Artificial Intelligence

Adversarial Search (Game Playing)

Chapter 5

Adapted from materials by Tim Finin, Marie desJardins, and Charles R. Dyer

Outline

- Game playing
 - State of the art and resources
 - Framework
- Game trees
 - Minimax
 - Alpha-beta pruning
 - Adding randomness

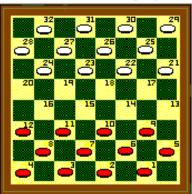
State of the art

- How good are computer game players?
 - Chess:
 - Deep Blue beat Gary Kasparov in 1997
 - Garry Kasparav vs. Deep Junior (Feb 2003): tie!
 - Kasparov vs. X3D Fritz (November 2003): tie!
 - **Checkers**: Chinook (an AI program with a *very large* endgame database) is the world champion. Checkers has been solved exactly it's a draw!
 - Go: Computer players are decent, at best
 - **Bridge**: "Expert" computer players exist (but no world champions yet!)
- Good place to learn more: http://www.cs.ualberta.ca/~games/

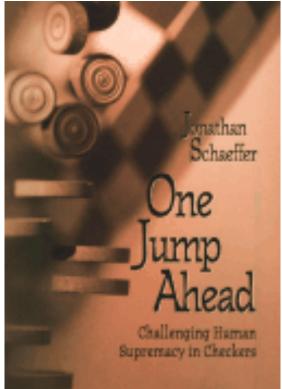
Chinook

- Chinook is the World Man-Machine Checkers Champion, developed by researchers at the University of Alberta.
- It earned this title by competing in human tournaments, winning the right to play for the (human) world championship, and eventually defeating the best players in the world.
- Visit http://www.cs.ualberta.ca/~chinook/ to play a version of Chinook over the Internet.
- The developers have fully analyzed the game of checkers and have the complete game tree for it.
 - Perfect play on both sides results in a tie.
- "One Jump Ahead: Challenging Human Supremacy in Checkers" Jonathan Schaeffer, University of Alberta (496 pages, Springer. \$34.95, 1998).

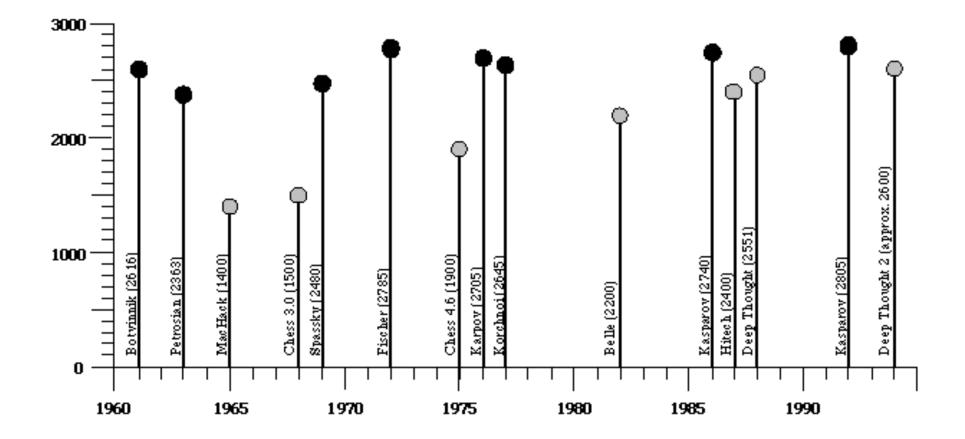
The board set for play



Red to play



Ratings of human and computer chess champions





Typical case

- 2-person game
- Players alternate moves
- Zero-sum: one player's loss is the other's gain
- **Perfect information**: both players have access to complete information about the state of the game. No information is hidden from either player.
- No chance (e.g., using dice) involved
- Examples: Tic-Tac-Toe, Checkers, Chess, Go, Nim, Othello
- Not: Bridge, Solitaire, Backgammon, ...

How to play a game

- A way to play such a game is to:
 - Consider all the legal moves you can make
 - Compute the new position resulting from each move
 - Evaluate each resulting position and determine which is best
 - Make that move
 - Wait for your opponent to move and repeat
- Key problems are:
 - Representing the "board"
 - Generating all legal next boards
 - Evaluating a position

Evaluation function

- Evaluation function or static evaluator is used to evaluate the "goodness" of a game position.
 - Contrast with heuristic search where the evaluation function was a non-negative estimate of the cost from the start node to a goal and passing through the given node
- The zero-sum assumption allows us to use a single evaluation function to describe the goodness of a board with respect to both players.
 - f(n) >> 0: position n good for me and bad for you
 - $f(n) \ll 0$: position n bad for me and good for you
 - f(n) near 0: position n is a neutral position
 - $\mathbf{f}(\mathbf{n}) = +\mathbf{infinity}$: win for me
 - $\mathbf{f}(\mathbf{n}) = -\mathbf{infinity}$: win for you

Evaluation function examples

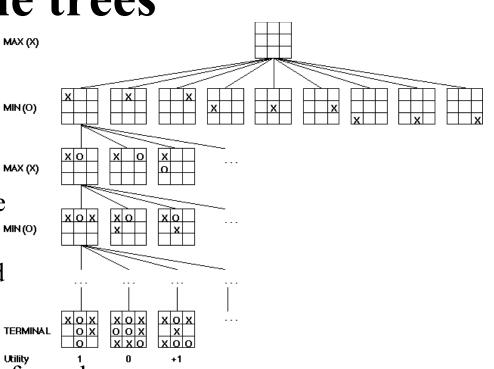
- Example of an evaluation function for Tic-Tac-Toe:
 f(n) = [# of 3-lengths open for me] [# of 3-lengths open for you]
 where a 3-length is a complete row, column, or diagonal
- Alan Turing's function for chess
 - f(n) = w(n)/b(n) where w(n) = sum of the point value of white's pieces and b(n) = sum of black's
- Most evaluation functions are specified as a weighted sum of position features:

 $f(n) = w_1^* feat_1(n) + w_2^* feat_2(n) + ... + w_n^* feat_k(n)$

- Example features for chess are piece count, piece placement, squares controlled, etc.
- Deep Blue had over 8000 features in its evaluation function

Game trees

- Problem spaces for typical games are represented as trees
- Root node represents the current board configuration; player must decide the best single move to make next
- Static evaluator function rates a board position. f(board) = real number with f>0 "white" (me), f<0 for black (you)
- Arcs represent the possible legal moves for a player
- If it is **my turn** to move, then the root is labeled a "**MAX**" node; otherwise it is labeled a "**MIN**" node, indicating **my opponent's turn**.
- Each level of the tree has nodes that are all MAX or all MIN; nodes at level i are of the opposite kind from those at level i+1

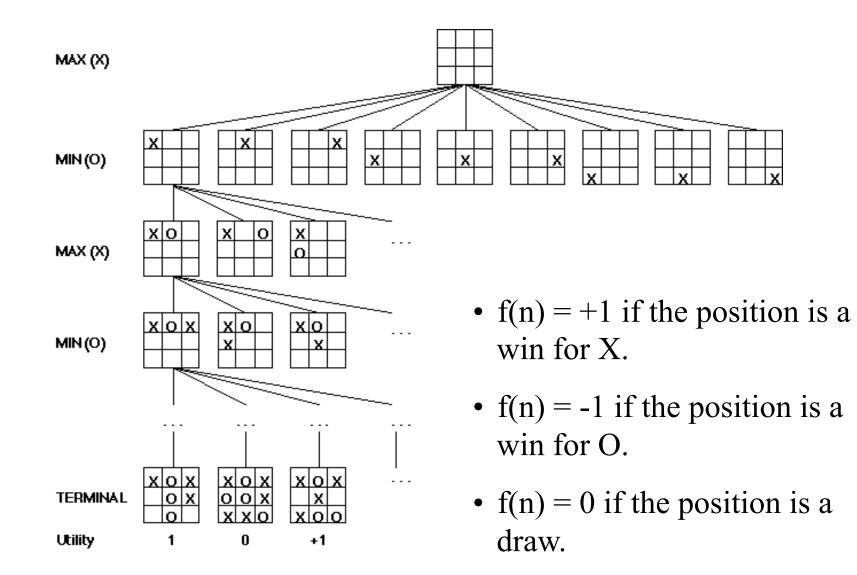


Minimax procedure

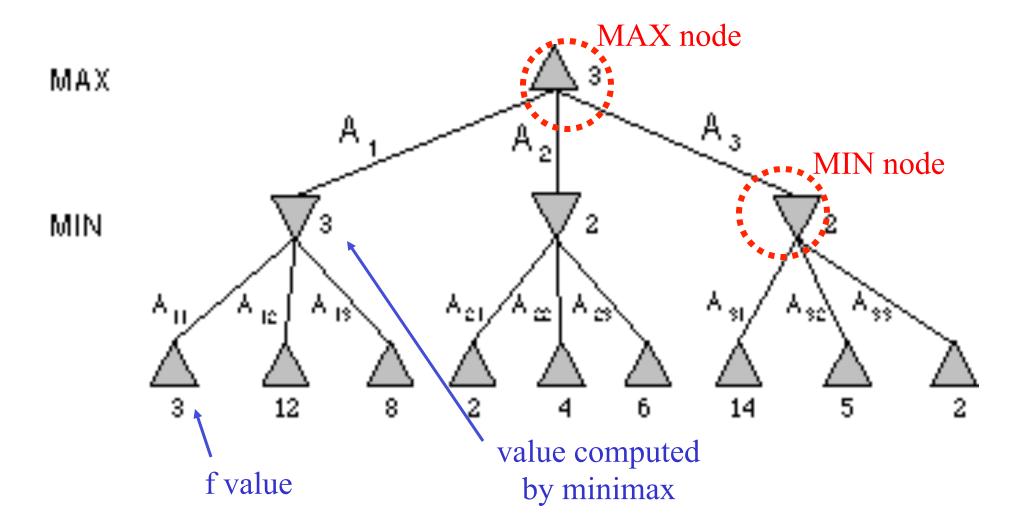
- Create start node as a MAX node with current board configuration
- Expand nodes down to some **depth** (a.k.a. **ply**) of lookahead in the game
- Apply the evaluation function at each of the leaf nodes
- "Back up" values for each of the non-leaf nodes until a value is computed for the root node
 - At MIN nodes, the backed-up value is the minimum of the values associated with its children.
 - At MAX nodes, the backed-up value is the maximum of the values associated with its children.
- Pick the operator associated with the child node whose backed-up value determined the value at the root

Minimax Algorithm This is the move selected by minimax Static evaluator value MAX MIN

Partial Game Tree for Tic-Tac-Toe

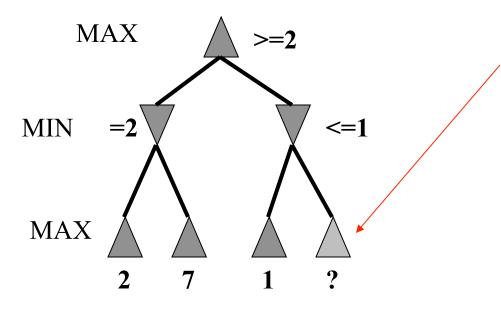


Minimax Tree



Alpha-beta pruning

- We can improve on the performance of the minimax algorithm through **alpha-beta pruning**
- Basic idea: *"If you have an idea that is surely bad, don't take the time to see how truly awful it is."* -- Pat Winston

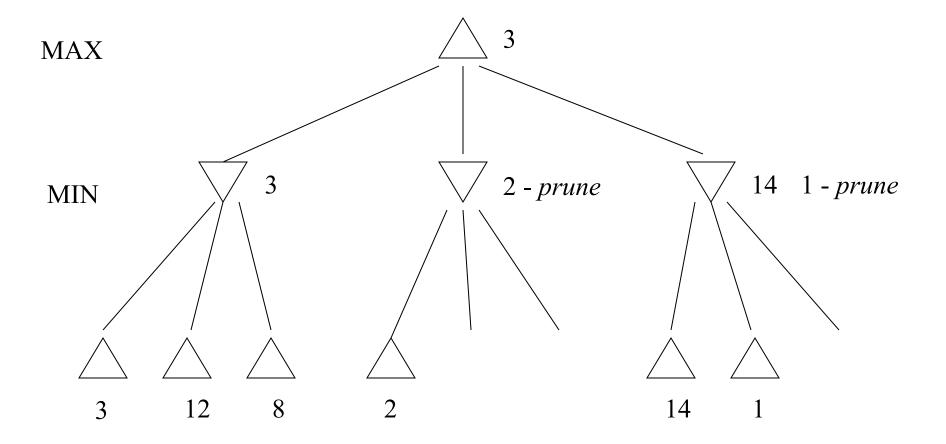


- We don't need to compute the value at this node.
- No matter what it is, it can't affect the value of the root node.

Alpha-beta pruning

- Traverse the search tree in depth-first order
- At each MAX node n, alpha(n) = maximum value found so far
- At each MIN node n, beta(n) = minimum value found so far
 Note: The alpha values start at -infinity and only increase, while beta values start at +infinity and only decrease.
- Beta cutoff: Given a MAX node n, cut off the search below n (i.e., don't generate or examine any more of n's children) if alpha(n) >= beta(i) for some MIN node ancestor i of n.
- Alpha cutoff: stop searching below MIN node n if beta(n) <= alpha(i) for some MAX node ancestor i of n.

Alpha-beta example



Alpha-beta algorithm

```
function MAX-VALUE (state, \alpha, \beta)
     ;; \alpha = best MAX so far; \beta = best MIN
if TERMINAL-TEST (state) then return UTILITY(state)
v := -\infty
for each s in SUCCESSORS (state) do
    v := MAX (v, MIN-VALUE (s, \alpha, \beta))
    if v \geq \beta then return v
    \alpha := MAX (\alpha, v)
end
return v
function MIN-VALUE (state, \alpha, \beta)
if TERMINAL-TEST (state) then return UTILITY(state)
V := ∞
for each s in SUCCESSORS (state) do
    v := MIN (v, MAX-VALUE (s, \alpha, \beta))
    if v \leq \alpha then return v
    \beta := MIN (\beta, v)
end
return v
```

Effectiveness of alpha-beta

- Alpha-beta is guaranteed to compute the same value for the root node as computed by minimax, with less or equal computation
- Worst case: no pruning, examining b^d leaf nodes, where each node has b children and a d-ply search is performed
- Best case: examine only $(2b)^{d/2}$ leaf nodes.
 - Result is you can search twice as deep as minimax!
- **Best case** is when each player's best move is the first alternative generated
- In Deep Blue, they found empirically that alpha-beta pruning meant that the average branching factor at each node was about 6 instead of about 35!

Games of chance

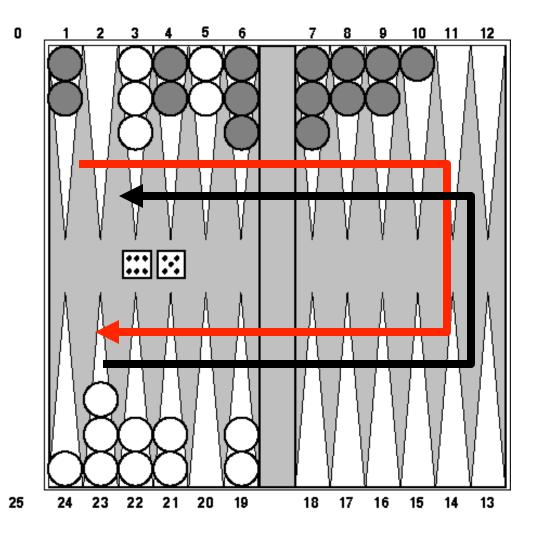
• Backgammon is a two-player game with **uncertainty**.

•Players roll dice to determine what moves to make.

•White has just rolled *5 and 6* and has four legal moves:

• 5-10, 5-11 •5-11, 19-24 •5-10, 10-16 •5-11, 11-16

•Such games are good for exploring decision making in adversarial problems involving skill and luck.

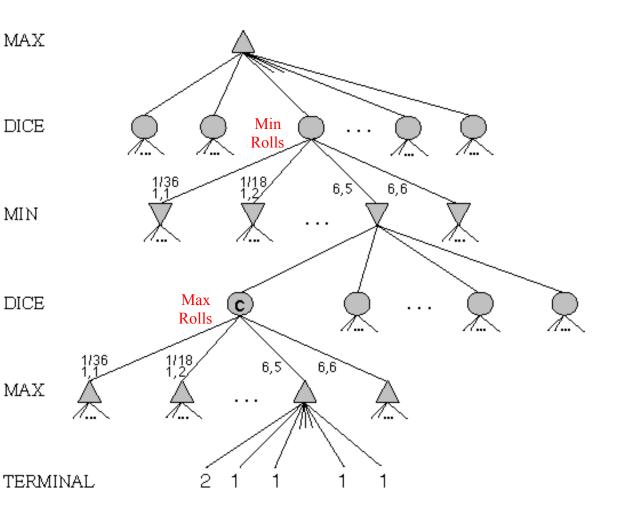


Game trees with chance nodes

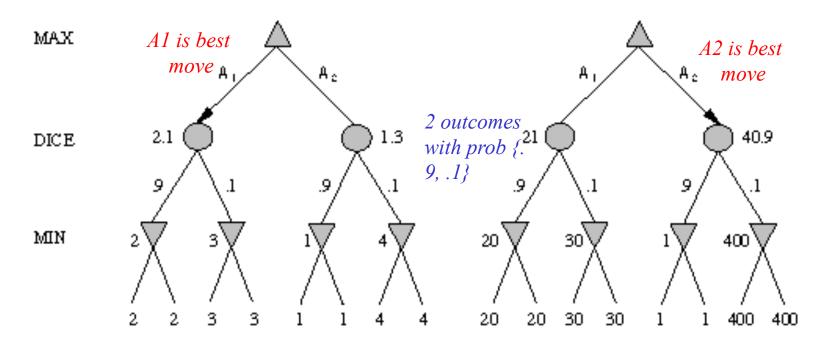
- Chance nodes (shown as circles) represent random events
- For a random event with N outcomes, each chance node has MAX N distinct children; a probability is associated with each
- (For 2 dice, there are 21 distinct outcomes)
- Use minimax to compute values MIN for MAX and MIN nodes
- Use **expected values** for chance DICE nodes
- For chance nodes over a max node, as in C:

 $expectimax(C) = \sum_{i} (P(d_{i}) * maxvalue(i))$

• For chance nodes over a min node: expectimin(C) = $\sum_{i} (P(d_i) * minvalue(i))$



Meaning of the evaluation function



- Dealing with probabilities and expected values means we have to be careful about the "meaning" of values returned by the static evaluator.
- Note that a "relative-order preserving" change of the values would not change the decision of minimax, but could change the decision with chance nodes.
- Linear transformations are OK