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CS206

Trees

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# Lab

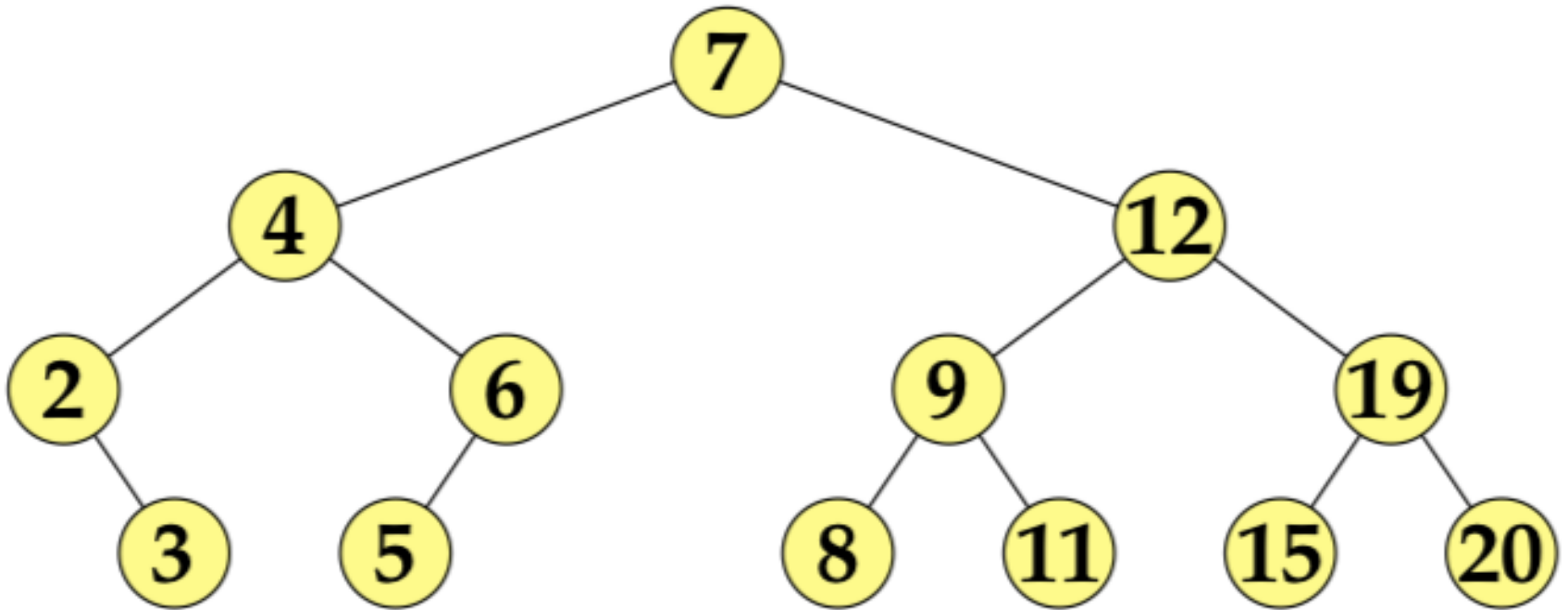
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- Given the data:
- 7, 4, 12, 2, 6, 9, 19, 3, 5, 8, 11, 15, 20
- Draw the binary tree
- What the height of the tree?
- If the data were re-arranged, what is the shortest possible tree?

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# Traversals / Printing

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# Postorder traversal

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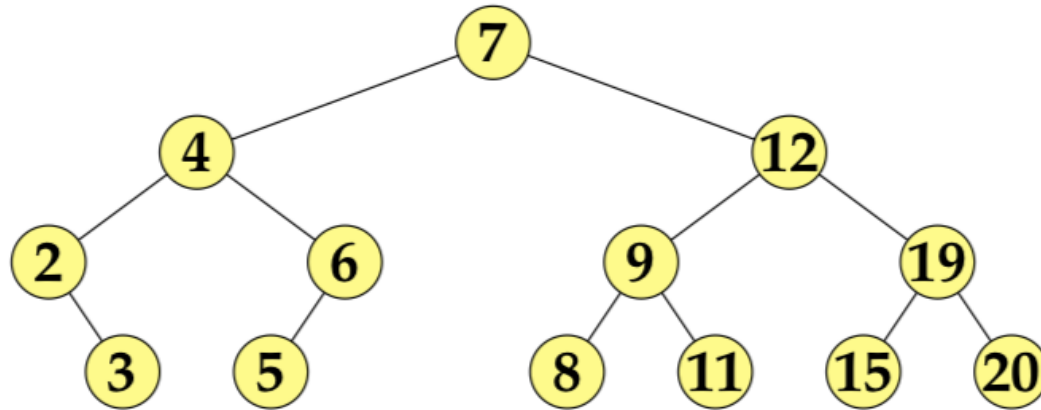
```
public void printPostOrder() {  
    iPrintPostOrder(root, 0);  
    System.out.println();  
}
```

```
private void iPrintPostOrder(Node treePart, int depth) {  
    if (treePart==null) return;  
    iPrintPostOrder(treePart.left, depth+1);  
    iPrintPostOrder(treePart.right, depth+1);  
    System.out.print("[ "+treePart.payload+", "+depth+" ]");  
}
```

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# Breadth First traversal

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0 [7]  
1 [4 12]  
2 [2 6 9 19]  
3 [3 5 8 11 15 20]

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# Remove

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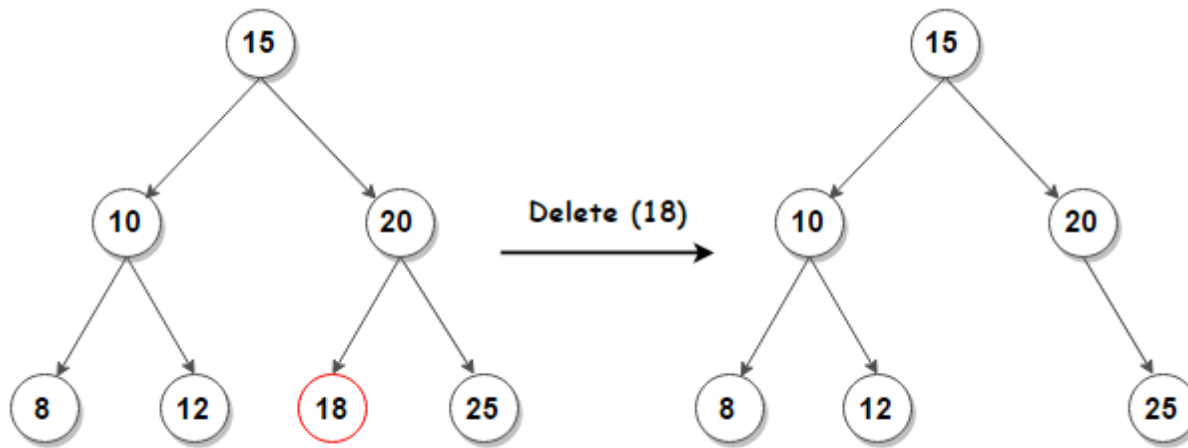
- `boolean remove(E element);`
- returns true if element existed and was removed and false otherwise
- Cases
  - element not in tree
  - element is a leaf
  - element has one child
  - element has two children

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# Leaf

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- Just delete

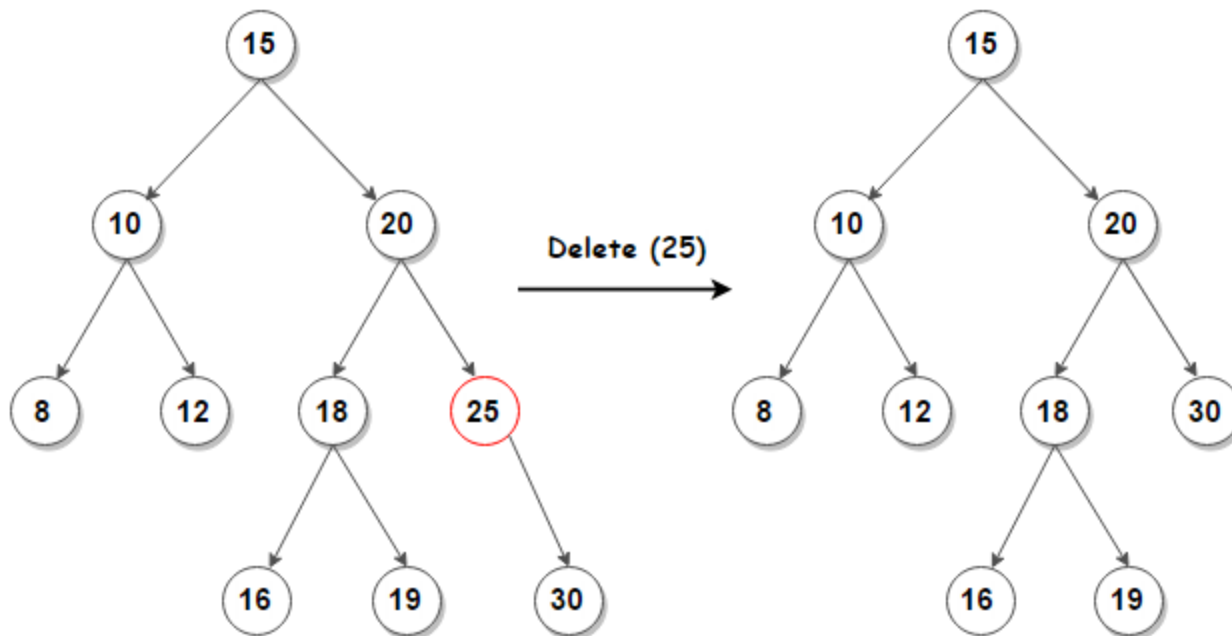


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# One child

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- Replace with child – skip over like in linked list





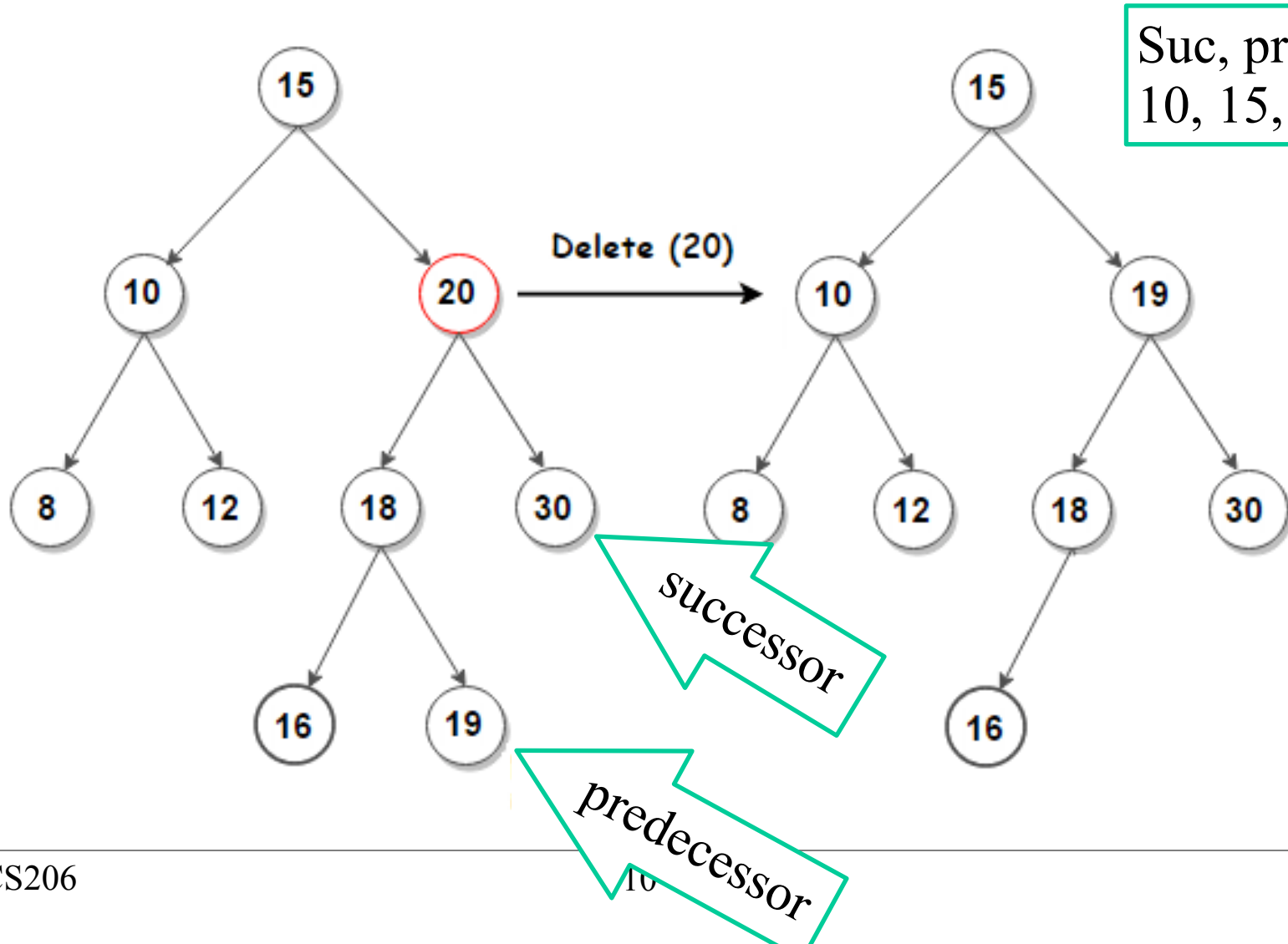
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# Two Children

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- Replace with in-order predecessor or in-order successor
- in-order predecessor
  - rightmost child in left subtree
  - max-value child in left subtree
- in-order successor
  - leftmost child in right subtree
  - min-value child in right subtree

# 2 child replacement



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# remove pseudocode

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```
boolean remove(element)
    return removeUtil(element, root, null);

boolean removeUtil(element, node, parent)
    if (node==null) return false;
    if (node.payload>element)
        removeUtil(element, node.left, node);
    else if (node.payload<element)
        removeUtil(element, node.right, node);
    else
```

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# remove pseudocode 2

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```
// found the node to delete
if (node.right==null && node.left==null)
    // at a leaf
    parent.remove(node)
    return true
if (node.right==null)
    // one descendent on left
    attach node.left to parent
    return true;
if (node.left==null)
    // one descendent on right
    attach node.right to parent
    return true;
```

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# remove pseudocode 3

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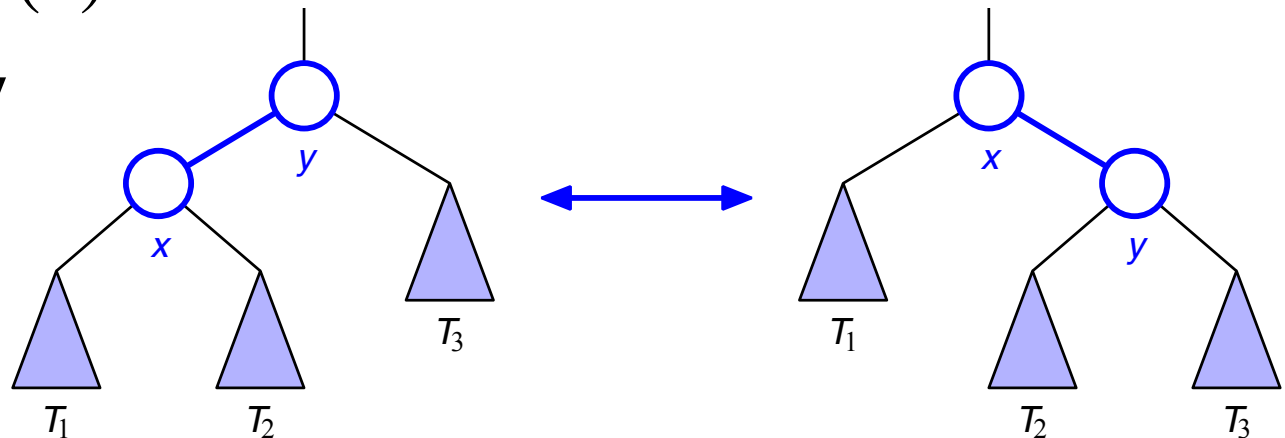
```
// two children
successorNode = inorderSucessor(node.right)
node.payload=successorNode.payload
removeUtil(successorNode.payload, node.right
node);
return true;
```

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# Balanced Search Trees

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- A variety of algorithms augment a standard BST with occasional operations to reshape, reduce height and maintain balance.
- General approach: Rotation — moves a child to be above its parent,
  - ideally  $O(1)$
  - certainly  $O(\lg n)$



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# AVL Trees

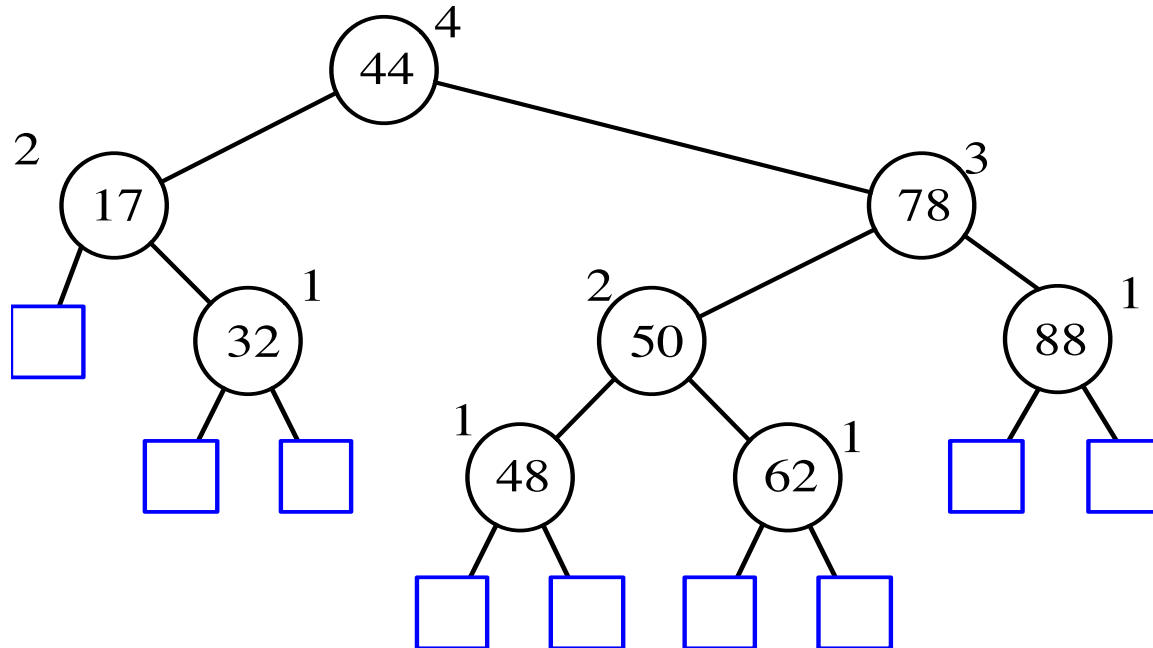
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- Height-balance property
  - For every internal node, the `avlHeight` of the two children differ by at most 1
    - `avlHeight` = max distance from null endpoint
- Any binary tree satisfying the height-balance property is an AVL tree
- A height-balanced tree has height  $O(\lg n)$ 
  - max height is provably  $1.44 * \lg(n)$

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# AVL Tree Example

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# Insertion

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- Maintain with each node the `avlHeight`.
- On insertion, first recur down through tree to insert.
- Then as you unwind recursion, update the `avlHeight` of each node.
- If height changes, check the height of other child
  - if not in balance then fix

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# Insertion code to maintain height

(the only code today!!!)

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```
private class Node {
    Comparable<E> element;
    int avlHight;
    Node right;
    Node left;

    public Node(Comparable<E> e) {
        avlHight = 1;
        element=e;
        right=null;
        left=null;
    }
}
```

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# More insertion (pseudo)code

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```
int insertUtil(node, element):  
    if element==node.payload  
        return -1;  
  
    avlD=2;  
    if node.payload > element:  
        if node.left==null  
            node.left=new Node(payload)  
        else  
            avlD = 1+insertUtil(node.left,element);  
    else  
        // same but for right  
  
    node.avlHieght = greater of avlD and  
                    node.avlHeight  
  
    return node.avlHeight
```

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# Fixing height imbalances

## Rotation!!

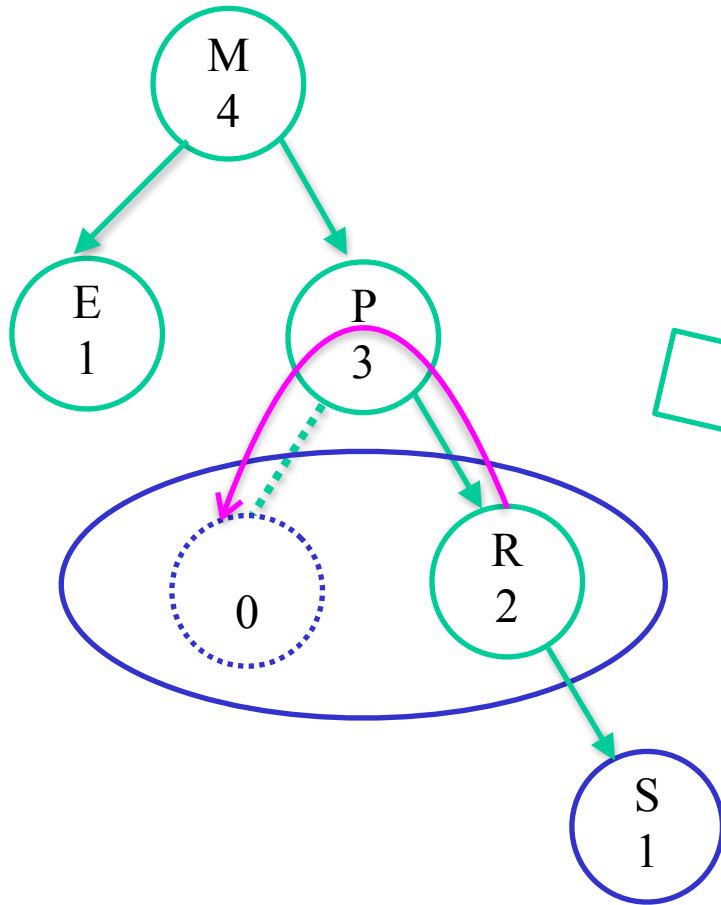
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- Two types of rotation
- Single
  - left subtree of left node causes imbalance
  - right subtree of right node causes imbalance
- Double
  - right subtree of left node causes imbalance
  - left subtree of right node causes imbalance
    - The first rotation of a double puts the tree into position for a single rotation!

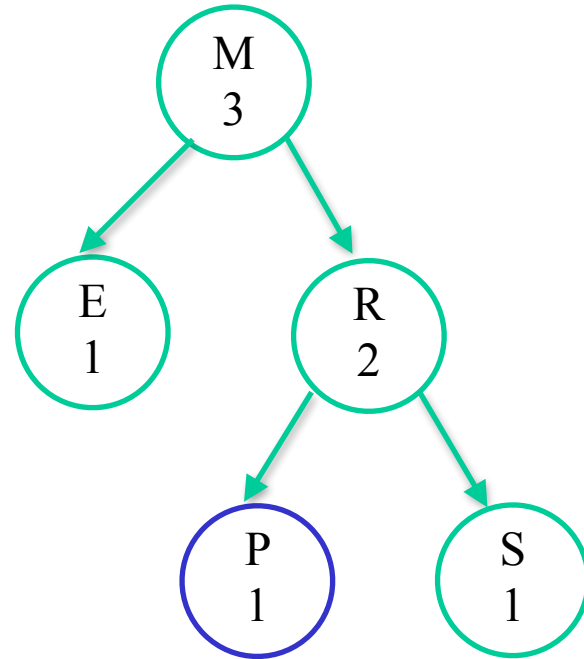
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# Single Rotation

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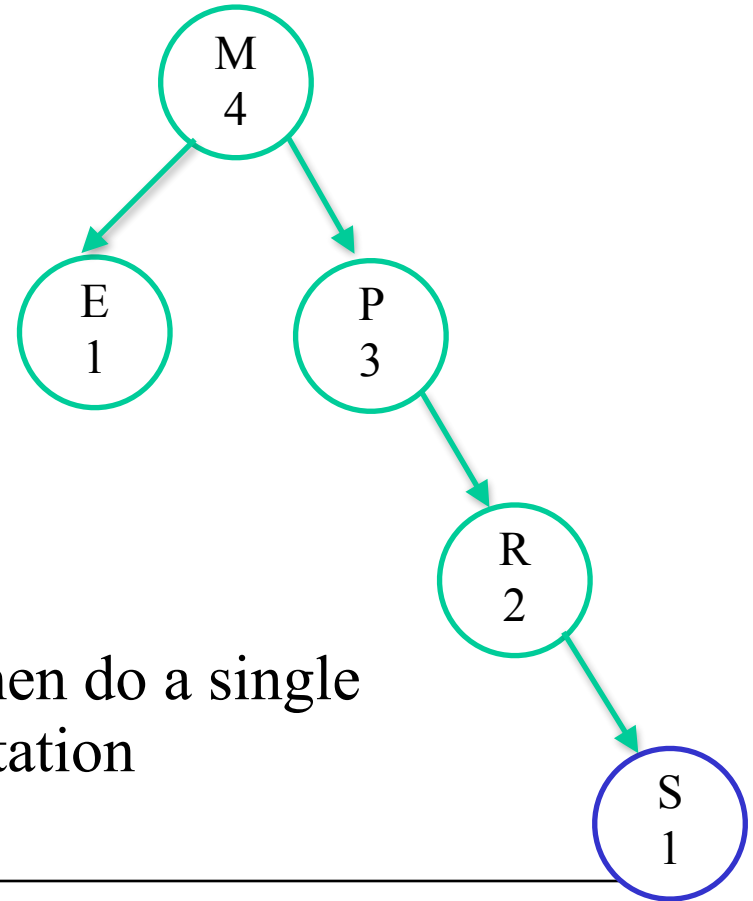
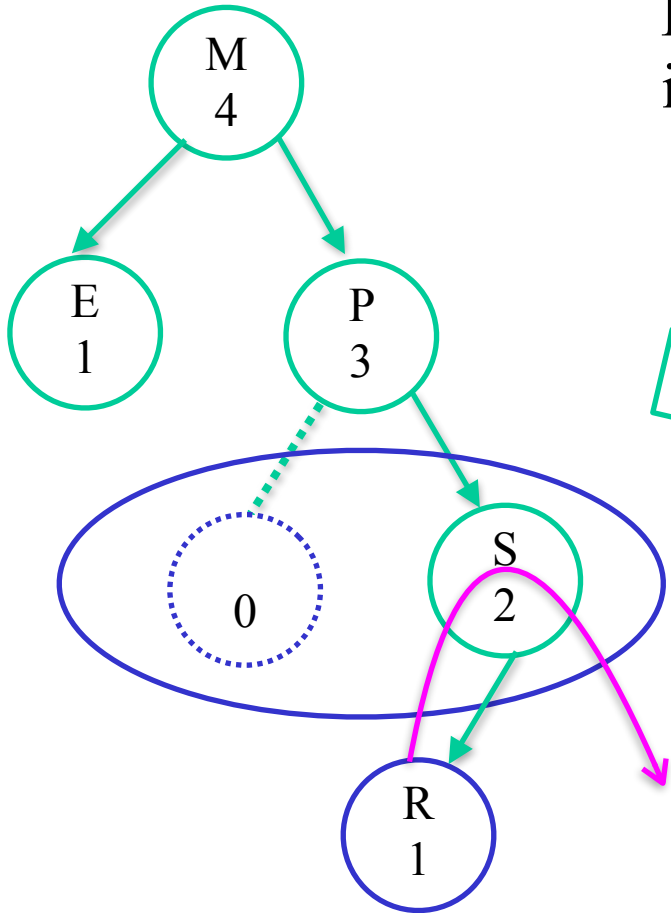


Rotate across parent at the lowest imbalance



# Double Rotation

First rotate across the point imbalance, This shifts from R-L to R-R



Then do a single rotation

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# Groups

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- Given the following data so the tree after each operation, while keeping the tree balanced using AVL
- for deletions, always delete using inorder predecessor.
- iXXX == insert XXX into tree
- dXXX == delete XXX from tree
  
- i1024, i512, i256, i128, i64, i32, i16, i750, i875, d125, d32

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# Homework!!!!

Its more than just a lab

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Show the tree after each insertion / deletion  
Before and after AVL rotation

insert	100
insert	200
insert	300
insert	400
insert	500
insert	600
insert	700
insert	800
insert	900
insert	750
insert	1000
insert	850
delete	400
delete	300
delete	200
delete	700
delete	500