## CS206

## Final set of review questions

Hashtables and Maps

- Which of the hash table collision-handling schemes could tolerate a load factor above 1 and which could not?
- Draw the 11 -entry hash table that results from using the hash function, $h(i)=$ $(3 i+5) \bmod 11$, to hash the keys $12,44,13,88,23,94,11,39,20,16$, and 5 , assuming collisions are handled by chaining.
- What is the result of the previous exercise, assuming collisions are handled by linear probing?
- Show the result of Q2 assuming collisions are handled by quadratic probing, up to the point where the method fails.
- What is the result of Exercise R-10.6 when collisions are handled by double hashing using the secondary hash function $h^{\prime}(k)=7-(k \bmod 7)$ ?
- What is the worst-case time for putting $n$ entries in an initially empty hash table, with collisions resolved by chaining? What is the best case?
- Explain why a hash table is not suited to implement a sorted map.
- Describe how to perform a removal from a hash table that uses linear probing to resolve collisions where we do not use a special marker to represent deleted elements. That is, we must rearrange the contents so that it appears that the removed entry was never inserted in the first place.


## Binary Search Trees and AVL Trees

- If we insert the entries $(1, A),(2, B),(3, C),(4, D)$, and $(5, E)$, in this order, into an initially empty binary search tree, what will it look like?
- Insert, into an empty binary search tree, entries with keys $30,40,24,58,48,26$, 11, 13 (in this order). Draw the tree after each insertion.
- How many different binary search trees can store the keys $\{1,2,3\}$ ?
- Dr. Amongus claims that the order in which a fixed set of entries is inserted into a binary search tree does not matter-the same tree results every time. Give a small example that proves he is wrong.
- Dr. Amongus claims that the order in which a fixed set of entries is inserted into an AVL tree does not matter-the same AVL tree results every time. Give a small example that proves he is wrong.

- Draw the AVL tree resulting from the insertion of an entry with key 52 into the AVL tree of the above figure.
- Draw the AVL tree resulting from the removal of the entry with key 62 from the AVL tree of Figure 11.13b.
- Explain why you would get the same output in an inorder listing of the entries in a binary search tree, $T$, independent of whether $T$ is maintained to be an AVL tree, or not.
- Explain how to use an AVL tree or a to sort $n$ comparable elements in $O(n \log n)$ time in the worst case.


## Graphs

- Would you use the adjacency matrix structure or the adjacency list structure in each of the following cases? Justify your choice.
- The graph has 10,000 vertices and 20,000 edges, and it is important to use as little space as possible.
- The graph has 10,000 vertices and $20,000,000$ edges, and it is important to use as little space as possible.
- You need to answer the query getEdge $(u, v)$ as fast as possible, no matter how much space you use.
- Explain why the DFS traversal runs in $O\left(n^{2}\right)$ time on an $n$-vertex simple graph that is represented with the adjacency matrix structure.
- Computer networks should avoid single points of failure, that is, network vertices that can disconnect the network if they fail. We say an undirected, connected graph $G$ is biconnected if it contains no vertex whose removal would divide $G$ into two or more connected components. Give an algorithm for adding at most $n$ edges to a connected graph $G$, with $n \geq 3$ vertices and $m \geq n-1$ edges, to guarantee that $G$ is biconnected. Your algorithm should run in $O(n+m)$ time.

- Draw an adjacency matrix representation of the undirected graph shown above. You may add labels to links if you find it convenient
- Draw an adjacency list representation of the undirected graph shown above. You may add labels to links if you find it convenient
- $\quad$ Suppose we represent a graph $G$ having $n$ vertices and $m$ edges with the edge list structure. Why, in this case, does the insertVertex method run in $O(1)$ time while the removeVertex method runs in $O(m)$ time
- Let $G$ be an undirected graph with $n$ vertices and $m$ edges. Write a method traversing each edge of $G$ exactly once in each direction. The method should run in $\mathrm{O}(\mathrm{m}+\mathrm{n})$ time. You may use choose a graph representation. How would this algorithm be changed for a directed graph?
- Let $G$ be an undirected graph whose vertices are the integers 1 through 8 , and let the adjacent vertices of each vertex be given by the table below:
vertex
- $\quad 1(2,3,4)$
- $2(1,3,4)$
- $3(1,2,4)$
- $4(1,2,3,6)$
- $5(6,7,8)$
- $6(4,5,7)$
- $7(5,6,8)$
- $8(5,7)$

Assume that, in a traversal of $G$, the adjacent vertices of a given vertex are returned in the same order as they are listed in the table above.

- Draw $G$.
- Give the sequence of vertices of $G$ visited using a DFS traversal starting at vertex 1 .
- Give the sequence of vertices visited using a BFS traversal starting at vertex 1 .
- The time delay of a long-distance call is determined by the number of communication links on the telephone network between the caller and callee. The engineers of RT\&T want to compute the maximum possible time delay that may be experienced in a long-distance call. Write method that computes the maximum number of links (when cycles are disallowed, between some caller A and every other caller.
- For the implementation of Graph, Link and Node add a method public void removeLink(Link<E> link) in an appropriate place. Where is that "appropriate" place?
- Reimplement Dijkstra's shortest path method (in the Graph class on the class website) using a Stack rather than a Priority Queue. Doing so saves some time in that stacks have $\mathrm{O}(1)$ time for push and pop whereas PriorityQueues have $\mathrm{O}(\lg \mathrm{n})$ time for the equivalent operations. How do the other adjustments you needed to make to the function to ensure that you got the shortest path affect its runtime?
- Write a method to do a breadth first traversal for a graph (starting at a node). Why might a breadth first traversal be preferred to a depth first traversal?

