
Queue

Based on the notes from David Fernandez-Baca and Steve Kautz

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CS206 Intro to Data Structures



Queue

- A **queue** is a list that operates under the **first-in first-out** (*FIFO*) policy:
 - read or remove only the item at the **front** (**head**) of the queue
 - add an item only to the **back** (**tail**) of the queue
 - examine the front item.
- `java.util` contains a `Queue<E>` interface that contains all the methods you would expect from a FIFO queues, as well as other kinds of queues. Java offers several implementations, for example the `LinkedList` class.

The Java Queue Interface

- **E element()**. Retrieves, but does not remove, the head of this queue. Throws `NoSuchElementException` if this queue is empty.
- **E peek()**. Retrieves, but does not remove, the head of this queue, or returns `null` if this queue is empty.
- **boolean add(E e)**. Inserts the specified element into this queue if it is possible to do so immediately without violating capacity restrictions, returning `true` upon success and throwing an `IllegalStateException` if no space is currently available

The Java Queue Interface (cont.)

- **boolean offer(E e)**. Inserts the specified element into this queue if it is possible to do so immediately without violating capacity restrictions.
- **E poll()**. Retrieves and removes the head of this queue, or returns null if this queue is empty.
- **E remove()**. Retrieves and removes the head of this queue. Throws `NoSuchElementException` if this queue is empty.

Since `Queue<E>` extends `Collection<E>`, it inherits all of the latter's methods, including `isEmpty()`, `size()`, and `iterator()`.

Implementation – Linked List

- A queue is easily implemented as a singly-linked list with a tail pointer.
- It is perhaps even better to use a *circular* list. In this case, a pointer to the last node also gives easy access to the first node, by following one link. Thus we can handle the structure by a single pointer, instead of two.

Implementation – Array-Based

We use an array a to store the elements. Additionally, we have two indices:

- first: points to first element of queue (front)
- last: points to first available slot in the array (just before the back)

We initialize $\text{first} = \text{last} = 0$. The queue is empty when $\text{first} == \text{last}$.

- To enqueue, put the new item in $A[\text{last}]$ and increment last.
- To dequeue, return $A[\text{first}]$ and increment first.

Implementation – Circular Array

- A potential problem using array:
After a series of enqueue/dequeue operations, both first and last are at the end of the array even though the array is not full.
- Solution: treat the array as being circular.
That is, when $\text{last} == \text{A.length}$ and we need to increment last to insert a new item, we just reset last to 0.
$$\text{last} = (\text{last} + 1) \% \text{A.length};$$

Note that we still use the convention that the queue is empty when $\text{first} == \text{last}$. This means that at least one entry of array A will always have to be left unused. Otherwise, we wouldn't be able to distinguish between an empty queue and a full one.