Biconditional} \\ \text{is equivalent to } & \forall x(x \in A \rightarrow x \in B) \wedge \forall x(x \in B \rightarrow x \in A) \quad \text{Distributive law} \\ \text{is equivalent to } & A \subseteq B \wedge B \subseteq A \quad \text{Def of subset} \end{aligned}$$

**Example 9** Analyze the logical forms of the following statements where the universe of discourse is  $\mathbb{N}$ .

1.  $x$  is a perfect square.
2.  $x$  is the smallest number that is a multiple of both  $y$  and  $z$ .

**Exercise 2** Analyze the logical form of the statement “Every positive number has exactly two square roots”. The universe of discourse is  $\mathbb{R}$ .

**Exercise 3** Show that the statements  $A \subseteq B$  and  $A \setminus B = \emptyset$  are equivalent by writing each in logical symbols and then showing that the resulting formulas are equivalent.

### 3 More on Sets

#### 3.1 Elementhood Test Notation and Logical Statements

We have seen the elementhood test notation for defining a set. For example, the set of all perfect squares can be written as  $S = \{n^2 \mid n \in \mathbb{N}\}$ . This also can be written as  $\{x \mid \exists n \in \mathbb{N}(x = n^2)\}$ . Therefore, the logical statement  $x \in \{n^2 \mid n \in \mathbb{N}\}$  means the same thing as  $\exists n \in \mathbb{N}(x = n^2)$ .

**Example 10** Analyze the logical forms of the following statements.

1.  $y \in \{\sqrt[3]{x} \mid x \in \mathbb{Q}\}$ .
2.  $\{n^2 \mid n \in \mathbb{N}\}$  and  $\{n^3 \mid n \in \mathbb{N}\}$  are not disjoint.

### 3.2 Power Set

**Definition 2 (power set)** Suppose that  $A$  is a set. The power set of  $A$ , denoted  $\mathcal{P}(A)$ , is the set whose elements are all the subsets of  $A$ . In other words,  $\mathcal{P}(A) = \{x \mid x \subseteq A\}$ .

For example, the set  $A = \{1, 2\}$  has four subsets:  $\emptyset$ ,  $\{1\}$ ,  $\{2\}$ , and  $\{1, 2\}$ . Therefore,  $\mathcal{P}(A) = \{\emptyset, \{1\}, \{2\}, \{1, 2\}\}$ .

Why is  $\emptyset$  a subset of any set?

What is  $\mathcal{P}(\emptyset)$ ?

**Example 11** Analyze the logical forms of the following statements.

1.  $x \in \mathcal{P}(A)$ .
2.  $\mathcal{P}(A) \subseteq \mathcal{P}(B)$ .
3.  $x \in \mathcal{P}(A \cap B)$ .
4.  $x \in \mathcal{P}(A) \cap \mathcal{P}(B)$ .

### 3.3 Cartesian Products

**Definition 3** Let  $n$  be a positive integer and let  $x_1, x_2, \dots, x_n$  be (not necessarily distinct) elements. The ordered  $n$ -tuple,  $(x_1, x_2, \dots, x_n)$ , consists of  $x_1, x_2, \dots, x_n$  together with the ordering: first  $x_1$ , then  $x_2$ , and so forth up to  $x_n$ . An ordered 2-tuple is called an ordered pair, and an ordered 3-tuple is called an ordered triple.

Two ordered  $n$ -tuple  $(x_1, x_2, \dots, x_n)$  and  $(y_1, y_2, \dots, y_n)$  are equal if, and only if,  $x_1 = y_1, x_2 = y_2, \dots, x_n = y_n$ .

Symbolically:

$$(x_1, x_2, \dots, x_n) = (y_1, y_2, \dots, y_n) \leftrightarrow x_1 = y_1, x_2 = y_2, \dots, x_n = y_n.$$

In particular,  $(a, b) = (c, d) \leftrightarrow a = c$  and  $b = d$ .

For example,  $(1, 2, 3, 4) \neq (1, 3, 2, 4)$  and  $(3, (-2)^2, \frac{1}{2}) = (\sqrt{9}, 4, \frac{3}{6})$ .

**Definition 4** Given sets  $A_1, A_2, \dots, A_n$ , the Cartesian product of  $A_1, A_2, \dots, A_n$  denoted  $A_1 \times A_2 \times \dots \times A_n$ , is the set of all ordered  $n$ -tuples  $(a_1, a_2, \dots, a_n)$  where  $a_1 \in A_1, a_2 \in A_2, \dots, a_n \in A_n$ .

Symbolically:

$$A_1 \times A_2 \times \dots \times A_n = \{(a_1, a_2, \dots, a_n) \mid a_1 \in A_1, a_2 \in A_2, \dots, a_n \in A_n\}.$$

In particular,

$$A_1 \times A_2 = \{(a_1, a_2) \mid a_1 \in A_1, a_2 \in A_2\}$$

is the Cartesian product of  $A_1$  and  $A_2$ .

**Example 12** Let  $A_1 = \{x, y\}$ ,  $A_2 = \{1, 2, 3\}$ , and  $A_3 = \{a, b\}$ .

1. Find  $A_1 \times A_2$ .
2. Find  $(A_1 \times A_2) \times A_3$ .
3. Find  $A_1 \times A_2 \times A_3$ .

## 4 Set Identities

Recall that in Propositional Logic, we have the following logical equivalences.

Given any statement variables  $p, q$ , and  $r$ , a tautology  $\mathbf{t}$  and a contradiction  $\mathbf{c}$ , the following logical equivalences hold.

Commutative laws:	$p \wedge q \equiv q \wedge p$	$p \vee q \equiv q \vee p$
Associative laws:	$(p \wedge q) \wedge r \equiv p \wedge (q \wedge r)$	$(p \vee q) \vee r \equiv p \vee (q \vee r)$
Distributive laws:	$p \wedge (q \vee r) \equiv (p \wedge q) \vee (p \wedge r)$	$p \vee (q \wedge r) \equiv (p \vee q) \wedge (p \vee r)$
Identity laws:	$p \wedge \mathbf{t} \equiv p$	$p \vee \mathbf{c} \equiv p$
Negation laws:	$p \vee \neg p \equiv \mathbf{t}$	$p \wedge \neg p \equiv \mathbf{c}$
Double negative law:	$\neg(\neg p) \equiv p$	
Idempotent laws:	$p \wedge p \equiv p$	$p \vee p \equiv p$
Universal bound laws:	$p \vee \mathbf{t} \equiv \mathbf{t}$	$p \wedge \mathbf{c} \equiv \mathbf{c}$
De Morgan's laws:	$\neg(p \wedge q) \equiv \neg p \vee \neg q$	$\neg(p \vee q) \equiv \neg p \wedge \neg q$
Absorption laws:	$p \vee (p \wedge q) \equiv p$	$p \wedge (p \vee q) \equiv p$
Negations of $\mathbf{t}$ and $\mathbf{c}$ :	$\neg \mathbf{t} \equiv \mathbf{c}$	$\neg \mathbf{c} \equiv \mathbf{t}$

In Set Theory, we have *set identities* which are equations universally true for all elements in some set.

**Theorem 1 (Set Identities)** *Let the universal set be  $U$ . For every set  $A \subseteq U$ , set  $B \subseteq U$ , and set  $C \subseteq U$ ,*

<i>Commutative laws:</i>	$A \cup B = B \cup A$	$A \cap B = B \cap A$
<i>Associative laws:</i>	$(A \cup B) \cup C = A \cup (B \cup C)$	$(A \cap B) \cap C = A \cap (B \cap C)$
<i>Distributive laws:</i>	$A \cup (B \cap C) = (A \cup B) \cap (A \cup C)$	$A \cap (B \cup C) = (A \cap B) \cup (A \cap C)$
<i>Identity laws:</i>	$A \cup \emptyset = A$	$A \cap U = A$
<i>Complement laws:</i>	$A \cup A^c = U$	$A \cap A^c = \emptyset$
<i>Double Complement law:</i>	$(A^c)^c = A$	
<i>Idempotent laws:</i>	$A \cup A = A$	$A \cap A = A$
<i>Universal bound laws:</i>	$A \cup U = U$	$A \cap \emptyset = \emptyset$
<i>De Morgan's laws:</i>	$(A \cup B)^c = A^c \cap B^c$	$(A \cap B)^c = A^c \cup B^c$
<i>Absorption laws:</i>	$A \cup (A \cap B) = A$	$A \cap (A \cup B) = A$
<i>Complements of <math>U</math> and <math>\emptyset</math>:</i>	$U^c = \emptyset$	$\emptyset^c = U$
<i>Set Difference Law:</i>	$A \setminus B = A \cap B^c$	

With set identities, we can prove some set properties algebraically (Section 6.3).

**Example 13** *Construct an algebraic proof that for all sets  $A$  and  $B$ ,  $A \setminus (A \cap B) = A \setminus B$ .*

## 5 Proof Involving Quantifiers

How to prove a goal of the form  $\forall xP(x)$ ? If your proof works no matter what  $x$  was, then you can conclude that  $\forall xP(x)$  is true. This means that your proof should work for any value of  $x$ . Thus, you should start your proof without any assumption about  $x$ .

**Strategy 1 (for all):** (a) To prove a goal of the form  $\forall xP(x)$ , let  $x$  be an arbitrary object and prove  $P(x)$ . The letter  $x$  must be a new variable in the proof. If  $x$  is already being used in the proof to stand for something, then you must choose an unused variable, say  $y$ , to stand for the arbitrary object, and prove  $P(y)$ . (b) To use a given assumption of the form  $\forall xP(x)$ , plug in any value, say  $m$ , for  $x$  and use this assumption to conclude that  $P(m)$  is true. This rule is called *universal instantiation*.

Form of the Proof

Let  $x$  be arbitrary.

[Proof of  $P(x)$  goes here]

Since  $x$  was arbitrary, we can conclude that  $\forall xP(x)$ .

**Example 14** Suppose  $A, B$ , and  $C$  are sets, and  $A \setminus B \subseteq C$ . Prove that  $A \setminus C \subseteq B$ .

*Proof sketch.*

*Given:*  $A \setminus B \subseteq C$

*Goal:*  $A \setminus C \subseteq B$

*Analyze the logical form of the goal:*  $\forall x(x \in A \setminus C \rightarrow x \in B)$

*Let  $x$  be arbitrary.*

*New Goal:*  $x \in A \setminus C \rightarrow x \in B$

*Let  $x$  be arbitrary.*

[Proof of  $x \in A \setminus C \rightarrow x \in B$  goes here]

*Since  $x$  was arbitrary, we can conclude that  $\forall x(x \in A \setminus C \rightarrow x \in B)$ , so  $A \setminus C \subseteq B$ .*

The main advantage of using this strategy to prove a goal of the form  $\forall xP(x)$  is to prove a goal about all objects by reasoning about only one object, as long as that object is arbitrarily chosen.

**Example 15** Suppose  $A$  and  $B$  are sets. Prove that if  $A \cap B = A$  then  $A \subseteq B$ .

*Proof sketch.*

*Given:*  $A \cap B = A$

*Goal:*  $\forall x(x \in A \rightarrow x \in B)$

*Given:*  $A \cap B = A, x \in A$

*Goal:*  $x \in B$

*Suppose that  $A \cap B = A$ .*

*Let  $x$  be arbitrary.*

*Suppose that  $x \in A$ .*

[Proof of  $x \in B$  goes here]

*Therefore  $x \in A \rightarrow x \in B$ .*

*Since  $x$  was arbitrary, we can conclude that  $\forall x(x \in A \rightarrow x \in B)$ , so  $A \subseteq B$ .*

*Therefore, if  $A \cap B = A$ , then  $A \subseteq B$ .*

*When we write up the final proof we can skip some or all of the last three sentences.*

**Strategy 2 (exists):** (a) To prove a goal of the form  $\exists xP(x)$ , we need to find a value of  $x$  where  $P(x)$  is true. Let  $x$  be the value we decide and show that  $P(x)$  is true for this value. Same as in the case of  $\forall xP(x)$ ,  $x$  should be a new variable. If necessary, rename  $x$  into some unused variable, say  $y$ , and rewrite the goal in the equivalent form  $\exists yP(y)$ . (b) To use a given assumption of the form  $\exists xP(x)$ , we introduce a new variable, say  $x_0$ , to stand for an object for which  $P(x_0)$  is true. Now assume that  $P(x_0)$  is true and use this as an assumption. This rule is called *existential instantiation*.

Form of the Proof

Let  $x =$  the value we decide. .

[Proof of  $P(x)$  goes here]

Therefore,  $\exists xP(x)$ .

**Example 16** Prove that for every real number  $x$ , if  $x > 0$  then there is a real number  $y$  such that  $y(y + 1) = x$ .

*Proof sketch.*

*Goal:*  $\forall x(x > 0 \rightarrow \exists y[y(y + 1) = x])$

*Let  $x$  be an arbitrary number*

*Given:*  $x > 0$

*Goal:*  $\exists y[y(y + 1) = x]$

*We need to find a value of  $y$  such that  $y(y + 1) = x$ .*

*Solve the quadratic equation:*

Recall that the roots of a quadratic equation  $ax^2 + bx + c = 0$  can be obtained by using Quadratic Formula:  $x = \frac{-b \pm \sqrt{b^2 - 4ac}}{2a}$ .

**Definition 5 (Divisible)** If  $n$  and  $d$  are integers and  $d \neq 0$  then  $n$  is divisible by  $d$  if, and only if,  $n$  equals  $d$  times some integer.

Instead of “ $n$  is divisible by  $d$ ”, we can say that

$n$  is a multiple of  $d$ , or

$d$  is a factor of  $n$ , or

$d$  is a divisor of  $n$ , or

$d$  divides  $n$ .

The notation  $d|n$  is read “ $d$  divides  $n$ .” Symbolically, if  $n$  and  $d$  are integers and  $d \neq 0$ :

$$d|n \text{ iff } \exists k \in \mathbb{Z} \text{ such that } n = dk.$$

For example,  $4|20$ , since  $5 \times 4 = 20$ .

**Remark.** When does  $d$  not divide  $n$ ? An integer  $d$  does not divide  $n$ , denoted by  $d \nmid n$ , if and only if, for every integer  $k$ ,  $n \neq dk$ , or, in other words, the quotient  $n/d$  is not an integer.

Does  $k$  divides 0 where  $k$  is any nonzero integer?

**Example 17 (Transitivity of Divisibility)** For all integers  $a, b$ , and  $c$ , if  $a|b$  and  $b|c$ , then  $a|c$ .

**Definition 6 (Prime and Composite)** An integer  $n$  is prime if, and only if,  $n > 1$  and for all positive integers  $r$  and  $s$ , if  $n = rs$ , then either  $r$  or  $s$  equals  $n$ . An integer  $n$  is composite if, and only if,  $n > 1$  and  $n = rs$  for some integers  $r$  and  $s$  with  $1 < r < n$  and  $1 < s < n$ .

Is every integer greater than 1 either prime or composite?

**Example 18 (Divisibility by a Prime)** Prove that any integer  $n > 1$  is divisible by a prime number.

**Example 19** Is the following statement true or false? For all integers  $a$  and  $b$ , if  $a|b$  and  $b|a$  then  $a = b$ .

**Theorem 2 (Unique Factorization of Integers Theorem (Fundamental Theorem of Arithmetic))**

Given any integer  $n > 1$ , there exist a positive integer  $k$ , distinct prime numbers  $p_1, p_2, \dots, p_k$ , and positive integers  $e_1, e_2, \dots, e_k$  such that

$$n = p_1^{e_1} p_2^{e_2} p_3^{e_3} \cdots p_k^{e_k},$$

and any other expression for  $n$  as a product of prime numbers is identical to this except, perhaps, for the order in which the factors are written.



**Definition 7** Given any integer  $n > 1$ , the standard factored form of  $n$  is an expression of the form

$$n = p_1^{e_1} p_2^{e_2} p_3^{e_3} \cdots p_k^{e_k},$$

where  $k$  is a positive integer;  $p_1, p_2, \dots, p_k$  are prime numbers;  $e_1, e_2, \dots, e_k$  are positive integers; and  $p_1 < p_2 < \cdots < p_k$ .

**Example 20** Suppose  $m$  is an integer such that

$$8 \cdot 7 \cdot 6 \cdot 5 \cdot 4 \cdot 3 \cdot 2 \cdot m = 17 \cdot 16 \cdot 15 \cdot 14 \cdot 13 \cdot 12 \cdot 11 \cdot 10.$$

Does  $17|m$ ?

**Theorem 3 (The Quotient-Remainder Theorem)** Given any integer  $n$  and positive integer  $d$ , there exist unique integers  $q$  and  $r$  such that

$$n = dq + r \text{ and } 0 \leq r < d.$$

**Definition 8 (div and mod)** Given an integer  $n$  and a positive integer  $d$ ,

$n \text{ div } d$  = the integer quotient obtained when  $n$  is divided by  $d$   
 $n \text{ mod } d$  = the nonnegative integer remainder obtained when  $n$  is divided by  $d$ .

Symbolically, if  $n$  and  $d$  are integers and  $d > 0$ , then

$$n \text{ div } d = q \text{ and } n \text{ mod } d = r \leftrightarrow n = dq + r$$

where  $q$  and  $r$  are integers and  $0 \leq r < d$ .

**Definition 9 (Parity)** The parity of an integer refers to whether the integer is even or odd. We call the fact that any integer is either even or odd the parity property.

**Example 21** Prove that the square of any odd integer has the form  $8m + 1$  for some integer  $m$ .

**Definition 10 (mutually disjoint)** Sets  $A_1, A_2, A_3, \dots$  are mutually disjoint (or pairwise disjoint or nonoverlapping) if, and only if, no two sets  $A_i$  and  $A_j$  have any elements in common, where  $i \neq j$ . Symbolically, for all  $i, j = 1, 2, 3, \dots$ ,  $A_i \cap A_j = \emptyset$  whenever  $i \neq j$ .

**Definition 11 (partition)** A finite or infinite collection of nonempty sets  $\{A_1, A_2, A_3, \dots\}$  is a partition of a set  $A$  if, and only if,

1.  $A$  is the union of all the  $A_i$
2. The set  $A_1, A_2, A_3, \dots$  are mutually disjoint.

**Example 22** Let  $\mathbb{Z}$  be the set of all integers and let

- $T_0 = \{n \in \mathbb{Z} \mid n = 3k, \text{ for some integer } k\}$ ,
- $T_1 = \{n \in \mathbb{Z} \mid n = 3k + 1, \text{ for some integer } k\}$ , and
- $T_2 = \{n \in \mathbb{Z} \mid n = 3k + 2, \text{ for some integer } k\}$ .

Is  $\{T_0, T_1, T_2\}$  a partition of  $\mathbb{Z}$ ?

**Definition 12** For any real number  $x$ , the absolute value of  $x$ , denoted  $|x|$ , is defined as follows:

$$|x| = \begin{cases} x & \text{if } x \geq 0 \\ -x & \text{if } x < 0. \end{cases}$$

**Lemma 1** For all real numbers  $r$ ,  $-|r| \leq r \leq |r|$ .

**Lemma 2** For all real numbers  $r$ ,  $|-r| = |r|$ .

**Example 23 (The Triangle Inequality)** Prove that for all real numbers  $x$  and  $y$ ,  $|x + y| \leq |x| + |y|$ .

More examples on how to apply these strategies will be given when we learn predicate logic.

Sometimes when proving a biconditional  $P \leftrightarrow Q$ , the steps for proving  $P \rightarrow Q$  is same as the steps for proving  $Q \rightarrow P$  in a reversed order. In this case, we could simplify the proof by writing it as a string of equivalences, starting with  $P$  and ending with  $Q$ .

**Example 24** Suppose  $A, B$ , and  $C$  are sets. Prove that  $A \cap (B \setminus C) = (A \cap B) \setminus C$ .

*Proof Sketch.*

$x \in A \cap (B \setminus C)$  iff  $x \in A \wedge x \in B \setminus C$  iff  $x \in A \wedge x \in B \wedge x \notin C$ ;

$x \in (A \cap B) \setminus C$  iff  $x \in A \cap B \wedge x \notin C$  iff  $x \in A \wedge x \in B \wedge x \notin C$ .

Proof by cases is one of the strategies we discussed before when a given assumption is of the form  $P \vee Q$ .

**Example 25** Suppose that  $A, B$ , and  $C$  are sets. Prove that if  $A \subseteq C$  and  $B \subseteq C$  then  $A \cup B \subseteq C$ .

*Proof Sketch.*

*Given:*  $A \subseteq C, B \subseteq C$

*Goal:*  $\forall x(x \in A \cup B \rightarrow x \in C)$

*Let  $x$  be arbitrary and assume that  $x \in A \cup B$*

*Given:*  $A \subseteq C, B \subseteq C, x \in A \vee x \in B$

*Goal:*  $x \in C$

## Acknowledgement

Much of the content of the notes is based on the following books:

- *Discrete Mathematics with Applications, 4th Edition* by Susanna Epp, Cengage Learning 2010.
- *How to Prove It: A Structured Approach, 2nd Edition* by Daniel J. Velleman, Cambridge University Press 2006