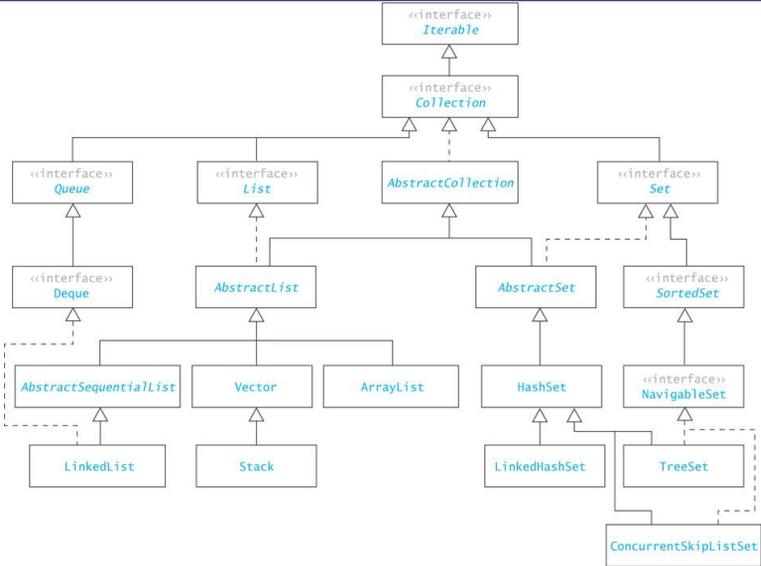


# CHAPTER 3 & 4

## Stacks & Queues

### The Collection Framework



## Stack Abstract Data Type

- A stack is one of the most commonly used data structures in computer science
- A stack can be compared to a Pez dispenser
  - ▣ Only the top item can be accessed
  - ▣ You can extract only one item at a time
- The top element in the stack is the last added to the stack (most recently)
- The stack's storage policy is *Last-In, First-Out*, or *LIFO*



## Java Collections: Stack

- The Java API includes a `Stack` class as part of the package

`java.util :`

```
Stack<String> myStringStack = new Stack<String>();
Stack<Place> myPlacesStack = new Stack<Places>();

myStringStack.push("Deepak");

myPlacesStack.push(new Place("19010", "Bryn Mawr", "PA"));

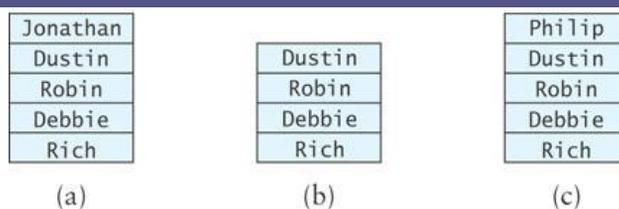
etc.
```

## Specification of the Stack Abstract Data Type

- Only the top element of a stack is visible; therefore the number of operations performed by a stack are few
- We need the ability to
  - ▣ test for an empty stack (`empty()`)
  - ▣ inspect the top element (`peek()`)
  - ▣ retrieve the top element (`pop()`)
  - ▣ put a new element on the stack (`push()`)

Methods	Behavior
<code>boolean empty()</code>	Returns <b>true</b> if the stack is empty; otherwise, returns <b>false</b> .
<code>E peek()</code>	Returns the object at the top of the stack without removing it.
<code>E pop()</code>	Returns the object at the top of the stack and removes it.
<code>E push(E obj)</code>	Pushes an item onto the top of the stack and returns the item pushed.

## A Stack of Strings



- “Rich” is the oldest element on the stack and “Jonathan” is the youngest (Figure a)
- `String last = names.peek();` stores a reference to “Jonathan” in `last`
- `String temp = names.pop();` removes “Jonathan” and stores a reference to it in `temp` (Figure b)
- `names.push(“Philip”);` pushes “Philip” onto the stack (Figure c)

## Other examples of stacks

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- Back button in browser
- Palindrome checker  
Go hang a salami, I'm a lasagna hog!
- Matching parentheses
- Expression evaluation
- `printStackTrace()`

## Queue

- The queue, like the stack, is a widely used data structure
- A queue differs from a stack in one important way
  - A stack is LIFO list – *Last-In, First-Out*
  - while a queue is FIFO list, *First-In, First-Out*

## Queue Abstract Data Type

- A queue can be visualized as a line of customers waiting for service
- The next person to be served is the one who has waited the longest
- New elements are placed at the end of the line



## Print Queue

- Operating systems use queues to
  - ▣ keep track of tasks waiting for a scarce resource
  - ▣ ensure that the tasks are carried out in the order they were generated
- Print queue: printing is much slower than the process of selecting pages to print, so a queue is used

Document Name	Status	Owner	Pages	Size	Submitted	P
Microsoft Word - Queues_Paul_1007.doc		Paul Wolfgang	52	9.75 MB	1:53:18 PM 10/7/2003	
Microsoft Word - Stacks.doc		Paul Wolfgang	46	9.05 MB	1:53:57 PM 10/7/2003	
Microsoft Word - Trees2.doc		Paul Wolfgang	54	38.4 MB	1:54:41 PM 10/7/2003	

3 document(s) in queue

## Specification for a Queue Interface

Method	Behavior
<code>boolean offer(E item)</code>	Inserts <code>item</code> at the rear of the queue. Returns <b>true</b> if successful; returns <b>false</b> if the item could not be inserted.
<code>E remove()</code>	Removes the entry at the front of the queue and returns it if the queue is not empty. If the queue is empty, throws a <code>NoSuchElementException</code> .
<code>E poll()</code>	Removes the entry at the front of the queue and returns it; returns <b>null</b> if the queue is empty.
<code>E peek()</code>	Returns the entry at the front of the queue without removing it; returns <b>null</b> if the queue is empty.
<code>E element()</code>	Returns the entry at the front of the queue without removing it. If the queue is empty, throws a <code>NoSuchElementException</code> .

- The `Queue` interface implements the `Collection` interface (and therefore the `Iterable` interface), so a full implementation of `Queue` must implement all required methods of `Collection` (and the `Iterable` interface)

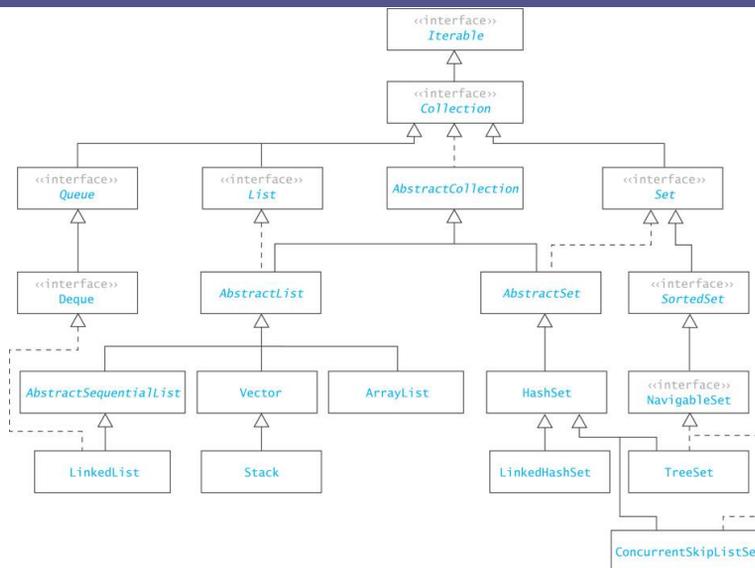
## Class `LinkedList` Implements the `Queue` Interface

- The `LinkedList` class provides methods for inserting and removing elements at either end of a double-linked list, which means all `Queue` methods can be implemented easily
- The Java 5.0 `LinkedList` class implements the `Queue` interface
 

```
Queue<String> names = new LinkedList<String>();
```

  - ▣ creates a new `Queue` reference, `names`, that stores references to `String` objects
  - ▣ The actual object referenced by `names` is of type `LinkedList<String>`, but because `names` is a type `Queue<String>` reference, you can apply only the `Queue` methods to it

# The Collection Framework



## Java Collections: Queue

- The Java API includes a `Queue` interface as part of the package

`java.util :`

```

Queue<String> myStringQueue = new LinkedList<String>();
Queue<Place> myPlacesQueue = new LinkedList<Places>();

myStringQueue.offer("Deepak");

myPlacesQueue.offer(new Place("19010", "Bryn Mawr", "PA"));

etc.

```

## Examples of Queues

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- Simulations of real life situations: Service Queues
- Scheduling processes in Operating Systems
- Keep track of state in systematic searches

## Stacks & Queues

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### java.util.Stack<E>

boolean empty()  
 E peek()  
 E pop()  
 ■ Both raise  
   EmptyStackException  
 E push(e)  
 + all List<E> operations

### java.util.Queue<E>

boolean add(e)  
 boolean offer(e)  
 E remove()  
 E poll()  
 E peek()  
 E element()  
 - Return T/F/null  
 ■ Raise  
 NoSuchElementException



## Stack Applications

### Section 3.2

## Finding Palindromes

- **Palindrome:** a string that reads identically in either direction, letter by letter (ignoring case)
  - ▣ kayak
  - ▣ "I saw I was I"
  - ▣ "Able was I ere I saw Elba"
  - ▣ "Level madam level"
  
- **Problem:** Write a program that reads a string and determines whether it is a palindrome

## Finding Palindromes (cont.)

Data Fields	Attributes
<code>private String inputString</code>	The input string.
<code>private Stack&lt;Character&gt; charStack</code>	The stack where characters are stored.
Methods	Behavior
<code>public PalindromeFinder(String str)</code>	Initializes a new <code>PalindromeFinder</code> object, storing a reference to the parameter <code>str</code> in <code>inputString</code> and pushing each character onto the stack.
<code>private void fillStack()</code>	Fills the stack with the characters in <code>inputString</code> .
<code>private String buildReverse()</code>	Returns the string formed by popping each character from the stack and joining the characters. Empties the stack.
<code>public boolean isPalindrome()</code>	Returns <b>true</b> if <code>inputString</code> and the string built by <code>buildReverse</code> have the same contents, except for case. Otherwise, returns <b>false</b> .

## Finding Palindromes (cont.)

```
import java.util.*;

public class PalindromeFinder {
    private String inputString;
    private Stack<Character> charStack = new
        Stack<Character>();

    public PalindromeFinder(String str) {
        inputString = str;
        fillStack(); // fills the stack with the characters in
                    inputString
    }
    ...
}
```

## Finding Palindromes (cont.)

- Solving using a stack:
  - ▣ Push each string character, from left to right, onto a stack



k a y a k

```
private void fillStack() {
    for(int i = 0; i < inputString.length(); i++) {
        charStack.push(inputString.charAt(i));
    }
}
```

## Finding Palindromes (cont.)

- Solving using a stack:
  - ▣ Pop each character off the stack, appending each to the `StringBuilder` result



k a y a k

```
private String buildReverse(){
    StringBuilder result = new StringBuilder();
    while(!charStack.empty()) {
        result.append(charStack.pop());
    }
    return result.toString();
}
```

## Finding Palindromes (cont.)

...

```
public boolean isPalindrome() {  
    return inputString.equalsIgnoreCase(buildReverse());  
}  
}
```

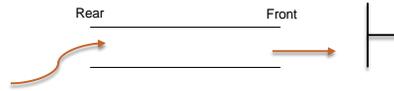
## Queue Applications

Discrete Event Simulation

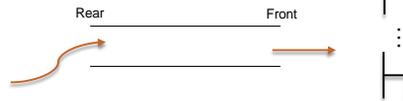
# Discrete Event Simulation

25

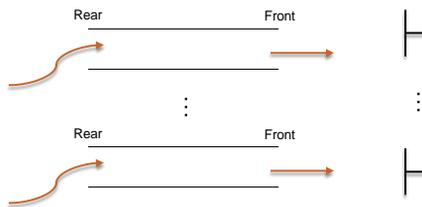
- Single Queue, single server



- Single Queue, multiple servers



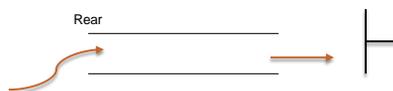
- Multiple Queue, multiple servers



## Example: Single Queue, Single Server

26

- Arrival process
  - ▣ How customers arrive: What is inter-arrival time?  
E.g. between 1-3 min
  - ▣ Service mechanism: How long will service take?  
E.g. 0.5 to 2.0 min
  - ▣ Queue characteristics: FIFO



## Example Data

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Customer	Inter-arrival Time	Service Time (min)
C1	1.9	1.7
C2	1.3	1.8
C3	1.1	1.5
C4	1.0	0.9

### Queue Simulation

T	Arrival	Queue	Server	Depart
0		[]	Idle	
1.9	C1	[]	C1	
3.2	C2	[C2]	C1	
3.6		[]	C2	C1
4.3	C3	[C3]	C2	
5.3	C4	[C4, C3]	C2	
5.4		[C4]	C3	C2
6.9		[]	C4	C2
7.8		[]		C4

## Application: Lab Printer Simulation

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- There is one printer in the Computer Science Lab
- At any given time, there may be as many as 10 students working in the lab
- Each student may print up to twice in an hour
- Print jobs are 1-20 pages long
- ∴ There are up to 20 print jobs in an hour
- Question: What is the chance that in any given second there will be a print job scheduled?

## Application: Lab Printer Simulation

29

- There is one printer in the Computer Science Lab
- At any given time, there may be as many as 10 students working in the lab
- Each student may print upto twice in an hour
- Print jobs are 1-20 pages long
- ∴ There are up to 20 print jobs in an hour
- Question: What is the chance that in any given second there will be a print job scheduled?

$$\frac{20}{1 \text{ hour}} * \frac{1 \text{ hour}}{60 \text{ min}} * \frac{1 \text{ min}}{60 \text{ sec}} = \frac{20}{3600} = \frac{1 \text{ task}}{180 \text{ sec}}$$

## Application: Lab Printer Simulation

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- There is one printer in the CS Lab (10 ppm)
- At any given time, there may be as many as 10 students working in the lab
- Each student may print upto twice in an hour
- Print jobs are 1-20 pages long
- ∴ There are up to 20 print jobs in an hour
- Question: What will the average wait time be for students to receive their printouts?
- Question: What would the average wait time be if the printer were upgraded to 20 ppp?

## Implementing a Stack

### Section 3.3

## Java Collections: Stack

- The Java API includes a `Stack` class as part of the package

`java.util :`

```
Stack<String> myStringStack = new Stack<String>();
Stack<Place> myPlacesStack = new Stack<Places>();

myStringStack.push("Deepak");

myPlacesStack.push(new Place("19010", "Bryn Mawr", "PA"));

etc.
```



## Implementing a Stack with a List Component

- We can write a class, `ListStack`, that has a `List` component (in the example below, `theData`)
- We can use either the `ArrayList`, or the `LinkedList` classes, as all implement the `List` interface. The `push` method, for example, can be coded as

```
public E push(E obj) {
    theData.add(obj);
    return obj;
}
```

- A class which adapts methods of another class by giving different names to essentially the same methods (`push` instead of `add`) is called an *adapter class*
- Writing methods in this way is called *method delegation*

## Implementing a Stack Using an Array

- If we implement a stack as an array we would need . . .

```
public class ArrayStack<E> implements Stack<E> {
    private E[] theData;
    int topOfStack = -1;
    private static final int INITIAL_CAPACITY = 10;

    public ArrayStack() {
        this(INITIAL_CAPACITY);
    }

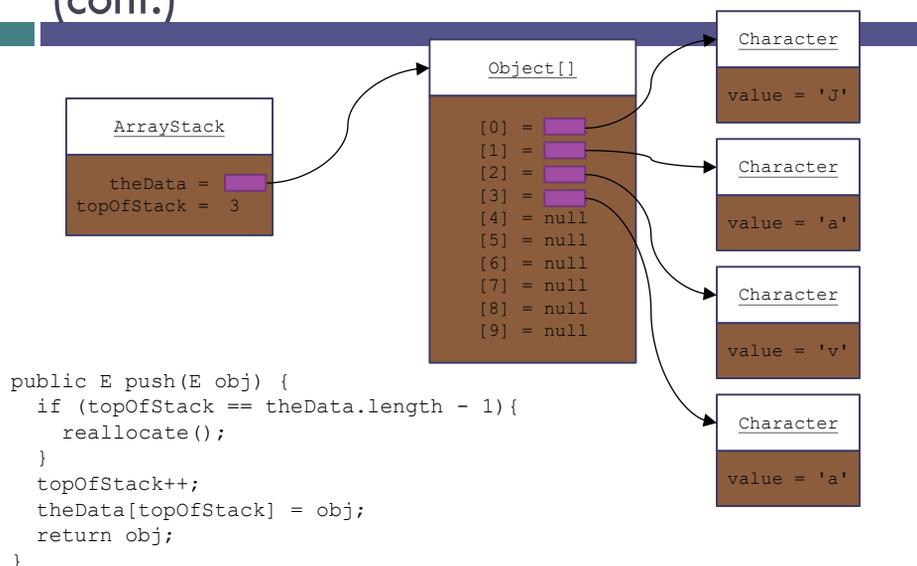
    public ArrayStack(int capacity) {
        theData = new E[capacity];
    } //
}
```

Allocate storage for an array with a default

Keep track of the top of the stack (subscript of the element at the top of the stack; for empty stack = -1)

There is no size variable or method

## Implementing a Stack Using an Array (cont.)



## Implementing a Stack Using an Array (cont.)

```

public E pop() {
    if (empty()) {
        throw new EmptyStackException();
    }
    return theData[topOfStack--];
} // pop()

```

## Implementing a Stack using an array

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```
import java.util.EmptyStackException;
public class ArrayStack<E> implements StackInt<E> {

    E[] theData;
    int topOfStack = -1; // Initially empty stack.
    private static final int INITIAL_CAPACITY = 10;

    public ArrayStack() {
        theData = (E[]) new Object[INITIAL_CAPACITY];
    } // ArrayStack()

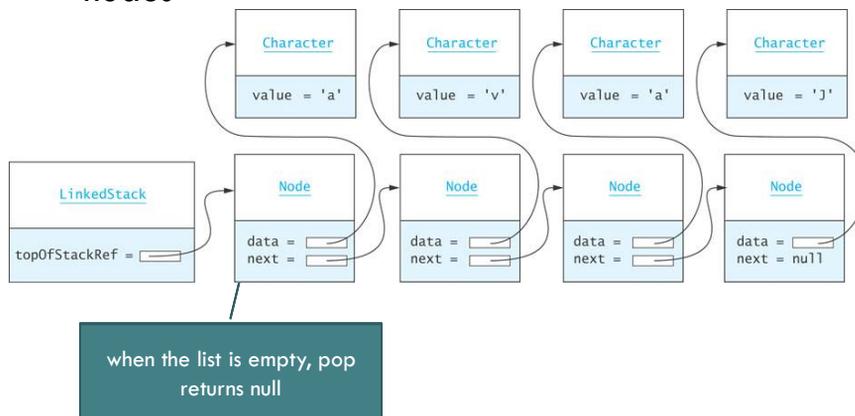
    public E push(E obj) {
        if (topOfStack == theData.length - 1) {
            reallocate();
        }
        topOfStack++;
        theData[topOfStack] = obj;
        return obj;
    } // push()

    public E pop() {
        if (empty()) {
            throw new EmptyStackException();
        }
        return theData[topOfStack--];
    } // pop()

} // class ArrayStack<E>
```

## Implementing a Stack as a Linked Data Structure

- We can also implement a stack using a linked list of nodes



## Implementing a Stack as a Linked Data Structure (cont.)

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- [Listing 3.5](#) (`LinkedStack.java`, pages 168 - 169)

## Comparison of Stack Implementations

- The easiest implementation uses a `List` component (`ArrayList` is the simplest) for storing data
  - ▣ An underlying array requires reallocation of space when the array becomes full, and
  - ▣ an underlying linked data structure requires allocating storage for links
  - ▣ As all insertions and deletions occur at one end, they are constant time,  $O(1)$ , regardless of the type of implementation used

## Additional Stack Applications

### Section 3.4

## Additional Stack Applications

- Postfix and infix notation
  - ▣ Expressions normally are written in infix form, but
  - ▣ it easier to evaluate an expression in postfix form since there is no need to group sub-expressions in parentheses or worry about operator precedence

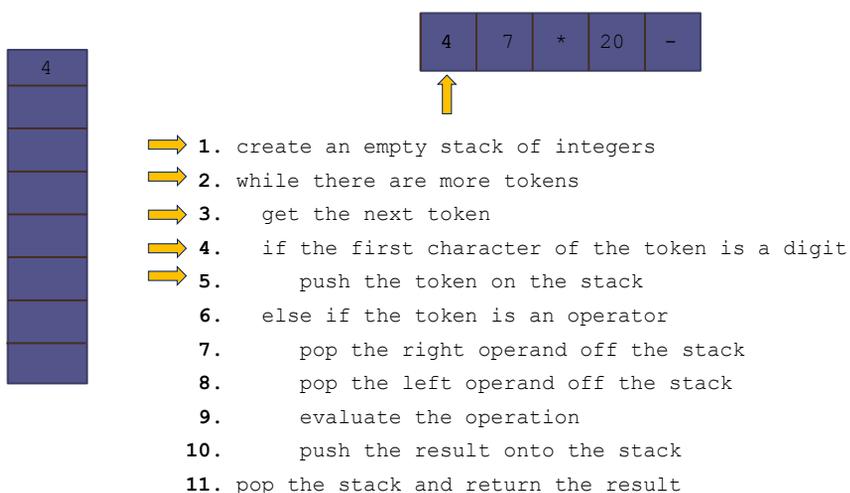
Postfix Expression	Infix Expression	Value
$4\ 7\ *$	$4 * 7$	28
$4\ 7\ 2\ +\ *$	$4 * (7 + 2)$	36
$4\ 7\ * \ 20\ -$	$(4 * 7) - 20$	8
$3\ 4\ 7\ * \ 2\ / \ +$	$3 + ((4 * 7) / 2)$	17

## Evaluating Postfix Expressions

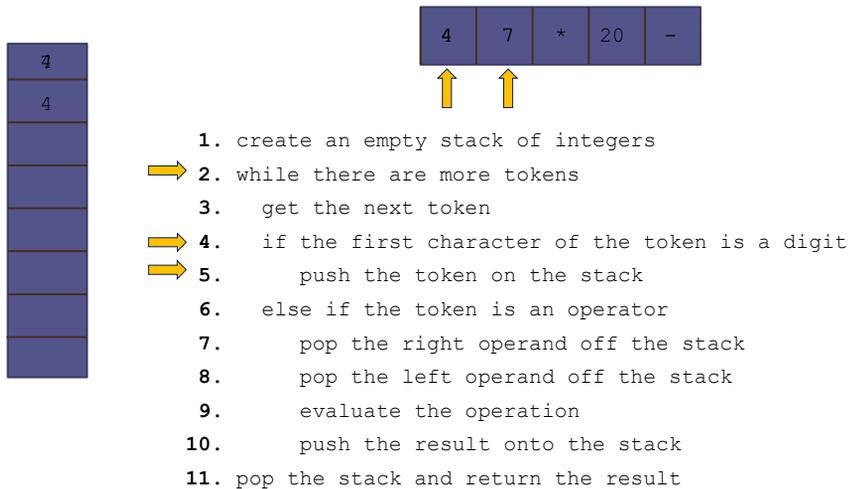
- Write a class that evaluates a postfix expression
- Use the space character as a delimiter between tokens

Data Field	Attribute
Stack<Integer> operandStack	The stack of operands (Integer objects).
Method	Behavior
public int eval(String expression)	Returns the value of expression.
private int evalOp(char op)	Pops two operands and applies operator op to its operands, returning the result.
private boolean isOperator(char ch)	Returns <b>true</b> if ch is an operator symbol.

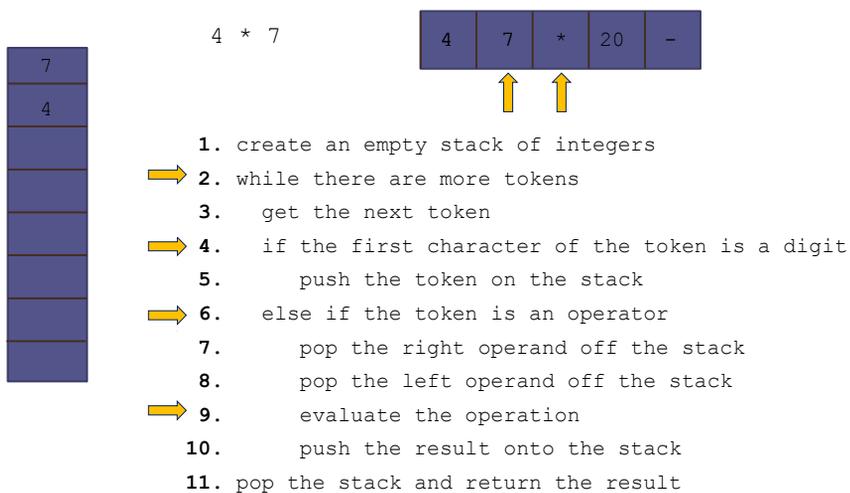
## Evaluating Postfix Expressions (cont.)



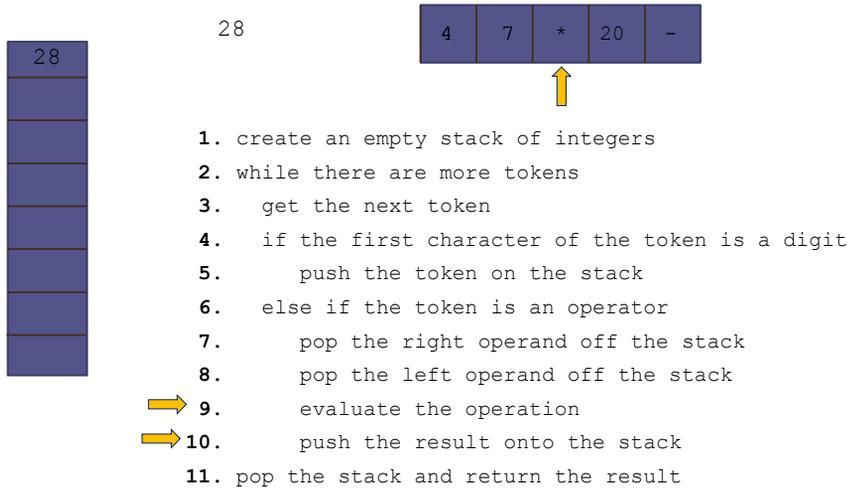
## Evaluating Postfix Expressions (cont.)



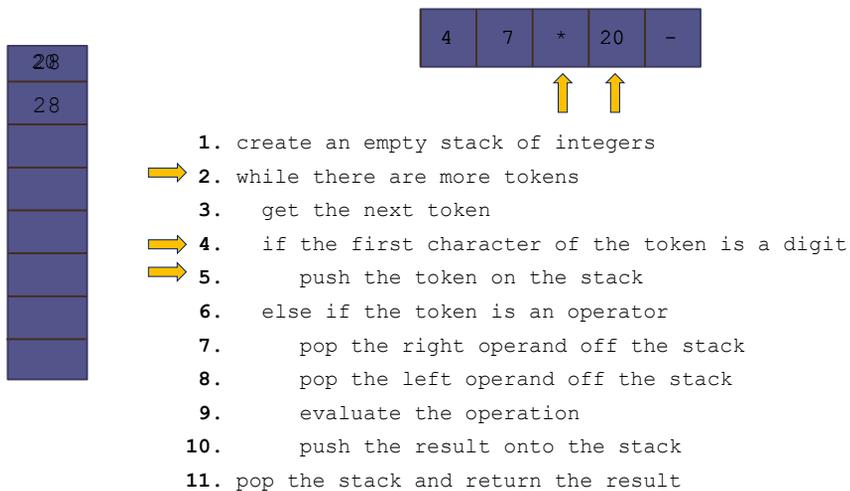
## Evaluating Postfix Expressions (cont.)



## Evaluating Postfix Expressions (cont.)



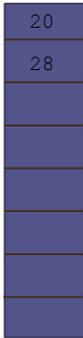
## Evaluating Postfix Expressions (cont.)





## Evaluating Postfix Expressions (cont.)

28 - 20



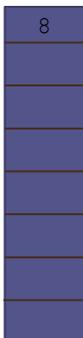
4	7	*	20	-
---	---	---	----	---

↑      ↑

1. create an empty stack of integers
- 2. while there are more tokens
3. get the next token
- 4. if the first character of the token is a digit
5. push the token on the stack
- 6. else if the token is an operator
7. pop the right operand off the stack
8. pop the left operand off the stack
- 9. evaluate the operation
10. push the result onto the stack
11. pop the stack and return the result

## Evaluating Postfix Expressions (cont.)

8

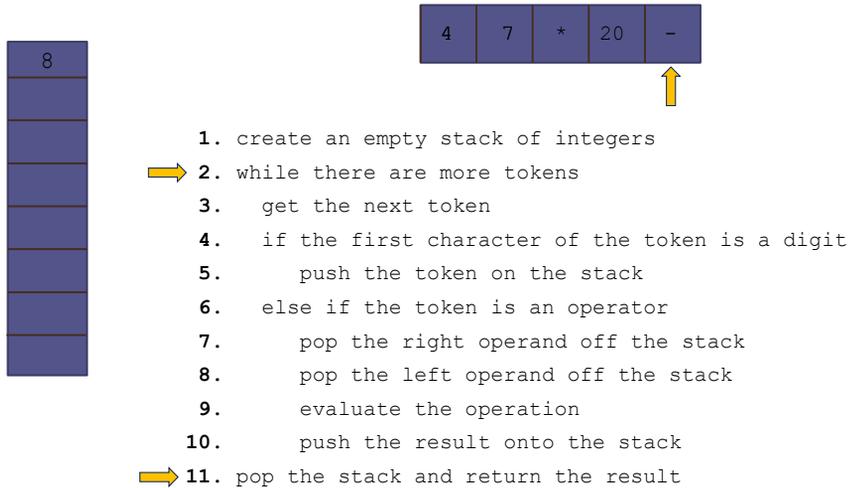


4	7	*	20	-
---	---	---	----	---

↑

1. create an empty stack of integers
2. while there are more tokens
3. get the next token
4. if the first character of the token is a digit
5. push the token on the stack
6. else if the token is an operator
7. pop the right operand off the stack
8. pop the left operand off the stack
- 9. evaluate the operation
- 10. push the result onto the stack
11. pop the stack and return the result

## Evaluating Postfix Expressions (cont.)



## Evaluating Postfix Expressions (cont.)

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- [Listing 3.6 \(PostfixEvaluator.java, pages 173 - 175\)](#)

## Converting from Infix to Postfix

- Convert infix expressions to postfix expressions
- Assume:
  - ▣ expressions consists of only spaces, operands, and operators
  - ▣ space is a delimiter character
  - ▣ all operands that are identifiers begin with a letter or underscore
  - ▣ all operands that are numbers begin with a digit

Data Field	Attribute
private Stack<Character> operatorStack	Stack of operators.
private StringBuilder postfix	The postfix string being formed.
Method	Behavior
public String convert(String infix)	Extracts and processes each token in infix and returns the equivalent postfix string.
private void processOperator(char op)	Processes operator op by updating operatorStack.
private int precedence(char op)	Returns the precedence of operator op.
private boolean isOperator(char ch)	Returns <b>true</b> if ch is an operator symbol.

## Converting from Infix to Postfix

(cont.)

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- Example: convert
 
$$w - 5.1 / \text{sum} * 2$$
 to its postfix form
 
$$w 5.1 \text{sum} / 2 * -$$

## Converting from Infix to Postfix (cont.)

Next Token	Action	Effect on operatorStack	Effect on postfix
w	Append w to postfix.	$\square$	w
-	The stack is empty Push - onto the stack	$\begin{array}{ c } \hline - \\ \hline \end{array}$	w
5.1	Append 5.1 to postfix	$\begin{array}{ c } \hline - \\ \hline \end{array}$	w 5.1
/	precedence(/) > precedence(-), Push / onto the stack	$\begin{array}{ c } \hline / \\ \hline - \\ \hline \end{array}$	w 5.1
sum	Append sum to postfix	$\begin{array}{ c } \hline / \\ \hline - \\ \hline \end{array}$	w 5.1 sum
*	precedence(*) equals precedence(/) Pop / off of stack and append to postfix	$\begin{array}{ c } \hline - \\ \hline \end{array}$	w 5.1 sum /

## Converting from Infix to Postfix (cont.)

Next Token	Action	Effect on operatorStack	Effect on postfix
*	precedence(*) > precedence(-), Push * onto the stack	$\begin{array}{ c } \hline * \\ \hline - \\ \hline \end{array}$	w 5.1 sum /
2	Append 2 to postfix	$\begin{array}{ c } \hline * \\ \hline - \\ \hline \end{array}$	w 5.1 sum / 2
End of input	Stack is not empty, Pop * off the stack and append to postfix	$\begin{array}{ c } \hline - \\ \hline \end{array}$	w 5.1 sum / 2 *
End of input	Stack is not empty, Pop - off the stack and append to postfix	$\square$	w 5.1 sum / 2 * -

## Converting from Infix to Postfix (cont.)

### Algorithm for Method `convert`

1. Initialize `postfix` to an empty `StringBuilder`.
2. Initialize the operator stack to an empty stack.
3. **while** there are more tokens in the infix string
4.     Get the next token.
5.     **if** the next token is an operand
6.         Append it to `postfix`.
7.     **else if** the next token is an operator
8.         Call `processOperator` to process the operator.
9.     **else**
10.         Indicate a syntax error.
11. Pop remaining operators off the operator stack and append them to `postfix`.

## Converting from Infix to Postfix (cont.)

### Algorithm for Method `processOperator`

1. **if** the operator stack is empty
2.     Push the current operator onto the stack.
3.     **else**
4.         Peek the operator stack and let `topOp` be the top operator.
5.         **if** the precedence of the current operator is greater than the precedence of `topOp`
6.             Push the current operator onto the stack.
7.         **else**
8.             **while** the stack is not empty and the precedence of the current operator is less than or equal to the precedence of `topOp`
9.                 Pop `topOp` off the stack and append it to `postfix`.
10.             **if** the operator stack is not empty
11.                 Peek the operator stack and let `topOp` be the top operator.
12.             Push the current operator onto the stack.

## Converting from Infix to Postfix (cont.)

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- [Listing 3.7](#) (`InfixToPostfix.java`, pages 181 - 183)

## Converting from Infix to Postfix (cont.)

- Testing
  - ▣ Use enough test expressions to satisfy yourself that the conversions are correct for properly formed input expressions
  - ▣ Use a driver to catch `InfixToPostfix.SyntaxErrorException`
- [Listing 3.8](#) (`TestInfixToPostfix.java`, page 184)

## Converting Expressions with Parentheses

---

- The ability to convert expressions with parentheses is an important (and necessary) addition
- Modify `processOperator` to push each opening parenthesis onto the stack as soon as it is scanned
- When a closing parenthesis is encountered, pop off operators until the opening parenthesis is encountered
- [Listing 3.9 \(InfixToPostfixParens.java, pages 186 - 188\)](#)